

GE863-PRO³ Linux Software User Guide

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1 Introduction

1.1 Scope

This user guide serves the following purpose:

- Provides details about the GE863-PRO³ software architecture
- Describes how software developers can use the functions of Linux device drivers to configure, manage and use GE863-PRO³ hardware resources and system peripherals

1.2 Audience

This User Guide is intended for software developers who develop applications on the ARM processor of GE863-PRO³ module.

1.3 Contact Information, Support

Our aim is to make this guide as helpful as possible. Keep us informed of your comments and suggestions for improvements.

For general contact, technical support, report documentation errors and to order manuals, contact Telit's Technical Support Center at:

TS-EMEA@telit.com or <http://www.telit.com/en/products/technical-support-center/contact.php>

Telit appreciates feedback from the users of our information.

1.4 Open Source Licenses

Linux system is made up of many Open Source device drivers licensed as follows:

GNU GENERAL PUBLIC LICENSE
Version 2, June 1991

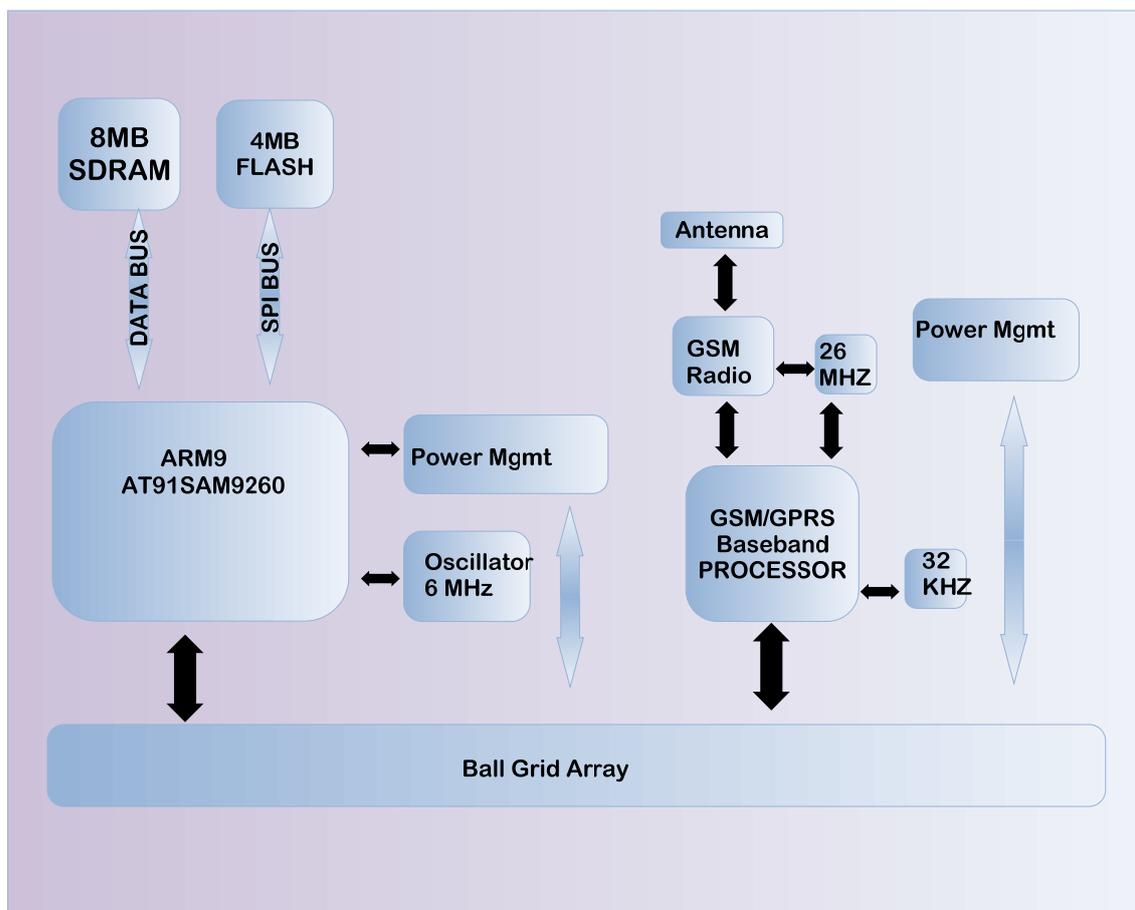
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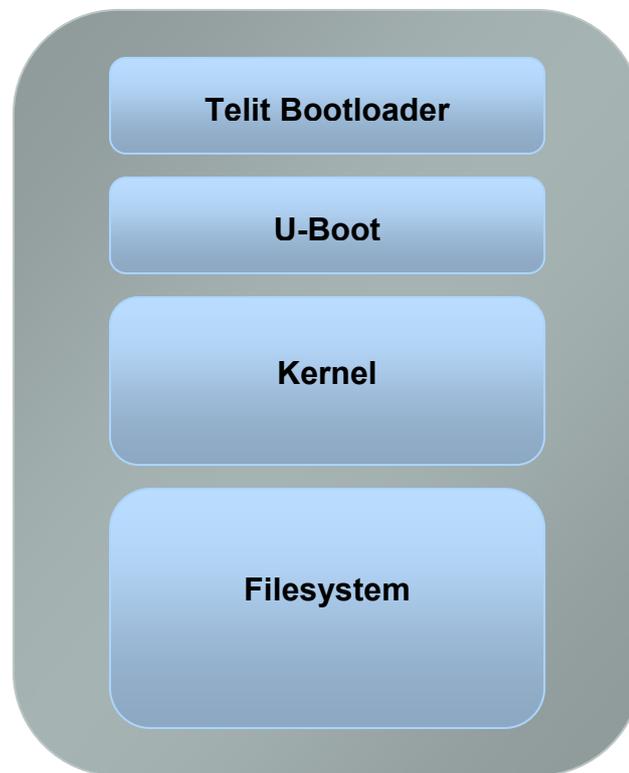


2 GE863-PRO³ Architecture

2.1 Hardware

The GE863-PRO³ is an innovation to the quad-band, RoHS compliant GE863 product family which includes a powerful ARM9™ processor core exclusively dedicated to customer applications. The GE863-PRO³ incorporates much of the necessary hardware for communicating microcontroller solutions, including the critical element of memory, significant simplification of the bill of material, vendor management, and logistics effort are achieved. Below there is a simple plot and a list of the GE863-PRO³ key elements:





2.2.1 Telit Bootloader

Telit Bootloader is a small binary used for hardware-related management: its main task is to load and start Telit customized U-Boot. For further details refer to document [2] .

2.2.2 Telit customized U-boot

U-Boot is the Open Source "universal" cross-platform bootloader supporting hundreds of embedded boards and a wide variety of microcontrollers (the Atmel AT91SAM9260 included).

Some of the features provided by U-Boot are:

- Hardware initialization.
- Providing boot parameters for the Linux kernel.
- Starting the Linux kernel.
- Reading and writing memory locations.
- Uploading new binary images to the board's RAM.
- Copying binary images from RAM to FLASH memory.
- Storing environment variables which can be used to configure the system.

The GE863-PRO³ has a customized version of U-Boot to get the most out of the board. For further details refer to document [2] .



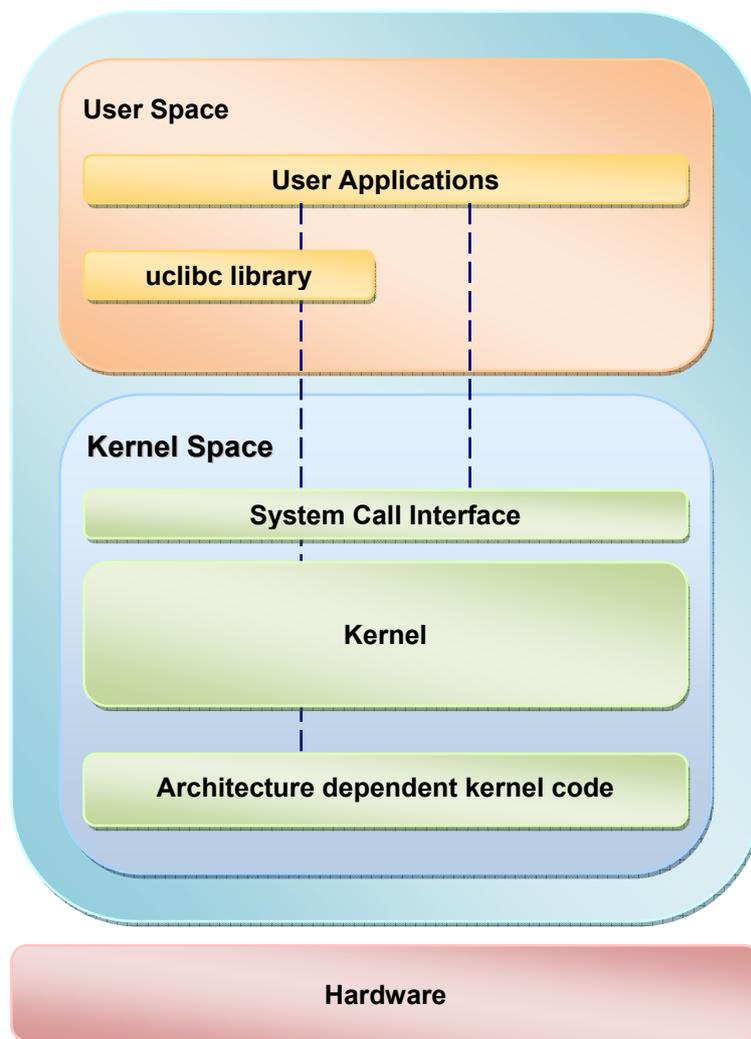
2.2.3 Linux kernel

2.2.3.1 Overview

The kernel is the central part of the GNU/Linux operating system: its main task is to manage system's resources in order to make the hardware and the software communicate. A kernel usually deals with process management (including inter-process communication), memory management and device management.

The Linux kernel belongs to the family of Unix-like operating system kernel; created in 1991, it has been developed by a huge number of contributors worldwide during these years, becoming one of the most common and versatile kernel for embedded systems.

Below there is a schematic representing, from a high level perspective, the architecture of a GNU/Linux operating system:



Two regions can be identified:

- 1) User space: where the user applications are executed.
- 2) Kernel space: where the kernel (with all its components such as device drivers) works.

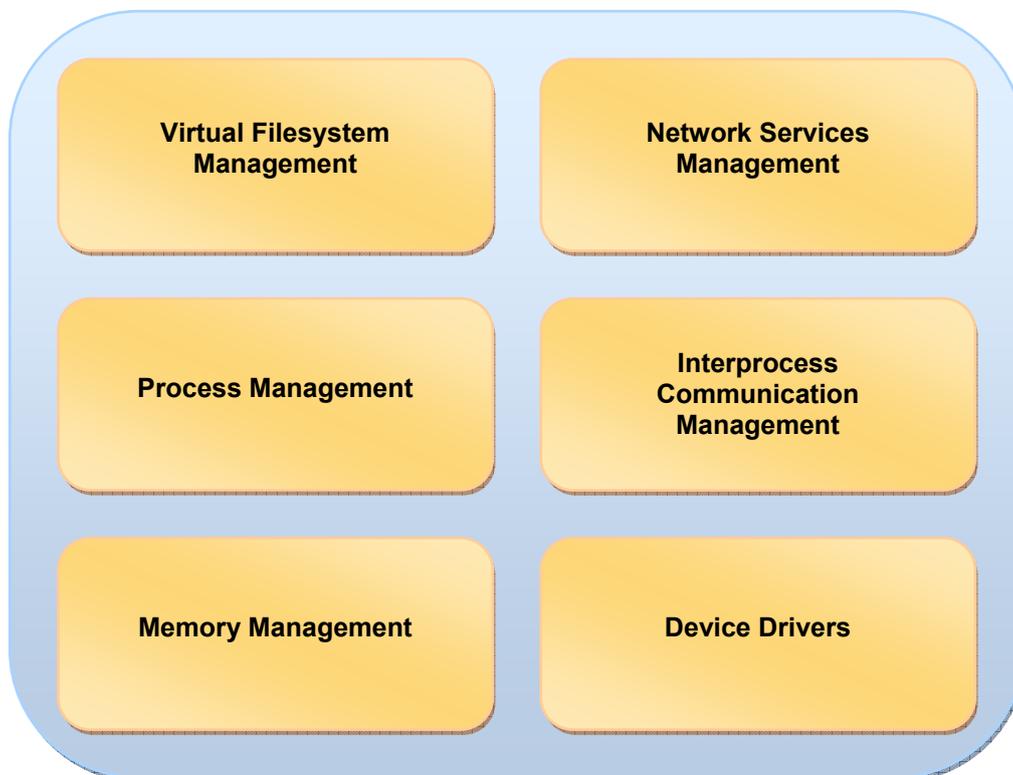
These two regions are separated and have different memory address spaces; there are several methods for user/kernel interaction:

- Using the System Call Interface that connects user to the kernel and provides the mechanism to communicate between the user-space application and the kernel through the C library.
- Using kernel calls directly from application code leaping over the C library.
- Using the virtual filesystem.

The ordinary C library in Linux system is the glibc. Uclibc is a C library mainly targeted for developing embedded Linux systems; despite being much smaller than the glibc it has almost all its features (including shared libraries and threading), making easy to port applications from glibc to uclibc.

The Linux kernel architecture-independent code stays on the top of platform specific code for the GE863-PRO³ board: this code is allows exploiting all hardware features of the GE863-PRO³.

Inside the kernel we can find, among the others, the following fundamental components:

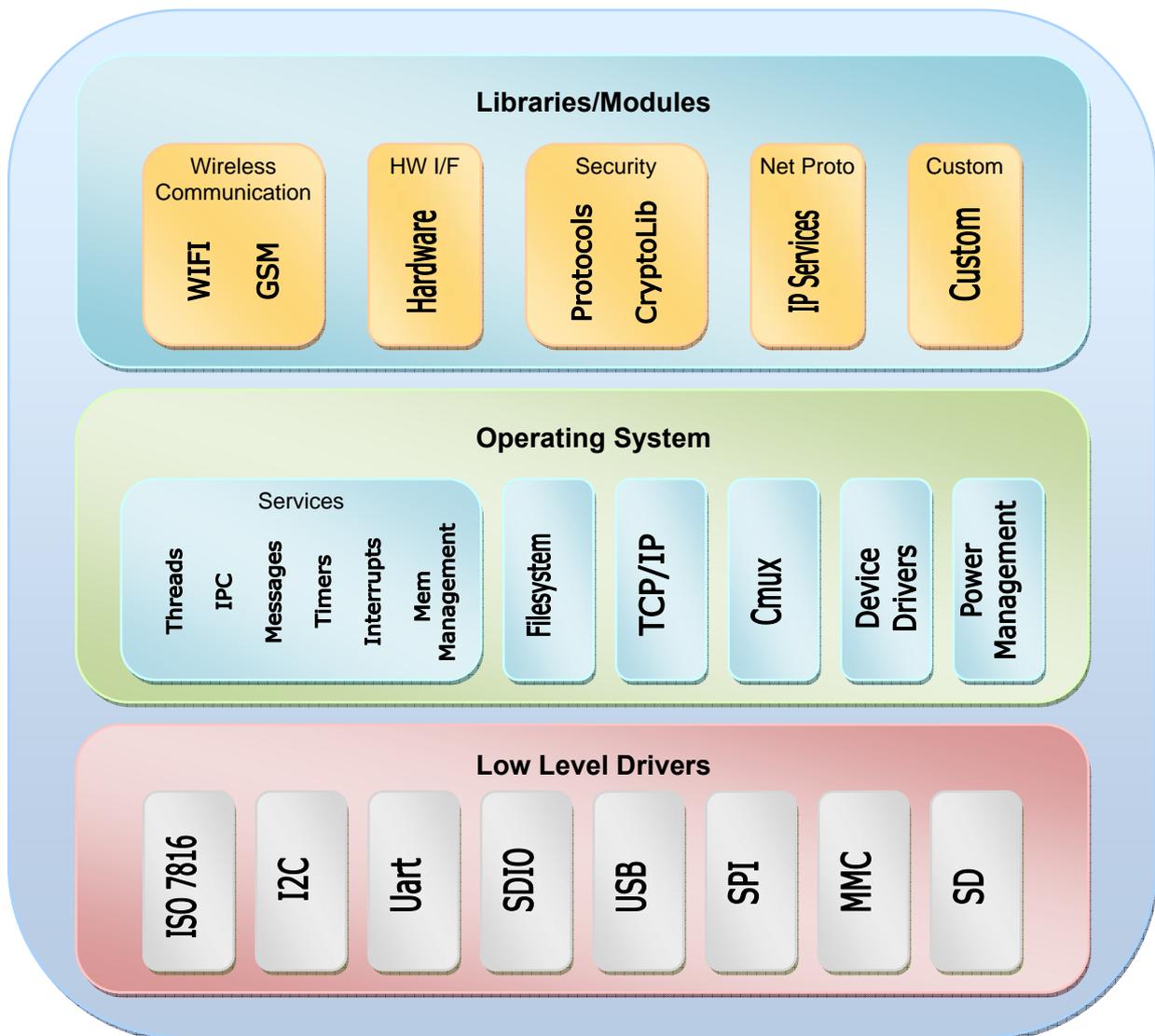


The Linux kernel is a monolithic one (i.e. all OS services run along with the main kernel thread, thus also residing in the same memory area), but it has the capability to dynamically load/unload some of its components called modules.

A kernel module is a compiled piece of code which can be dynamically linked to the kernel when is needed (becoming part of the kernel as the other normal kernel code) then removed when it is no longer needed; modules are mainly used for device drivers.

2.2.3.2 The GE863-PRO³ Linux kernel

The Linux kernel on the GE863-PRO³ is based on version number 2.6.24. Below there is a picture showing the kernel main components:



The GE863-PRO³ Linux kernel comes with some of its features linked statically; while others are compiled as modules (see the table below).

At the moment, available drivers are:

Driver Name	Module	Module Name
GPIO	N	N/A
MMC/SD	N	N/A
Ethernet	Y	macb
Usb Human Interface Device	Y	usbhid
Support for Host-side USB	Y	ohci-hcd
Support for USB Mass Storage	Y	usb-storage
USB Ethernet Gadget	Y	g-ether
I2C	Y	i2c-core i2c-algo-bit i2c-dev i2c-gpio
SPI	Y	Spidev

2.2.4 Filesystem

2.2.4.1 Overview

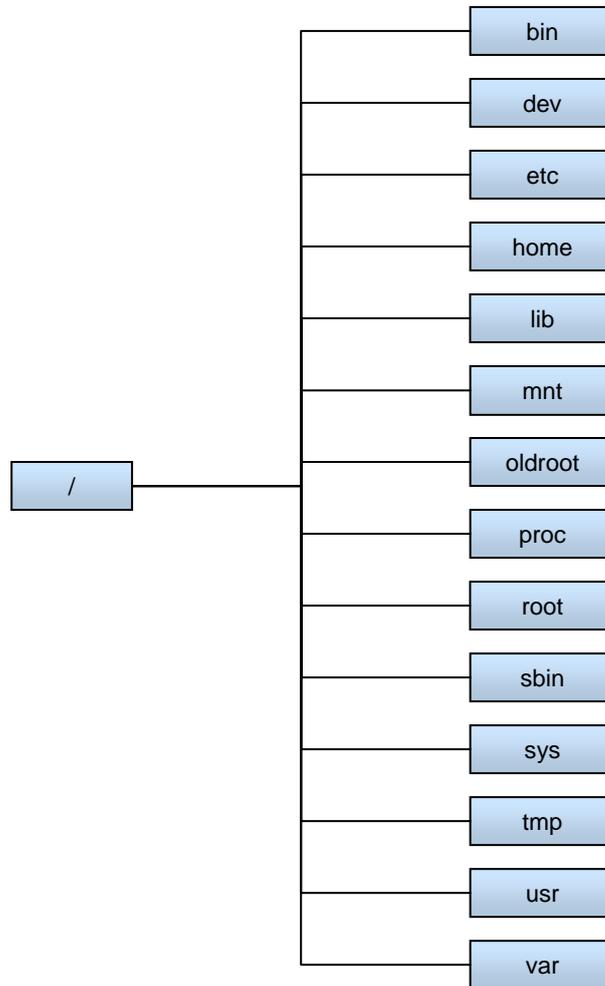
A filesystem is the entity where the files are placed logically for storage and retrieval. The filesystem also specifies conventions for naming files (e.g. maximum number of characters for the file name). Usually the file system has a hierarchical-tree structure: files are placed in directories inside this tree structure.

The GE863-PRO³ has the JFFS2, a log-structured file system specifically for use on flash devices in embedded systems. It implements file compression with the following formats: Zlib, LZO, Real time.

2.2.4.2 The filesystem structure

The image below shows the filesystem root tree.





`/bin`

This directory contains essential tools and other programs (so called binaries): note that, by default in the GE863-PRO³, a most of these files are symbolic link that depend on Busybox (see paragraph 2.2.4.3). Moreover the Busybox binary is also placed under this directory.

`/dev`

This directory contains files representing the system's various hardware devices. Here, for example, can be found all the devices representing the various GPIOs.

`/etc`

This directory contains system configuration files, startup files other. Here, for example, can be found the file `fstab` used for mounting the various devices in the filesystem.

`/home`

This directory contains the home directories of the various users.

`/lib`



This directory contains shared libraries required by programs; for example, here there can be found the uclibc libraries. Moreover in its subdirectories there are the various loadable kernel modules (for example the modules for USB Mass Storage Devices): they can be easily recognized by the .ko extension.

`/mnt`

This directory contains the mount points¹ for the various devices.

`/oldroot`

By default this directory is empty. It is used by the system for the startup process.

`/proc`

This directory contains virtual folders and files which represent the current state of the kernel; for example here you can find info about the cpu (type `cat /proc/cpuinfo`).

`/root`

This directory is the home of the user root. By default it is empty.

`/sbin`

This directory contains executable files are mainly used by the root user for administration task. By default in the GE863-PRO³, most of these files are symbolic links that depend on Busybox (see paragraph 2.2.4.3).

`/sys`

This directory contains system files used for device configuration.

`/tmp`

This directory contains temporary files.

`/usr`

This directory contains files and directories related to user tools and applications. By default in the GE863-PRO³ some of these files are symbolic links that depend on Busybox (see paragraph 2.2.4.3).

`/var`

This directory contains variable data files.

¹ The mount point is the location in the operating system's directory structure where a mounted file system appears



2.2.4.3 Busybox

Inside the filesystem in the directory `/bin` there is the Busybox. Busybox gives the user a set of commands useful to interact with the system; with it the user can carry out normal tasks such as copying files (`cp` command), listing directories (`ls` command), deleting files (`rm` command) etc. etc.

BusyBox combines tiny versions of many common UNIX utilities into a single small executable. It provides replacements for most of the utilities usually found in a GNU/Linux system. Note that the utilities in BusyBox generally have fewer options than their full-featured GNU counterparts; however, the options that are included usually satisfy the user's needs and provide a fairly complete POSIX environment.

Busybox is composed of:

- A single binary placed in `/bin`, which is the real executor of the various utilities.
- A number of symbolic links through all the filesystem, having the name of the various utilities (`cp`, `ls`, `kill` etc.) which refer to the Busybox binary.

To discover the currently defined commands for the GE863-PRO³, open a terminal (refer to 3.1 for further details) and type:

```
# cd /bin
# ./busybox
```

To see help content for any command type in the terminal:

```
# <command name> --help
```



3 System Startup

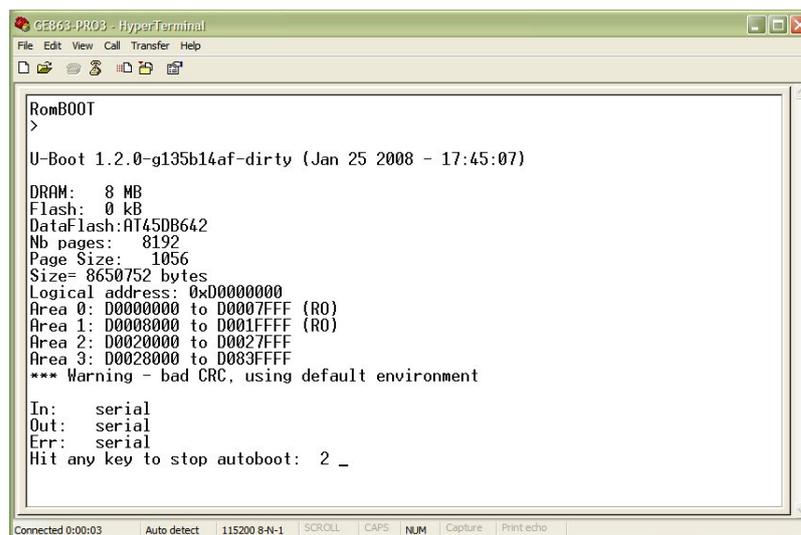
3.1 Startup process

Connect the GE863-PRO³ to your host system via serial cable (use Debug port of the EVK, for further details refer to document [3]). In your host system open a terminal program (such as Hyper Terminal) and use the following parameters for the connection:

Bits per second: 115200
 Data bits: 8
 Parity: None
 Stop bits: 1
 Flow Control: None

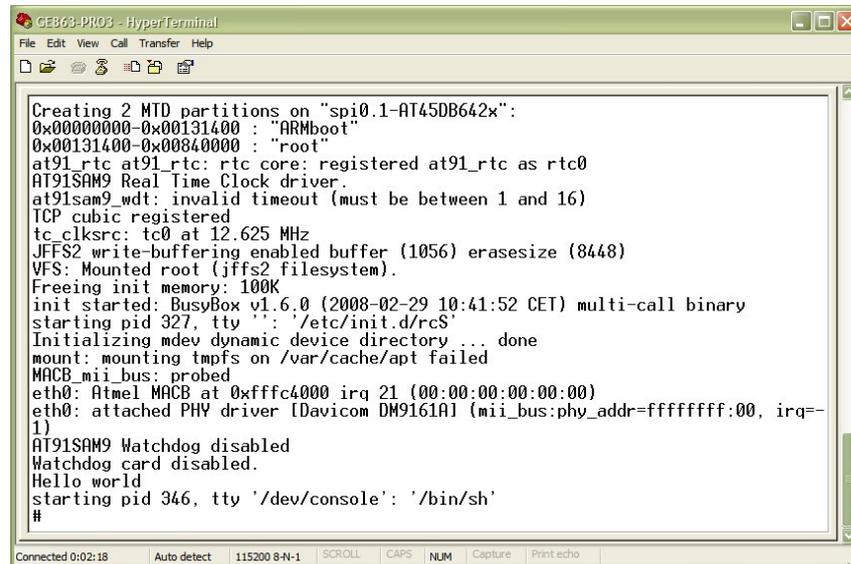
After turning on the GE863-PRO³ the following operations occur:

- The Telit Bootloader (refer to 2.2.1) starts and after initializing the hardware it loads the U-Boot image from flash into RAM and runs it.
- Telit customized U-Boot (refer to 2.2.2) starts and in the terminal you should see a countdown. If you wish to enter in U-Boot command mode (for further details refer to document [2]) press any key, otherwise the startup process will continue within a few seconds: U-Boot loads the kernel image in RAM then starts it.



- The kernel image decompresses itself and starts. After the initial sequence the control is passed to the shell (refer to 3.2) which can be used to interact with the target. At this point the system is ready to be used.





```

GE863-PRO3 - HyperTerminal
File Edit View Call Transfer Help
Creating 2 MTD partitions on "spi0.1-AT45DB642x":
0x00000000-0x00131400 : "ARMboot"
0x00131400-0x00840000 : "root"
at91_rtc at91_rtc: rtc core: registered at91_rtc as rtc0
AT91SAM9 Real Time Clock driver.
at91sam9_wdt: invalid timeout (must be between 1 and 16)
TCP cubic registered
tc clksrc: tc0 at 12.625 MHz
JFFS2 write-buffering enabled buffer (1056) erasesize (8448)
VFS: Mounted root (jffs2 filesystem).
Freeing init memory: 100K
init started: BusyBox v1.6.0 (2008-02-29 10:41:52 CET) multi-call binary
starting pid 327, tty '': '/etc/init.d/rcS'
Initializing mdev dynamic device directory ... done
mount: mounting tmpfs on /var/cache/apt failed
MACB_mii_bus: probed
eth0: Atmel MACB at 0xfffc4000 irq 21 (00:00:00:00:00:00)
eth0: attached PHY driver [Davicom DM9161A] (mii_bus:phy_addr=ffffff:00, irq=-1)
AT91SAM9 Watchdog disabled
Watchdog card disabled.
Hello world
starting pid 346, tty '/dev/console': '/bin/sh'
#
Connected 0:02:18 Auto detect 115200 8-N-1 SCROLL CAPS NUM Capture Print echo

```

3.2 The Linux shell

The Linux shell is a user program that allows you to interact with the target by entering commands from the keyboard; the shell parses the command, executes it and display the output of the command on the screen.

When the target has finished booting, in the terminal the shell prompt will appear:

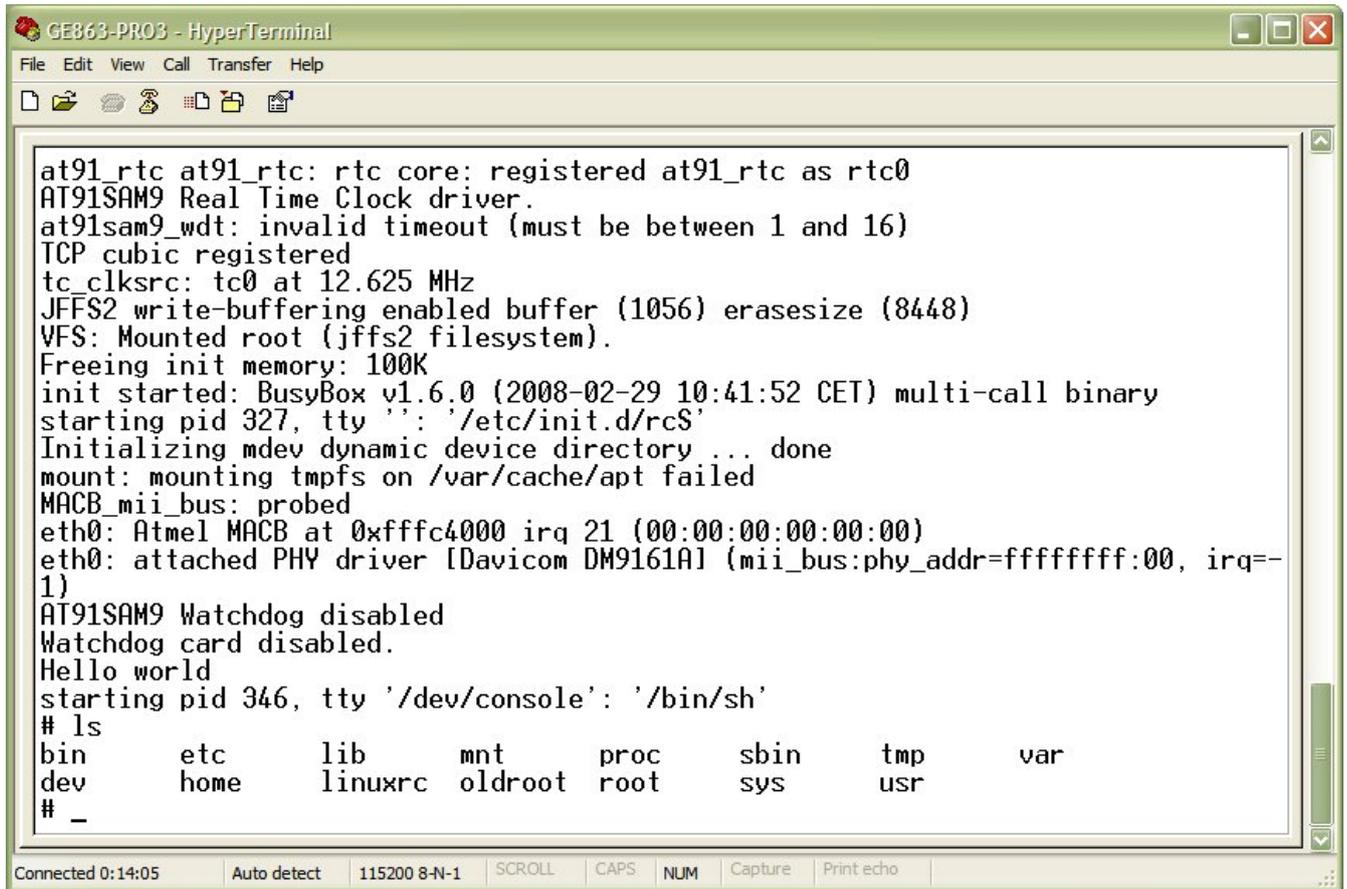
```
#
```

This means that the shell is ready to accept command: you can type in the terminal the command you want to execute; for example:

```
# ls
```

You should see the listing of the “/” directory as in the image below.





```

GE863-PRO3 - HyperTerminal
File Edit View Call Transfer Help
[at91_rtc at91_rtc: rtc core: registered at91_rtc as rtc0
AT91SAM9 Real Time Clock driver.
at91sam9_wdt: invalid timeout (must be between 1 and 16)
TCP cubic registered
tc clksrc: tc0 at 12.625 MHz
JFFS2 write-buffering enabled buffer (1056) erasesize (8448)
VFS: Mounted root (jffs2 filesystem).
Freeing init memory: 100K
init started: BusyBox v1.6.0 (2008-02-29 10:41:52 CET) multi-call binary
starting pid 327, tty '': '/etc/init.d/rc$'
Initializing mdev dynamic device directory ... done
mount: mounting tmpfs on /var/cache/apt failed
MACB_mii_bus: probed
eth0: Atmel MACB at 0xffffc4000 irq 21 (00:00:00:00:00:00)
eth0: attached PHY driver [Davicom DM9161A] (mii_bus:phy_addr=ffffffff:00, irq=-1)
AT91SAM9 Watchdog disabled
Watchdog card disabled.
Hello world
starting pid 346, tty '/dev/console': '/bin/sh'
# ls
bin      etc      lib      mnt      proc     sbin     tmp      var
dev      home    linuxrc  oldroot  root     sys      usr
# _

```

Connected 0:14:05 Auto detect 115200 8-N-1 SCROLL CAPS NUM Capture Print echo

3.3 Loading a module

To load a module type in the terminal:

```
# modprobe <module name>
```

refer to 2.2.3.2 for modules' name.

By default `modprobe` tries to load the module from the directory (and subdirectories) where modules are usually stored. Some of the modules are dependent form others and require loading of other modules first. `modprobe` commits to solve these dependencies before loading the specified module.

To see which modules are currently loaded in the kernel use the `lsmod` command as shown below:



3.5 Downloading a file into GE863-PRO³

There are several ways to download a file in the target. In the following paragraphs you can find explained two of these methods; for further details refer to document [4] .

3.5.1 Downloading a file using the Ethernet connection

Be sure to have the Telit Development Environment correctly installed (with an Ethernet connection up) and coLinux started (refer to document [4] for further details). In the host system go to **Start** → **Telit Development Platform** → **Console**; and the Linux console will be opened. Type:

```
# cd /mnt/windows
```

Now the current directory is where the Windows partition has been mounted.



Once identified the file to be copied, use the cp command in the following form:

```
# cp -r <path where the file is>/<file name> /var/www/apache2-default/
```

For example:

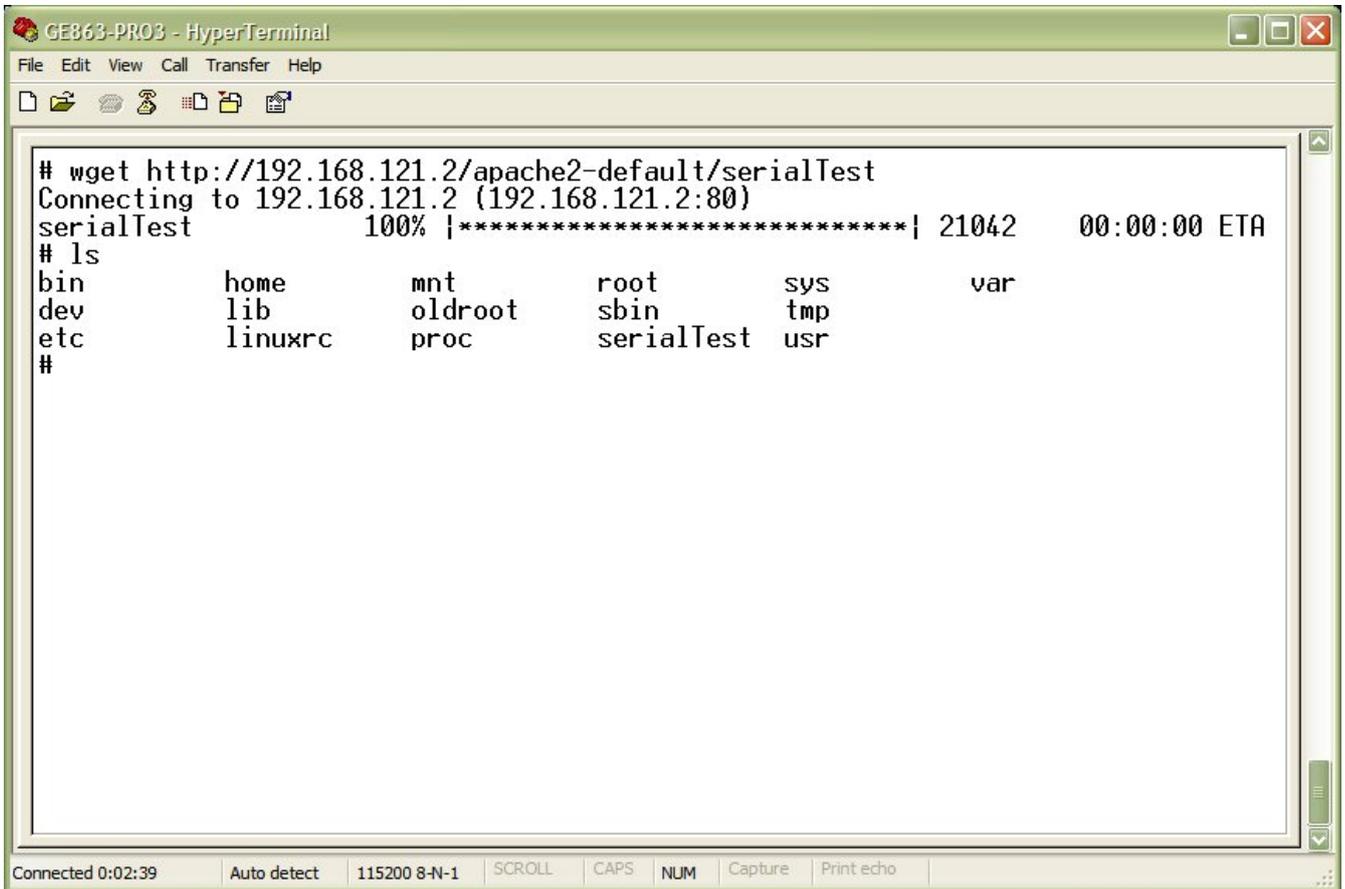
```
# cp -r ARM/binaries/executable /var/www/apache2-default/
```

The file is ready to be uploaded in the target; open the terminal software (as explained in 3.1) and, when the shell is ready, within the directory where you want to place the uploaded file, type:

```
# wget http://192.168.121.2/apache2-default/<file name>
```



You should see an output similar to that in the image below (supposing that the file is called serialTest):



The `ls` command shows that the file has been downloaded.

To remove the file from the development environment type in the Linux console:

```
# rm -rf /var/www/apache2-default/<file name>
```

3.5.2 Downloading a file using an USB mass storage device

Follow the procedure described in 4.8.1 for mounting the device (for example an USB memory key). Then, type in the terminal:

```
# cd /mnt/usbdev0
```

in order to have the current working directory in the USB mass storage device. Now the file can be copied using the `cp` command:

```
# cp <file name> <directory where the file is to be copied>
```



4 Device Drivers

Under Linux OS devices distinguish among three fundamental types: *char*, *block* and *network*. Each kernel module (see paragraph 2.2.3) usually implements one of these types, and thus is classifiable as a *char module*, a *block module*, or a *network module*.

A *char* (character) device is one that can be accessed as a stream of bytes (like a file); serial ports (e.g. `/dev/ttyS0`) are examples of char devices. A char driver usually implements at least the *open*, *close*, *read*, and *write* system calls allowing this type of communication.

Like *char* devices, *block* devices are accessed by filesystem nodes in the `/dev` directory. A block device is a device (e.g., a disk) that can host a filesystem. *Block* and *char* devices differ only in the way data is managed internally by the kernel, thus *block* drivers have a completely different interface to the kernel than *char* drivers.

A *network* interface is a device that is able to exchange data with other hosts and is controlled by a *network* driver. Not being a stream-oriented device, a *network* interface doesn't have a corresponding entry in the filesystem.

Some types of drivers work with additional layers of kernel support functions for dedicated device interface: for example USB devices are driven by a USB module that works with the USB subsystem, but the devices themselves can be char devices (USB serial ports), block devices (USB memory cards), or network devices (USB Ethernet interfaces).

4.1 Serial port

As discussed above serial ports are *char* devices that can be accessed through the following filesystem nodes available on the GE863-PRO³:

<code>ttyS0</code>	used by the shell
<code>ttyS1</code>	available (e.g. used for debugging)
<code>ttyS2</code>	available
<code>ttyS3</code>	available (e.g. for accessing the GSM modem)

Refer to document [3] for further information about hardware setup of serial ports before using them.

If you want to test the `ttyS1` serial port you can, for example, connect with a serial cable your host pc and the GE863-PRO³ on `ttyS1`.

On GE863-PRO³ shell type:

```
# cat /dev/ttyS1 &
```

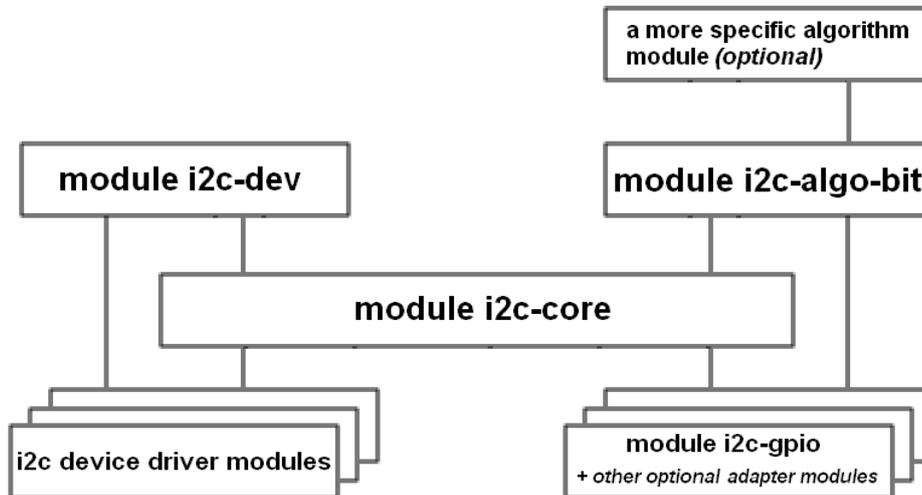
in order to print on standard output any character received over this serial port.

Now open a terminal on you host pc (by default configured for 9600 data rate) and type, for example,

```
hello<carriage return>
```



Although there are many I²C devices and chips, we can use a unique I²C kernel interface composed by a core module used by both master and slave device modules. In the following picture we can see the whole module structure:



So, in our system, in order to use the I²C Interface, the following steps shall be followed:

1. Load the i2c-core module;
2. Load the i2c-algo-bit module;
3. Load the i2c-dev module;
4. Load the i2c-gpio module and/or other adapter driver module(s);
5. Load your I²C device driver module(s).

However, you can find the list of all the available I²C loadable modules in the directory:
 /lib/modules/your_kernel_version/kernel/drivers/i2c

4.2.1 Loading i2c modules

1. The module i2c-core is used by every other I²C modules, so it must be loaded first.

To load this module type the command

```
# modprobe i2c-core
```

2. The i2c-algo-bit module implements a generic algorithm for the communications on the bus. If you need a more specific algorithm, you can modify the i2c-algo-bit source file or implement another module that uses i2c-algo-bit.

To load the i2c-algo-bit module type the command

```
# modprobe i2c-algo-bit
```



- The `i2c-dev` module implements a useful tool that allows the user to access all devices on an adapter from userspace, through the `/dev` interface. So we'll have one item in the `/dev` directory for each bus adapter in our system. Each registered `i2c` adapter gets a number, counting from 0. I²C device files are character device files with major device number 89 and a minor device number corresponding to the number assigned as explained above. So, for example, we'll have in `/dev` the following files: `i2c-0`, `i2c-1`, ..., `i2c-10`, ... All 256 minor device numbers are reserved for `i2c`.

To load the `i2c-dev` module type the command

```
# modprobe i2c-dev
```

- The `i2c-gpio` module is an adapter driver or, in other words, a bus controller for the GPIO bus. If we don't load adapter driver we won't be able to find any item in the `/dev` directory. So we have to load one adapter driver for each I²C bus in our system. In this case, we'll have only one bus and its correspondent driver.

To load the `i2c-gpio` module type the command

```
# modprobe i2c-gpio
```

- Of course you'll need to use the I²C Interface to connect I²C chips or devices to your system. So, after loading the previous modules, you may also need to load the driver module(s) of your specific device(s). In order to do this, you can use the `modprobe` command followed by your module's name.

4.2.2 open()

The `open()` function establishes the connection between a file and a file descriptor. The file descriptor is used by other I/O functions to refer to opened file.

Header file:

`fcntl.h`

Prototype:

```
int open(const char *pathname, int flags)
```

Parameters:

`pathname` – file name with its own path

`flags` – an *int* specifying file opening mode: that can be `O_RDONLY`, `O_WRONLY` or `O_RDWR` which requests opening the file read-only, write-only or read/write, respectively

Returns:

The new file descriptor *fdes* if operation is completed successfully, otherwise it is -1

Example:

Open the `/dev/i2c-0`.

```
int fd; // file descriptor for the /dev/watchdog entry
```




```
I2C_FUNC_SMBUS_READ_BLOCK_DATA
I2C_FUNC_SMBUS_WRITE_BLOCK_DATA
I2C_FUNC_SMBUS_READ_I2C_BLOCK
I2C_FUNC_SMBUS_WRITE_I2C_BLOCK
I2C_FUNC_SMBUS_BYTE
I2C_FUNC_SMBUS_BYTE_DATA
I2C_FUNC_SMBUS_WORD_DATA
I2C_FUNC_SMBUS_BLOCK_DATA
I2C_FUNC_SMBUS_I2C_BLOCK
```

Of course, if the `funcs` mask is different from `FFF800F`, the system will support different functionalities. See the `i2c.h` file in the linux source tree for the available functionalities.

- **I2C_RDWR**

Do a combined read/write transaction without break in between. This is valid only if the adapter has `I2C_FUNC_I2C`.

The difference between doing normal read/write calls and using an `ioctl` with the `I2C_RDWR` request is that in the second case you can do the several read or write using the repeated start condition between two messages. If you use multiple calls to read and write functions instead of the `ioctl` with the `I2C_RDWR` request, you'll send messages with a stop condition between them. The argument of the `ioctl` call is a pointer to a

```
struct i2c_rdwr_ioctl_data {
    struct i2c_msg *msgs; /* ptr to array of simple messages */
    int nmsgs;           /* number of messages to exchange */
}
```

where `nmsgs` is the number of messages to exchange. These messages are contained in the `struct i2c_msg` pointer that points to an array of structures with following definition:

```
struct i2c_msg {
    __u16 addr;           /* slave address */
    __u16 flags;
    __u16 len;           /* msg length */
    __u8 *buf;          /* pointer to msg data */
};
```

In the `addr` field you have to put the address of the slave device you want to communicate with. In the `flags` field you have to set the options among the ones listed below:

- `I2C_M_TEN`: setting this option you indicate that the address in the `addr` field is composed by ten bits, so you'll use the ten bit address mode;
- `I2C_M_RD`: if set, a read operation will be performed; if you don't set it, it'll automatically do a write operation;
- `I2C_M_NOSTART`: if set, there won't be any start condition during the sending of this message with the `i2c` protocol;
- `I2C_M_REV_DIR_ADDR`: you should set this flag if you want to do a write, but need to simulate the process of a Read instead of a normal Write, or vice versa;




```

__u32 size;
union i2c_smbus_data {
    __u8 byte;
    __u16 word;
    __u8 block[I2C_SMBUS_BLOCK_MAX + 2];
                // block[0] is used for length
                // and one more for user-space
compatibility
}
};

```

The `read_write` variable can be configured with `I2C_SMBUS_READ` or with `I2C_SMBUS_WRITE`. The command values are always valid.

The `size` value can be `I2C_SMBUS_BLOCK_PROC_CALL`, `I2C_SMBUS_QUICK`, `I2C_SMBUS_PROC_CALL`, `I2C_SMBUS_BYTE`, `I2C_SMBUS_BYTE_DATA`, `I2C_SMBUS_WORD_DATA`, `I2C_SMBUS_BLOCK_DATA`, `I2C_SMBUS_I2C_BLOCK`

Example:

```

struct i2c_smbus_ioctl_data mycmd;
mycmd.read_write=0;
mycmd.command=0;
mycmd.size=I2C_SMBUS_QUICK; // this means "probe if this address
replies"
mycmd.data=NULL;
ret=ioctl(f, I2C_SMBUS, &mycmd);
if (ret < 0)
    printf("An error occurred in ioctl");

```

- **I2C_RETRIES**

Sets the number of times a device address should be polled when not acknowledging.
Example:

```

unsigned long num_retries = 4;
res = ioctl(file_descriptor, I2C_RETRIES, num_retries);

```

- **I2C_TIMEOUT**

With this macro you can set the timeout value in jiffies.
Example:

```

unsigned long jiffies = 100;
res = ioctl(file_descriptor, I2C_TIMEOUT, jiffies);

```

4.2.4 read()

The `read()` function reads *nbyte* bytes from the file associated with the open file descriptor, *fildev*, and copies them in the buffer that is pointed to by *buf*.



Header file:

unistd.h

Prototype:

ssize_t read(int fildes, void *buf, size_t nbyte);

Parameters:

fildes – file descriptor
 buf – destination buffer pointer
 nbyte – number of bytes that read() attempts to read

Returns:

The number of bytes actually read if if operation is completed successfully, otherwise it is -1.

Example:

Read *sizeof(read_buff)* bytes from the file associated with *fd* and stores them into *read_buff*.

```
char read_buff[BUFF_LEN];

if(read(fd, read_buff, sizeof(read_buff)) < 0)
{
    /* Error Management Routine */
} else {
    /* Value Read */
}
```

4.2.5 write()

The *write()* function writes *nbyte* bytes from the buffer that are pointed by *buf* to the file associated with the open file descriptor, *fildes*.

Header file:

unistd.h

Prototype:

ssize_t write(int fildes, const void *buf, size_t nbyte);

Parameters:

fildes – file descriptor
 buf – destination buffer pointer
 nbyte – number of bytes that write() attempts to write

Returns:

The number of bytes actually written if operation is completed successfully (this number shall never be greater than *nbyte*), otherwise it is -1.

Example:



Write `sizeof(value_to_be_written)` bytes from the buffer pointed by `value_to_be_written` to the file associated with the open file descriptor, `fd`.

```
char value_to_be_written[] = "dummy_write";

if (write(fd, value_to_be_written, sizeof(value_to_be_written)) < 0)
{
    /* Error Management Routine */
} else {
    /* Value Written */
}
```

4.2.6 close()

The `close()` function deallocates the file descriptor indicated by `filides`. To deallocate means to make the file descriptor available for subsequent calls to `open()` or other functions that allocate file descriptors.

Header file:

unistd.h

Prototype:

int close(int filides);

Parameters:

filides – file descriptor

Returns:

0 if operation is completed successfully, otherwise it is -1

Example:

Close the I²C device.

```
if(close(fd) < 0)
{
    /* Error Management Routine */
} else {
    /* File Closed */
}
```

4.2.7 A Test Program

The following simple C program is useful to test the I²C interface. It opens the `/dev/i2c` interface and calls the write function in an infinite loop to write some random values on the device.

```
#include <sys/types.h>
#include <sys/stat.h>
#include <fcntl.h>
```



```

#include <sys/ioctl.h>
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <linux/i2c.h>
#include <linux/i2c-dev.h>

int main()
{
    int fd;
    char val;
    char buf[5];
    int addr = 0x40; /* This I2C address is a random value
                     because no device is attached
                     on the bus; if you connect a device,
                     change this value with the correct one*/
    fd = open("/dev/i2c-0",O_RDWR);
    if (fd < 0) {
        fprintf(stderr, "Error during open\n");
        exit(1);
    }

    /* The following ioctl call sets in slave mode
       the device located at the addr address */
    if (ioctl(fd, I2C_SLAVE, addr) < 0) {
        fprintf(stderr, "Error during IOCTL\n");
        exit(1);
    }
    buf[0] = 'T';
    buf[1] = 'e';
    buf[2] = 's';
    buf[3] = 't';

    for(val=1;;val=(val+1)%255) {
        buf[4] = val;
        if (write(fd, buf, 5) != 5) {
            fprintf(stderr, "Write error (%d)\n",val);
        } else {
            fprintf(stderr, "Write OK (%d)\n",val);
        }
        usleep(500000);
    }
    return 0;
}

```



}

If you have some I²C device connected to the board, you'll have to set the *addr* variable with the correct address of the I²C device you want to write. Otherwise, you can initialize that variable with a random value. In this case, you can connect an oscilloscope to the SDA and SCL pins of the board to observe the behaviour of the signal on the pins.

If you connect an I²C device to the board the write call should be successful and the program should report "Write OK". Otherwise the program displays "Write error", but this is not a problem: it only means that there is no device connected on the bus.

4.3 SPI

The "Serial Peripheral Interface" (SPI) is a synchronous four wire serial link used to connect microcontrollers to sensors, memory, and peripherals.

The Serial Peripheral Interface is essentially a shift register that serially transmits data bits to other SPIs. During a data transfer, one SPI system acts as the "master" and controls the data flow, while the other devices act as "slaves".

The SPI system consists of two data lines and two control lines:

- Master Out Slave In (MOSI): This data line supplies the output data from the master shifted into the input(s) of the slave(s).
- Master In Slave Out (MISO): This data line supplies the output data from a slave to the input of the master. There may be no more than one slave transmitting data during any particular transfer.
- Serial Clock (SPCK): This control line is driven by the master and regulates the flow of the data bits. The master may transmit data at a variety of baud rates; the SPCK line cycles once for each bit that is transmitted.
- Slave Select (NSS): This control line allows slaves to be turned on and off by hardware (it is also called chipselect).

4.3.1 Loading the SPI module

The SPI driver is released to customer under in the form of a loadable module.

To load the SPI module type in the terminal:

```
# modprobe spidev
```

When loaded, the SPI driver will install a new device named `spidev1.1` under the `/dev/` directory. The "1.1" name extension represents the number of the selected SPI bus and the selected chipselect.

Once installed the SPI device can be used as a normal character device and can be accessed by any application running in userspace.



4.3.2 open()

The *open()* function shall establish the connection between a file and a file descriptor. The file descriptor is used by other I/O functions to refer to the opened file.

Header file:

fcntl.h

Prototype:

int open(const char *pathname, int flags)

Parameters:

pathname – file name with its own path

flags – is an *int* specifying file opening mode: is one of O_RDONLY, O_WRONLY or O_RDWR which request opening the file read-only, write-only or read/write, respectively

Returns:

The new file descriptor *filides* if operation is completed successfully, otherwise it is -1.

Example:

Open the /dev/spidev1.1.

```
int fd;           // file descriptor for the /dev/watchdog entry

if((fd = open("/dev/spidev1.1", O_RDWR) < 0)
{
    /* Error Management Routine */
} else {
    /* SPI Device Opened */
}
```

4.3.3 ioctl()

The *ioctl()* function manipulates the underlying device parameters. In particular, many operating characteristics can be controlled with *ioctl()* requests.

Header file:

sys/ioctl.h
linux/spi/spidev.h

Prototype:

int ioctl(int fildes, int request, ...)

Parameters:

fildes – file descriptor

request – device-dependent request code.

The following *ioctl*s request codes can be used for SPI:

- SPI_IOC_RD_MODE and SPI_IOC_WR_MODE to get/set the SPI mode



▪ SPI_IOC_MESSAGE

Standard read() and write() operations are obviously only half-duplex, and the chipselect is deactivated between those operations. Full-duplex access, and composite operation without chipselect de-activation, is available using the SPI_IOC_MESSAGE(N) request.

Please note that the SPI_IOC_MESSAGE(N) request needs, as parameter, a pointer to struct spi_ioc_transfer whose fields speed_hz and bits_per_word must be set to 0.

Example:

```
uint8_t tx[] = {
    0xFF, 0xFF, 0xFF, 0xFF, 0xFF, 0xFF,
    0x40, 0x00, 0x00, 0x00, 0x00, 0x95,
    0xFF, 0xFF, 0xFF, 0xFF, 0xFF, 0xFF,
    0xFF, 0xFF, 0xFF, 0xFF, 0xFF, 0xFF,
    0xFF, 0xFF, 0xFF, 0xFF, 0xFF, 0xFF,
    0xDE, 0xAD, 0xBE, 0xEF, 0xBA, 0xAD,
    0xF0, 0x0D,
};
uint8_t rx[ARRAY_SIZE(tx)] = {0, };
struct spi_ioc_transfer tr = {
    .tx_buf = (unsigned long)tx,
    .rx_buf = (unsigned long)rx,
    .len = ARRAY_SIZE(tx),
    .delay_usecs = delay, /*Delay before chipselect
deactivation*/
    .speed_hz = 0, /* Must be set to 0 */
    .bits_per_word = 0, /* Must be set to 0*/
};
for (i = 0; i < 100 ; i++)
{
    ret = ioctl(fd, SPI_IOC_MESSAGE(1), &tr);
    if (ret == -1)
        printf("can't send spi message");
}
```

4.3.4 read()

The read() function reads nbyte bytes from the file associated with the open file descriptor, fildes, and copies them in the buffer that is pointed to by buf.

Header file:

unistd.h

Prototype:

ssize_t read(int fildes, void *buf, size_t nbyte);

Parameters:



filides – file descriptor
buf – destination buffer pointer
nbyte – number of bytes that `read()` attempts to read

Returns:

The number of bytes actually read if operation is completed successfully, otherwise it is -1.

Example:

Read `sizeof(read_buff)` bytes from the file associated with `fd` and stores them in `read_buff`.

```
char read_buff[BUFF_LEN];

if(read(fd, read_buff, sizeof(read_buff)) < 0)
{
    /* Error Management Routine */
} else {
    /* Value Read */
}
```

4.3.5 write()

The `write()` function writes *nbyte* bytes from the buffer that are pointed by *buf* to the file associated with the open file descriptor, *filides*.

Header file:

unistd.h

Prototype:

ssize_t write(int *filides*, const void **buf*, size_t *nbyte*);

Parameters:

filides – file descriptor
buf – destination buffer pointer
nbyte – number of bytes that `write()` attempts to write

Returns:

The number of bytes actually written if operation is completed successfully, (this number shall never be greater than *nbyte*), otherwise it is -1

Example:

Write `sizeof(value_to_be_written)` bytes from the buffer pointed by `value_to_be_written` to the file associated with the open file descriptor, `fd`.

```
char value_to_be_written[] = "dummy_write";

if (write(fd, value_to_be_written, sizeof(value_to_be_written)) < 0)
{
    /* Error Management Routine */
} else {
```



```

        /* Value Written */
    }

```

4.3.6 close()

The *close()* function deallocates the file descriptor indicated by *filides*. To deallocate means to make the file descriptor available for subsequent calls to *open()* or other functions that allocate file descriptors.

Header file:

unistd.h

Prototype:

int close(int filides);

Parameters:

filides – file descriptor

Returns:

0 if operation is completed successfully, otherwise it is -1

Example:

Close the SPI device.

```

if(close(fd) < 0)
{
    /* Error Management Routine */
} else {
    /* File Closed */
}

```

4.3.7 A Test Program

Below it is reported a simple piece of code that opens the device, writes some random data, reads some data if available and then closes the devices. Please note that the following code works for half duplex communication. For typical full duplex communication the SPI_IOC_MESSAGE ioctl will be used.

In order to receive data when executing the read() function a transmitting SPI peripheral must be connected.

```

#include <stdint.h>
#include <unistd.h>
#include <stdio.h>
#include <stdlib.h>
#include <getopt.h>
#include <fcntl.h>
#include <sys/ioctl.h>
#include <linux/types.h>
#include <linux/spi/spidev.h>

```




```

        return;
    }

    int main(int argc, char *argv[])
    {
        int ret = 0;
        int fd;

        fd = open(device, O_RDWR);
        if (fd < 0)
        {
            printf("can't open device");
            return -1;
        }
        myTransfer(fd);

        close(fd);

        return ret;
    }

```

4.4 GPIO

The GPIO driver creates a series of devices under the `/dev/` directory. Each device is named as `at91sam9260_gpio.<pin>`.

GPIO pins are numbered from 0 to 95. The first 32 pins refer to the A bank, the second 32 pins refer to the B bank and the others refer to the C bank.

GPIO pins can be read and written thorough simple read/write operations.

The read operation returns the status and the level of the pin if it is configured as a GPIO, otherwise it fails. The write operation can change the configuration and output value. The configuration can be changed using:

- O for output enable with pull up
- o for output enable without pull up
- I input enabled with pull up
- i input enabled without pull up
- 1 set the level up
- 0 set the level low

To test a gpio you can use the shell, for example (gpio number 92):

```
# echo 01 > /dev/at91sam9260_gpio.92
```

to raise up the signal.



The following subparagraphs show all the functions that can be used from C source code to perform read/write operations onto GPIOs pins.

4.4.1 open()

The *open()* function establishes the connection between a file and a file descriptor. The file descriptor is used by other I/O functions to refer to the opened file.

Header file:

fcntl.h

Prototype:

```
int open(const char *pathname, int flags)
```

Parameters:

pathname – file name with its own path

flags – is an *int* specifying file opening mode: is one of O_RDONLY, O_WRONLY or O_RDWR which request opening the file read-only, write-only or read/write, respectively

Returns:

The new file descriptor *filides* if operation is completed successfully, otherwise it is -1.

Example:

Open the GPIO device (for example number 92).

```
int fd;          // file descriptor for GPIO /dev entries

if((fd = open("/dev/at91sam9260_gpio.92", O_RDWR)) < 0)
{
    /* Error Management Routine */
} else {
    /* Gpio Opened */
}
```

4.4.2 read()

The *read()* function reads *nbyte* bytes from the file associated with the open file descriptor, *filides*, and copies them in the buffer that is pointed to by *buf*.

Header file:

unistd.h

Prototype:

```
ssize_t read(int fildes, void *buf, size_t nbyte);
```

Parameters:

fildes – file descriptor

buf – destination buffer pointer

nbyte – number of bytes that read() attempts to read



Returns:

The number of bytes actually read if operation is completed successfully, otherwise it is -1

Example:

Read *sizeof(read_buff)* bytes from the file associated with *fd* and stores them into *read_buff*.

```
char read_buff[BUFF_LEN];

if(read(fd, read_buff, sizeof(read_buff)) < 0)
{
    /* Error Management Routine */
} else {
    /* Value Read */
}
```

4.4.3 write()

The *write()* function writes *nbyte* bytes from the buffer that are pointed by *buf* to the file associated with the open file descriptor, *filides*.

Header file:

unistd.h

Prototype:

ssize_t write(int filides, const void *buf, size_t nbyte);

Parameters:

- filides – file descriptor
- buf – destination buffer pointer
- nbyte – number of bytes that write() attempts to write

Returns:

The number of bytes actually written if operation is completed successfully (this number shall never be greater than *nbyte*), otherwise it is -1.

Example:

Write *sizeof(value_to_be_written)* bytes from the buffer pointed by *value_to_be_written* to the file associated with the open file descriptor, *fd*.

```
char value_to_be_written[] = "01";

if (write(fd, value_to_be_written, sizeof(value_to_be_written)) < 0)
{
    /* Error Management Routine */
} else {
    /* Value Written */
}
```



4.4.4 close()

The *close()* function deallocates the file descriptor indicated by *filides*. To deallocate means to make the file descriptor available for subsequent calls to *open()* or other functions that allocate file descriptors.

Header file:
unistd.h

Prototype:
int close(int filides);

Parameters:
filides – file descriptor

Returns:
0 if operation is completed successfully, otherwise it is -1.

Example:
Close the GPIO device.

```
if(close(fd) < 0)
{
    /* Error Management Routine */
} else {
    /* File Closed */
}
```

4.5 Watchdog

A Watchdog Timer (WDT) is a hardware circuit that can reset the computer system in case of a software fault.

Usually a userspace daemon will notify the kernel watchdog driver via the */dev/watchdog* special device file that the userspace is still alive, at regular intervals. When such a notification occurs, the driver will usually tell the hardware watchdog that everything is in order, and that the watchdog should wait before resetting the system. If userspace fails (RAM error, kernel bug, whatever), the notifications cease to occur, and the hardware watchdog will reset the system (causing a reboot) after the timeout expires.

Userspace can interact with the kernel watchdog driver through the functions shown in the paragraphs below.

4.5.1 open()

The *open()* function establishes the connection between a file and a file descriptor. The file descriptor is used by other I/O functions to refer to the opened file.

Header file:



fcntl.h

Prototype:

int open(const char *pathname, int flags)

Parameters:

pathname – file name with its own path

flags – an *int* specifying file opening mode: that can be O_RDONLY, O_WRONLY or O_RDWR which request opening the file read-only, write-only or read/write, respectively

Returns:

The new file descriptor *filides* if operation is completed successfully, otherwise it is -1

Example:

Open the /dev/watchdog.

```
int fd;           // file descriptor for the /dev/watchdog entry

if((fd = open("/dev/watchdog", O_WRONLY) < 0)
{
    /* Error Management Routine */
} else {
    /* Watchdog Device Opened */
}
```

4.5.2 ioctl()

The ioctl() function manipulates the underlying device parameters. In particular, many operating characteristics can be controlled with ioctl() requests.

Header file:

sys/ioctl.h
linux/watchdog.h

Prototype:

int ioctl(int fildes, int request, ...)

Parameters:

fildes – file descriptor

request – device-dependent request code.

The following ioctls request codes can be used for watchdog device:

- WDIOCG_KEEPALIVE: to notify the watchdog that userspace is still alive
- WDIOCG_SETTIMEOUT: to set a timeout
- WDIOCG_SETOPTIONS: to enable or disable the watchdog
- WDIOCG_GETTIMEOUT: to query the timeout

The third argument is a *void ** and depends on the ioctl request code used: see the examples below.

Returns:



0 if operation is completed successfully, otherwise it is -1

Examples:

- WDIOC_KEEPALIVE

The AT91Sam9260 watchdog driver supports the WDIOC_KEEPALIVE ioctl. It notifies the watchdog that userspace is still alive. In this case the third argument in the ioctl is ignored. This function call resets the timer and leaves the system alive (is used to avoid reset when the watchdog is active). Example:

```
fd = open("/dev/watchdog", O_WRONLY);
int dummy;
for(;;){
    ioctl(fd, WDIOC_KEEPALIVE, &dummy);
    sleep(1)
}
```

- WDIOC_SETTIMEOUT

Setting Timeout is performed by the SETTIMEOUT ioctl. The third argument is an integer that represents the timeout in seconds (max timeout is 16 seconds in at91sam9260 architecture). The driver returns the real timeout used in the same variable, and this timeout can be different from the one that has been set due to limitation of the hardware.

The watchdog timeout can be written only once (at91sam9260 only permits one program operation).

Example:

```
int timeout=15;
ioctl(fd, WDIOC_SETTIMEOUT, &timeout);
```

Notice: If the watchdog start is enabled the user must set the timeout otherwise it will use the default value.

- WDIOC_SETOPTIONS

It enables or disables the watchdog. The third argument is an integer indicating the option to be set.

Example to disable the watchdog:

```
int options = WDIOS_DISABLECARD;
ioctl(fd, WDIOC_SETOPTIONS, &options);
```

Example to enable the watchdog:

```
int options = WDIOS_ENABLECARD;
ioctl(fd, WDIOC_SETOPTIONS, &options);
```

- WDIOC_GETTIMEOUT



It is possible to query the timeout using the `WDIOC_GETTIMEOUT` ioctl. The third argument is an integer representing the timeout in seconds.

Example:

```
ioctl(fd, WDIOC_GETTIMEOUT, &timeout);
printf("The timeout is %d seconds\n", timeout);
```

4.5.3 close()

The `close()` function deallocates the file descriptor indicated by *filides*. To deallocate means to make the file descriptor available for subsequent calls to `open()` or other functions that allocate file descriptors.

Header file:

unistd.h

Prototype:

```
int close(int filides);
```

Parameters:

filides – file descriptor

Returns:

0 if operation is completed successfully, otherwise it is -1

Example:

Close the watchdog device.

```
if(close(fd) < 0)
{
    /* Error Management Routine */
} else {
    /* File Closed */
}
```

4.6 SDIO

GE863-PRO³ support for Secure Digital (SD) and Multimedia Card (MMC) is built in the kernel.

First create a directory where the device will be mounted, for example:

```
# mkdir /mnt/sdcard
```

Then, connect the device and mount it:

```
# mount /dev/mmcblk0p1 /mnt/sdcard
```



