1. PROCEDURAL PROGRAMMING

The Finite State Machine portions of IsoMaxTM are now covered. What remains to be discussed is the procedural portions of the conditions and actions.

The IsoMax procedural language is very similar to the programming language Forth. There are some significant changes because of the Finite State Machine functions, the Object-Oriented functions, and the architecture of the IsoPod[™] processor (which has separate data and program memories, and can't address memory as bytes). But if you are familiar with Forth, and particularly Max-Forth, you can skip most of this section.

1.1. The Dictionary

All of the commands and operations known to IsoMax are kept in a "dictionary." You will frequently hear these commands and operations called "words," because they are the words in the IsoMax language. In fact, you can print out a list of all the words known to IsoMax with the command

WORDS

This just prints a list of the words. Their definitions are much longer! You can find short definitions for all these words in section 18 of this manual (IsoMax Glossary).

Possible point of confusion: Don't confuse a "word" in the language, with a "word" of memory. A memory word (on the IsoPodTM) is 16 bits of storage. A word of the language can be any symbol made of non-blank characters. When there is risk of confusion, we will generally refer to 16 bits of storage as a "cell" of memory.

1.2. The Stacks

Numbers, addresses, and data which is being operated upon are normally held on *stacks*. Like the name implies, when you put something on a stack, it becomes the topmost item on the stack, and everything that was already on the stack is effectively "pushed down" one deeper. Likewise, when you take something off the top of the stack, the stuff underneath it "pops up" one position, and what was the second item on the stack becomes the new top item on the stack.

For the most part, the operation of stacks is invisible and automatic. One visible effect (which will be familiar to owners of Hewlett-Packard calculators) is that arithmetic operations require you to place the two operands on the stack first, and then specify the operation to be performed. In other words, instead of saying

2 + 3

in IsoMax you would say

2 3 +

We'll see more examples of this shortly.

There is one stack for fixed-point data (including integers, characters, and addresses), and a second stack for floating-point numbers. You don't need to specify this – each operation automatically uses the correct stack. But you might need to be aware of the two different stacks, if you're moving values to and from them.

1.2.1. Stack Notation

Whenever you type a number, IsoMax puts it on the stack. (Integers on the integer stack, floating-point numbers on the floating-point stack.) If you type a second number, it gets put on the stack and the previous number gets pushed down. So, for instance, if you type

1 2 3

the "3" will be the topmost item on the stack, "2" will be under it, and "1" will be on the bottom. You can see this with the command .S ("print stack"). Try typing the command

123.S

and you will see how the numbers are "stacked" (with the "3" on the top).

When we are describing parameters to be put on the stack, or values which are placed on the stack, we will use this left-to-right notation. The rightmost item in the description corresponds to the topmost item on the stack. So when you see parameters

a b c

you know it really means

c b a

with "c" on the top of the stack.

1.3. Arithmetic Operations

Most of the time you'll be using 16-bit integers. These can be treated as unsigned numbers, in the range 0..65535, or as signed numbers, in the range -32768..+32767. The four basic arithmetic operators are add, subtract, multiply, and divide:

10	2	+	adds 10 and 2, giving 12
10	2	-	subtracts 10-2, giving 8
10	2	*	multiplies 10 by 2, giving 20 (signed numbers)
10	2	/	divides 10 by 2, giving 5 (signed numbers)

Notice the order of the operands for subtract and divide. This is easy to remember, because it's the same left-to-right order you would use if you were writing these as algebraic equations. That is,

10	2	-	performs the computation	10	_	2
10	2	/	performs the computation	10	/	2

If you want to see the result of these computations, you can use the . S command. Or, you can use the . command (just a period character), which prints the topmost stack item and removes it from the stack. For example, try

10 3 / .

and you should see the result 3. Why 3 and not 3.33333? We're using integer math, so we get "3 with a remainder of 1," not 3.33333. To see the remainder. use the MOD command:

10 3 MOD .

When the IsoPodTM starts running, it expects decimal (base 10) numbers. But you can change this at any time. If you type the command HEX, all numbers from that point on will be entered and printed in hexadecimal (base 16). To change back to base 10, type the command DECIMAL. You can use this with the . (print) command to perform simple base conversions. For example:

HEX AO DECIMAL .

will print 160, the decimal equivalent of A0 hex.

You can type negative numbers, like -12 or even -FFF (in hexadecimal base). If you want to negate the result of a computation, you can use the command NEGATE. For example,

-12 3 * NEGATE .

will print 36, because -12 times 3 is -36, and the negative of -36 is 36.

1.4. Floating-Point Operations

To type a floating-point number into IsoMax, you *must* include an exponent, in the form Enn, as the suffix of the number. For example, all of these represent a floating-point value of 2.0:

2.0E0 2.000E0 2E0 0.2E1 .2E01 20.E-1

The "E" followed by a (positive or negative) number is required. The following will *not* work:

2.0 2. 002 2.000

Also, you *must* be in DECIMAL base to type a floating point number. This is because "E" is a valid hexadecimal digit. So,

DECIMAL	2E0	gives the floating-point value 2.0, but
HEX 2E0		gives the integer 2E0 hex (736 decimal)

You can use the command F.S to display the contents of the floating-point stack. The command F. will print (and remove) the topmost item on the floating-point stack in a "fixed point" notation, and the command E. will print (and remove) the topmost item in an "exponential" notation. To see this, try

```
DECIMAL 1.E1 2.E1 3.E1
F.S
F.
E.
E.
```

Floating-point arithmetic operations are similar to the integer operations, but have the prefix "F". Here are add, subtract, multiply, and divide:

10.E0	3.E0	F+	adds 10.0 and 3.0, giving 13.0
10.E0	3.E0	F-	subtracts 10.0-3.0, giving 7.0
10.E0	3.E0	F*	multiplies 10.0 by 3.0, giving 30.0
10.E0	3.E0	F/	divides 10.0 by 3.0, giving 3.3333

Other floating-point operations include trigonometric and transcendental functions. The complete list can be found in the glossary, Section 18.4.

1.5. Variables

Because IsoMax carries out computations on its stacks, you very rarely need to use "variables" such as X or Y or VELOCITY or VOLTAGE. But sometimes you do need to store a value between computations. So, IsoMax allows you to have named variables.

You must define a variable before you use it. This is done with the command VARIABLE (integer) or FVARIABLE (floating-point):

Integer		Floating-point			
VARIABLE	name	FVARIABLE	name		

In either case, "name" is a name you choose for the variable. This can be any combination of up to 31 non-blank characters. Even special characters and punctuation may be freely used. For example, the following are all valid variable names:

X Y1 Velocity \$PROFIT \$ 4TH_SPEED %#@!

Names can begin with numbers, and can be entirely non-alphabetic characters. Two restrictions, though. First, don't use a name that's already in use by IsoMax (as you can see with WORDS). This will cause confusion. IsoMax will allow it, but will warn you by telling you that your name is "not unique."

Second, don't use a name that's all numbers. IsoMax will allow that, and *won't* warn you, and then when you type that number, you'll get the variable instead of the number. As you can imagine, this will lead to no end of confusion. Be sure that all your names have one non-numeric character (and remember that A through F are digits in hex)!

You can use upper or lower case in your names, but remember that IsoMax is casesensitive. VELOCITY, VELocity, and Velocity are all different names.

When you have defined a variable, you can store a value into that variable with the ! or F! commands. You can fetch the stored value with the @ and F@ commands.

Integer Floating-point

name !name F!stores a value in variable "name"name @name F@fetches a value from variable "name"

This is not like other languages, which let you use just the name of a variable in place of a number in an equation. To get the value of a variable, you *must* use @ (for integer variables) or F@ (for floating-point variables). So, if you want to multiply (floating-point) Principal by Interest to compute a payment, you'd have to type

Principal F@ Interest F@ F* F.

Of course, before you could do this you would have had to define the variables

FVARIABLE Principal FVARIABLE Interest

and you would have had to store some values into these variables, e.g.,

```
10000.E0 Principal F!
0.05E0 Interest F!
```

1.6. Accessing Memory and I/O

When you define a VARIABLE, what you're really doing is reserving a memory location and giving that memory location a name. The operators @ and ! fetch from a memory location, and store to a memory location, respectively.

You can use @ and ! with *any* memory locations, not just variables. Suppose that you know a value is stored at memory address \$6A2. You can get that value with

HEX 6A2 @

Suppose you want to store a value of \$1234 into that location. You can use

HEX 1234 6A2 !

When might this be useful? Most of the time, you'll want to use named VARIABLES, because a variable will always be placed in an unused part of memory. If you try to choose memory addresses yourself, you might choose an address that's being used by IsoMax for something else.

But there is one situation when you might want to read or write a known memory location. The input and output of the IsoPod's DSP56F805 CPU is *memory-mapped*. This means that, instead of accessing the peripherals with IN and OUT instructions, you use normal memory fetch and store instructions. So, @ and ! give you access to the complete I/O capability of the IsoPod processor!

For example, the Port A Data Register is located at address \$0FB1. The Port A Data Direction Register (DDR) is located at address \$0FB2, and the Peripheral Enable Register (PER) is at \$0FB3. Writing zero to both the DDR and PER will make Port A an input port, and you can then read the Data Register to read the eight input pins.

```
HEX
0 0FB2 !
0 0FB3 !
0FB1 @
```

This is not for the inexperienced user. There are *no* restrictions on @ and !, so it's quite possible for you to lock up the IsoPod completely by writing the wrong value to the wrong location. You should refer to Motorola's *DSP56F801/803/805/807 16-Bit Digital Signal Processor User's Manual* for a complete description of the on-chip I/O of the DSP56F805 processor, and its memory addresses.

Besides, isn't this easier?

PORTA GETBYTE

ServoPod owners note: The ServoPod uses the DSP56F807 processor, which has *different* I/O addresses from the IsoPod's DSP56F805.

1.7. Logical Operations

IsoMax also lets you perform logical operations on 16-bit values. The four basic arithmetic operators are AND, OR, XOR, and INVERT:

HEX	3A	0F	AND	bitwise logical AND, giving 0A
HEX	3A	0F	OR	bitwise logical OR, giving 3F
HEX	3A	0F	XOR	bitwise exclusive OR, giving 35
HEX	3A	INV	/ERT	bitwise inversion, giving FFC5

Notice that INVERT takes only one parameter. Also, all of the logical operations act on 16 bit values. If you try to print the result with the . operator, you may be surprised:

HEX 3A INVERT . prints -3B

This is because FFC5 is a negative number in two's complement notation, and . prints signed numbers. To print unsigned numbers, use the U. ("unsigned print") operator:

HEX 3A INVERT U. prints FFC5

The logical operations are is especially useful when you're working with I/O, when you need to act on specific bits. For example, suppose you need to read the low 4 bits of port A as a hex number from 0 to F. You could read the four bits individually, and write some code to merge them into a 4-bit value. But it's much easier to say

PORTA GETBYTE HEX OF AND

which reads all 8 bits of the port, and then "masks off" the unwanted bits.

What if you need the *high* 4 bits of the port? It's only an 8-bit port, so the "mask" should be F0 instead of 0F. You then need to shift the bits "down" four places:

PORTA GETBYTE HEX F0 AND 2/ 2/ 2/ 2/

The 2/ operator ("two-divide") gives a one-bit right shift. For integers, this is equivalent to dividing by two, hence the name. Applying it four times gives a total of four shifts to the right.

The corresponding left-shift operator is 2* ("two-times").

1.8. Adding New Definitions

Much of the time you will be using the IsoMax operations interactively -- as we have seen above -- or in the CONDITION or CAUSES phrases of an IsoMax state machine. What you have learned so far is sufficient for these uses, although you'll probably want to look at the IsoMax Glossary in Section 18 to see the full range of operations which are available to you.

But there may come a time when you want to create a procedural subroutine. This might be because

- There's a complex function you perform frequently, and you're tired of typing it all the time, or
- You want to write a computer program (in the traditional sense) and commit it to the IsoPod's memory.

In either case, you do this by adding a new word to the IsoMax dictionary. This new word will contain your complex function or your application program.

Just like with an English dictionary, you add a new word by first giving the name of the new word, and then defining that word *using only words which are already known*. IsoMax marks the start and end of a new definition with : and ; as follows:

: name-of-new-word ...definition... ;

The spaces after : and before ; are required. The name of the new word can be any combination of up to 31 non-blank characters, just like VARIABLE names.

Let's go back to our recent example, and assume that a 4-bit DIP switch is connected to the high 4 bits of Port A. We know how to read this port, mask the bits, and shift them to the low 4 bits. But we're going to be doing this a lot, and we don't want to type that long phrase every time. Also, to make the code more readable (and more maintainable), we'd like to call it something meaningful like GET-DIP-SWITCH. Here's how you can do it:

```
HEX
: GET-DIP-SWITCH PORTA GETBYTE F0 AND
2/ 2/ 2/ 2/;
```

Here we are telling IsoMax to add a new word, GET-DIP-SWITCH, to the dictionary. The "definition" of this new word is PORTA GETBYTE HEX F0 AND 2/ 2/ 2/ 2/. What this means in practice is that, whenever IsoMax sees GET-DIP-SWITCH, it will perform the action PORTA GETBYTE HEX F0 AND 2/ 2/ 2/ 2/. (Strictly speaking, we've created a *subroutine* containing those IsoMax instructions.)

Remember that IsoMax is *free-format* so you can split the definition across multiple lines, and use spaces to indent. **TAKE NOTE:** if you're going to use numbers inside the

definition, you must specify the number base *outside* the definition. In this example, we put HEX before we started the definition.¹

GET-DIP-SWITCH will have exactly the same stack effect as its definition. Since PORTA GETBYTE HEX F0 AND 2/ 2/ 2/ 2/. leaves a single value on the stack, GET-DIP-SWITCH will leave a single value on the stack.

Of course, we could also have written something which takes values from the stack. Maybe we need a four-bit right shift frequently:

: RIGHT-SHIFT-4 2/2/2/2/;

This will take a value on the stack, shift it right four times, and then leave the result on the stack. So you see, the stack is how we pass values to a function, and how we get results from a function. These are the *input parameters* and *output parameters* of the function.

Well, once we've told IsoMax what RIGHT-SHIFT-4 means, why can't we use that to define GET-DIP-SWITCH? We can:

```
: RIGHT-SHIFT-4 2/2/2/;
HEX
: GET-DIP-SWITCH PORTA GETBYTE F0 AND RIGHT-SHIFT-4;
```

1.8.1. Removing definitions

If you've been typing this example in, you've probably seen the warning GET-DIP-SWITCH NOT UNIQUE. This is IsoMax telling you that you've defined a word twice in its dictionary. This won't break IsoMax, but it will cause you some confusion, since you won't necessarily know what definition is being used at any given time.

It's better for all concerned if you tell IsoMax to forget your previous definition of the word. You do this with the command

FORGET GET-DIP-SWITCH

This gets rid of the old definition, and leaves you free to start a new one. (Strictly speaking, FORGET gets rid of the *most recent* definition of the word. If you've defined the word twice, you'll need to use FORGET twice to get rid of both definitions.)

TAKE NOTE: FORGET will not just forget the word you specify, it will forget *all words you have defined since that word*. In the last example, if you had typed FORGET RIGHT-SHIFT-4, you would *also* lose the definition of GET-DIP-SWITCH. This can be useful: if you've written a few dozen words, and you want to forget them all (so you

¹ If you put HEX inside the definition, that number base won't take effect until later, when you execute GET-DIP-SWITCH. This is sometimes useful, but usually is not what you want.

can start over), you don't need to type a few dozen FORGET commands. Just forget the first word, and all the following words will go too.

1.9. Program Control

We've seen how to write subroutines (as new word definitions), how to do arithmetic and logical functions, how to store data in memory variables, and how to do I/O. There's still one thing missing before we can write any computer program: how do we perform actions *conditionally*? That is, how do we change the *flow of control* of the program, based upon an input or the result of a calculation?

IsoMax offers six different constructs for program control. These correspond to the basic *control structures* from the discipline of structured programming.

1. IF some action... THEN

This performs an action *if* some condition is true. The condition is given by a value on the stack when IF is encountered. A zero value is "false", and any nonzero value will be considered "true." A true value causes the code between IF and THEN to be performed. A false value causes that code to be skipped.





2. IF ..true action.. ELSE ..false action.. THEN

This is similar to IF...THEN, except that it performs one action if the condition is true (nonzero), and a different action if the condition is false (zero). **Remember** that, unlike some other languages, THEN terminates the control structure. The code following then is always executed.

3. BEGIN some action... UNTIL

This performs an action repeatedly *until* some condition is true. The condition is given by a value on the stack when UNTIL is encountered. A zero (false) value means "do the action again," and will cause a loop from the UNTIL back to the BEGIN. A nonzero (true) value means "terminate the loop," and will cause execution to continue on to the code after the UNTIL. Note that it is the action inside the loop that produces the true/false value for UNTIL! Whatever else is done, this action *must* include code which leaves this "exit/loop" value on the stack. Note also that the action inside the loop will always be performed at least once!



4. BEGIN ... condition... WHILE ... some action... REPEAT



This performs an action repeatedly *while* some condition is true. This is similar to BEGIN. . UNTIL with the following differences:

- The code that produces the "exit/loop" value is placed *before* the WHILE, and the action to be taken is *after* the WHILE.
- A zero (false) value means "terminate the loop"; a nonzero (true) value means "do the action and keep looping."
- It's possible for the action to be performed *zero* times.

Another way to look at this: a false value at the WHILE will cause the program to immediately jump to the code after the REPEAT (thus exiting the loop). A true value at the WHILE will cause the code immediately following the WHILE to be executed, and then

REPEAT will loop back to BEGIN.

end start DO ...some action... LOOP

This performs an action repeatedly for a given number of times. This loops over values from "start" to "end-1". The "end" and "start" values are given on the stack when DO is encountered, with the "start" value on the top of stack, and the "end" value second on the stack. These values can be determined from a computation, but often will just be numeric constants. For example



will perform the action I. ten times, with the *loop index* going from 0 to 9 (inclusive). The operator I will always return the value of the current loop index. (Unlike other languages, you don't need to use a variable for this.) So in this example, the action I. will print the loop index on each pass through the loop.

end start DO ...some action... n +LOOP

10 0 DO I. LOOP

This is similar to DO..LOOP except that the loop index is incremented by "n" instead of 1. "n" is the value on top of the stack when +LOOP is encountered; it is usually a constant, but could be the result of a computation. It may be positive or negative. If the increment is negative, "end" must be less than "start," and the loop will proceed all the way to the end value (not end-1).

IMPORTANT LIMITATION: All of these control structures can be used inside a word definition, and inside a CONDITION or CAUSES phrase in an IsoMax state machine. But they *can not* be used interactively from the command line.

1.9.1. DO Loop Example

It's instructive to write some simple definitions which show how DO loops work:

```
DECIMAL
: TEST#1 10 0 DO I . LOOP ;
TEST#1
```

This will just print the value of the loop index as it goes from 0 to 9.

: TEST#2 DO I . LOOP ; 10 0 TEST#2 30 20 TEST#2 This has the same action as TEST#1, but instead of "hard coding" the loop limits inside the definition, we are passing them as parameters on the stack. So we can try the loop with a number of different start and end parameters.

```
: TEST#3 DO I . 3 +LOOP ;
10 0 TEST#3
```

This illustrates an increment greater than 1. You can try different end values -- say, 11, 12, and 13 -- and you'll see that the loop always stops short of the end value.

```
: TEST#4 DO I . -2 +LOOP ;
0 10 TEST#4
```

This illustrates a negative increment. You'll see that the loop *will include* the end value if it can, but it will not go past it. (Try end values of 1 and -1 instead of 0.)

Remember: The most common mistake made with DO loops is to get the order of the start and end values backwards. The "start" value is the last thing put on the stack.

1.10. Comparisons

Now that you have the ability to change the flow of your program based on a condition, you need some operators to create those true or false flags. IsoMax has four operators which will let you *compare* two numbers:

а	b	=	returns true if a=b
а	b	>	returns true if a>b (signed numbers)
а	b	<	returns true if a <b (signed="" numbers)<="" td="">
а	b	U<	returns true if a <b (unsigned="" numbers)<="" td="">

Here "a" and "b" refer to any two numbers on the stack. "b" is on top of the stack, exactly as though you had typed the numbers "a b" at the keyboard.

You'll note that there isn't a U> operator. We'll see in the next section how to construct one. There also isn't a U= operator, since = works for both signed and unsigned numbers.

IsoMax also has three operators which let you examine a single number:

а	0=	returns true if a is zero
а	0>	returns true if a is greater than zero
а	0<	returns true if a is less than zero (negative)

Finally, there is an operator which will turn true to false, and false to true:

a NOT *logical* inverse of a

Do not confuse NOT with INVERT. INVERT is a bitwise operator, which individually inverts all 16 bits of a value on the stack. NOT is a logical operator, which returns true (a nonzero value, actually \$FFFF) if the value on the stack is false (zero), and returns false (zero) if the value on the stack is true (any nonzero value). INVERT is for bits. NOT is for true/false values.

Of course, these comparisons are also useful in IsoMax state machines. The phrase between CONDITION and CAUSES must leave a true/false value on the stack. How this value is produced is up to you; it might come from testing an input bit, or it might come from comparing two numeric values.

1.11. Stack Operations

For most short pieces of code, like IsoMax CONDITION and CAUSES phrases, you'll only have one or two things on the stack, and they'll be in the right place at the right time. But sometimes values get put on the stack in the wrong order, or you have an extra value that you don't need, or maybe you need a value twice. To handle these little details, IsoMax provides an assortment of stack operators.

Another word about stack notation: recall that we use the notation

a b c

to signify that there are three values on the stack, with "a" on the bottom, "b" in the middle, and "c" on the top. (This is how they'd be on the stack if you typed three numbers on the command line, in the same left to right order.) Since the stack operators rearrange the values on the stack, we need "before" and "after" pictures to illustrate their operation. The common notation for this is

abc---xyz

where "a b c" is the stack data *before* the operation, and "x y z" is the stack data *after* the operation. We'll see how this works in a moment.

First, let's look at words that get rid of items on the stack:

Stack effect (before after)
a
a b

DROP simply takes whatever is on top of the stack, and gets rid of it. 2DROP gets rid of the top two items on the stack. Anything else that may be deeper on the stack is unaffected, so it's not shown in the stack notation. So, if you have 1 2 3 on the stack and you do a DROP, you'll wind up with 1 2 on the stack...only the 3 is DROPped.

There are also words to duplicate items on the stack:

DUP a --- a a OVER a b --- a b a

DUP just takes whatever is on top of the stack, and makes a second copy of it on the stack (so then you have two of them). You might want to do this if you need to test if a value is zero or nonzero, and then if it's nonzero, perform some computation with it.

OVER is trickier: it makes a copy of the *second* item on the stack, and pushes that copy onto the top of the stack (pushing everything else down). If that sounds confusing, just remember that OVER takes "a b" and gives you "a b a".

To rearrange the values on the stack, you can use:

SWAP	а	b	b	а		
ROT	а	b	C	b	С	а

SWAP just "swaps" the top two stack items (the second becomes first, and the first becomes second). ROT is short for "rotate"; it rotates the top *three* items on the stack, such that the deepest item becomes the topmost item. If you do three rotates -- that is, ROT ROT ROT -- you'll get right back to where you started.

IsoMax has many more stack operators, and you're encouraged to look at the IsoMax Glossary in Section 18 to learn about more of them. But these six will handle the most of the manipulations you'll need to perform.

1.11.1. Example

Remember that we don't have a U> operator? It ought to look like this:

a b U> should return true if a>b (unsigned numbers)

Well, if a>b, then it follows that b<a. And we *do* have an operator for unsigned-less-than (U<). So all we need to do is reverse the order of a and b, and then use U<:

a b SWAP U< returns true if b<a, and thus if a>b

We can take this one step further, and use our ability to make definitions to add this as a new part of the IsoMax language!

: U> SWAP U< ;

This defines a new word named U> which does the same thing as SWAP U<. Presto! What's nice is that this new word becomes a part of the language, just like all the words originally "known" to IsoMax. Any place you could use the built-in word U<, you can use your new word U>. There is no distinction between "built-in operators" and "user functions"; you can add new comparison, arithmetic, and logical operators as you please. For this reason, IsoMax is called an *extensible* language -- you can add new language elements at any time.

1.12. Word list

The complete word list is found in the IsoMax Glossary (Section 18) at the end of this manual.