

Allen-Bradley

PLC-3Communication Adapter Module

(Cat. No. 1775-KA)

Introduction

- Chapter 10 describes how to write a full–duplex line driver to enable a computer to communicate to the 1775–KA's RS–232–C channel.
- Chapter 11 describes how to write a half duplex line driver to enable a computer to communicate to the 1775–KA's RS–232–C channel.
- Chapter 12 describes the network and application layers of a software driver to enable a computer to communicate to the 1775–KA's RS–232–C channel.
- Appendix A shows detailed message formats.
- Appendix B lists error codes reported by the $1775–KA$, $1771–KA$, 1771–KG, 1771– KE/KF, 1773–KA, and 1774–KA modules.
- Appendix C lists diagnostic counters stored at the 1775–KA, 1771–KA, 1771–KG, 1771–KE/KF, 1773–KA and 1774–KA modules.
- Appendix D gives detailed flow charts of an example of software logic for implementing a full–duplex protocol.

Related Documentation

Read this manual in conjunction with the documentation listed in Table 1.A and Table 1.B. Table 1.A lists related PLC–3 documentation and Table 1.B lists related Data Highway documentation.

Publication Number (Old/New No.)	Title
1775-800/1775-6.7.1	PLC-3 Installation and Operations Manual
1775-801/1775-6.4.1	PLC-3 Programming Manual
1775-806/1775-6.5.3	I/O Scanner-Message Handling Module User's Manual
1775-900/1775-2.1	PLC-3 Controller Data Sheet
1775-901/1775-2.2	PLC-3 Main Processor Module Data Sheet
1775-902/--------	PLC-3 Memory Organization Data Sheet
1775-904/1775-2.4	Power Supply Data Sheet
1775-908/1775-2.6	PLC-3 Memory Modules Data Sheet
1775-910/1775-2.8	PLC-3 Main Chassis Data Sheet

Table 1.A Related PLC-3 Documentation

Terminology

In this manual you will read about the various commands the 1775–KA module can send and/or receive. To distinguish between commands, we use some of the following terms:

- a **protected** command can read or write only specified areas of PC data table. A switch on the PLC, PLC–2 Family, and PLC–4 Controllers determines if the PC will accept only protected commands from another PC or an RS–232–C device. When you use a protected command, you may have a limited area that you can read or write in the other station's memory.
- an **unprotected** command can read or write into any area of PC data table. A switch on the PC that receives the commands determines if the PLC, PLC–2 Family, and PLC–4 controller will accept unprotected commands from another PC or an RS–232–C device.
- **privileged** commands are sent by intelligent RS–232–C devices only. Such devices include computers and intelligent terminals. Allen–Bradley PC's do not send privileged commands, but receive and reply to them. A privileged command can read or write into any area in the memory of a PC, whether or not switches on the PC have been set to allow it to receive only protected commands. The term physical

command is sometimes used synonymously to mean privileged command.

non–privileged commands include any command that both PC's and \blacksquare RS–232–C device can send. The non–privileged commands include the protected write and unprotected read and write commands. The non–privileged commands are also referred to as **"PLC/PLC–2 type"** commands.

Figure 1.1 illustrates the front of the 1775–KA module. The module has the following hardware features: Module Description

- Self–test diagnostic indicators
- Thumbwheel switch for setting identification number
- Two ports– one for Data Highway and one for RS–232–C communication
- Two sets of indicators one for each port
- Switches for selecting fault responses and communication option

In addition, the module provides the following software features:

- **Programmable configuration parameters**
- Command language that allows for complex logic decisions, looping, and nesting
- Symbolic representation of data and addresses
- Embedded arithmetic expressions and logic operations
- Decimal, octal, or BCD (binary coded decimal) data entry

Specifications

Table 1.C lists the specifications for the 1775–KA module.

Applications

As already mentioned, the 1775–KA module serves two main purposes:

- Interfacing the PLC–3 processor with the Allen–Bradley Data Highway
- Interfacing the PLC–3 processor with an intelligent $RS-232-C$ device

You can use the module for both of these purposes simultaneously.

In Data Highway applications, the module serves as an interface between the PLC–3 programmable controller and the Allen–Bradley Data Highway. The Data Highway is an industrial communication network that links together as many as 64 distinct stations. Each station can consist of a programmable controller (such as the PLC–3), a computer, or an intelligent RS–232–C device. The central trunkline of the Data Highway may be up to 10,000 feet long, and each station may be as far as 100 feet from the trunkline. Figure 1.2 gives an example of a Data Highway configuration.

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The PLC–3 can support multiple 1775–KA modules in the same PLC–3 chassis. This provides the PLC–3 with concurrent access to several independent Data Highways.

The 1775–KA module can also serve as an interface between the PLC–3 programmable controller and an intelligent RS–232–C compatible device or any Allen–Bradley PC and its Data Highway module. Some examples of this application of the module are the following:

- Interfacing two PLC -3 controllers through a modem link
- Interfacing a PLC–3 controller with a computer (either directly or through modems)
- Interfacing a PLC–3 controller with a remote Data Highway through a modem link
- Interfacing a PLC–3 controller as a slave station on a multipoint modem link
- Interfacing a PLC–3 controller on a point–to–point link with $PLC-2$ Family processor through a 1771–KG module (The 1772–LR processor is not supported in this configuration.)

Figure 1.3 shows the 1775–KA module in a typical modem application.

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Chapter **2**

Installation

Rotate the thumbwheel to select the desired identification number.

Figure 2.1 Front View of 1775-KA Module

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If there is only one 1775–KA module in the PLC–3 chassis, set its thumbwheel switch to the number 1. If there are multiple 1775–KA modules in the same PLC–3 chassis, set their thumbwheel switches to consecutive numbers, starting with the number 1. You may write the selected number in the space provided beside the thumbwheel switch.

CAUTION: To guard against unpredictable operation of the PLC–3 processor, do not change the setting on any thumbwheel switch while the 1775–KA module is powered–up.

Option Switches

Figure 2.2 shows a set of four option switches on the bottom edge of the 1775– KA module. Switches 1 and 2 are used when the PLC–3 controller is programmed to operate in a backup configuration. Switch number 1 determines whether or not a fault in the 1775–KA module will cause the primary PLC–3 controller to switch over to the backup PLC–3. Switch number 2 determines whether or not the 1775–KA module will disable its Data Highway port when the PLC–3 becomes deactive. Switch 3 is for RS–232–C communication. Switch 4 is reserved for future use and should always be left open (up, or off). Use Table 2.A below to determine the appropriate switch setting:

Figure 2.2 Option Switches

Table 2.A 1775-KA Switch Settings

If this switch:	ls:	Then
1	OPEN	the PLC will switch over to backup whenever one of the following fault conditions occurs: 1. The 1775-KA module tries to hold control of the PLC-3 backplane for more than 138 microseconds. 2. The 1775-KA module experiences a execution timeout of more than 32 milliseconds 3. The 1775-KA module experiences an internal stack overflow 4. The 1775-KA module experiences severe Data Highway communication problems.
1	CLOSED	the primary PLC-3 will not switch to backup when a fault occurs with the 1775-KA module.
$\overline{2}$	OPEN	the 1775-KA module will disable is Data Highway port whenever the primary PLC-3 controller becomes deactive. The module will no longer be able to transmit or receivemessages through its Data Highway port. Also, setting switch 2 to open enables the backup operation feature.
\mathfrak{p}	CLOSED	the Data Highway port on the module will remain active if the primary PLC-3 becomes deactive.
3	OPEN	the module may be connected up to 7,000 cable feet away from a 1771-KF, a 1771-KG, 1773-KA or another 1775-KA module. In addition to setting switch 3 to the open position, you must also set switch 2 to closed position. This makes pin 25 on the RS-232-C port of the 1775-KA module active (refer to figures 2.8 to 2.10). Note that switch 3 must always be closed for communication with an RS-232-C device other than a 1771-KF, 1771-KG, 1773-KA, or 1775-KA module.
3	CLOSED	the MODEM INTERFACE port of the 1775-KA module may be connected to a standard RS-232-C device that is located within 50 cable feet of the module.
4	OPEN	Switch 4 is reserved for future use and should always be left open.

Module Placement

After setting the thumbwheel switch, insert the module into any one of the module slots in the PLC–3 processor chassis. Whenever you power–up the processor, the module will receive power also.

Indicators

There are three sets of LED indicators on the front of the 1775–KA module (Figure 2.1). The first group, labeled SELF–TEST, indicates the result of internal diagnostic tests that the module continuously performs on its own hardware and firmware. The second group, labeled MODEM INTERFACE, indicates the status of communication through the module's RS–232–C port. The last group, labeled DATA HWY, indicates the status of communication through the module's Data Highway port.

Table 2.B. tells what each indicator means.

Table 2.B LED Indicators

Data Highway Cable Connections

There are two cable connectors, or ports, on the front of the 1775–KA module (Figure 2.1). The bottom port, labeled DATA HWY., is for connection to the Allen–Bradley Data Highway. If you are using the 1775–KA module in a Data Highway application, plug the Data Highway dropline cable into this port. For details on the installation of the Data Highway cable, refer to the Data Highway Cable Assembly and Installation Manual (publication 1770–810).

RS-232-C Cable Connections

The RS–232–C port, labeled MODEM INTERFACE on the 1775–KA module, can interface with any RS–232–C device that is capable of understanding and generating the communication protocol described in this chapter. Some typical RS–232–C applications are:

Interfacing two PLC–3 controllers through a modem link (Figure 2.3)

Figure 2.3 Linking Two PLC-3 Controllers

> **NOTE:** Modems required only for distances greater than 50 feet.

 Interfacing a PLC–3 controller with a computer, either directly or through modems (Figure 2.4)

Figure 2.4 Linking a PLC-3 Station to a Computer

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■ Interfacing a PLC–3 controller with a remote Data Highway through a modem link (Figure 2.5)

10006–I

■ Interfacing a PLC–3 controller to a PLC–2 Family processor through a 1771–KG module in a point–to–point link (Figure 2.6)

Figure 2.6 Linking a PLC-3 to PLC-2 Family Controller

10007-I

■ Interfacing a PLC–3 controller as a slave station on a multipoint modem link (Figure 2.7)

The first four applications above use the module's RS–232–C port in the unpolled mode, while the last application uses the polled mode. You can select the mode of operation and other characteristics of the RS–232–C port through the LIST function.

Each mode of operation requires a different communication protocol. The unpolled mode uses full–duplex protocol (chapter 10) while the polled mode uses half–duplex protocol (chapter 11). In general, full–duplex protocol gives faster data throughput but is more difficult to implement; half–duplex protocol is easier to implement but gives slow data throughput.

NOTE: In other Data Highway documentation, full–duplex protocol might be referred to as DFI protocol, and half–duplex protocol might be referred to as polled–mode protocol.

Hardware Interface

The modem interface is based on EIA RS–232–C and related standards. This interface should be compatible with most dedicated and dial–up network RS–232 modems.

Mechanical

The RS–232 connector on the 1775–KA module is a 25–pin male connector.

Electrical

Input and output levels on the RS–232 connector conform to the RS–232–C standard. The transmitter has increased capability to drive a 7,000 foot isolated lines. This number depends on baud rate and refers to only direct wire connections. (Refer to Table 2.C.)

The receiver is designed to sense the signal generated by a similar transmitter, and is electrically isolated from all other circuitry on the module. It consists of an opto–isolater circuit with an input and return connection at the RS–232 connector. All other signals on the RS–232 connector are driven and received by standard RS–232 interface circuits, and have a maximum drive capability of 50 feet.

Pinout

The necessary RS–232–C port connections are described in Table 2.D below:

Table 2.D RS-232-C Port Connections

- TXD (transmitted data) caries serialized data. It is output from the RS–232 connector.
- RXD (received data) is serialized data input to the RS–232 connector. RXD and RXDRET are isolated from the rest of the circuitry on the module.
- RTS (request to send) is a request from the RS–232 connector to the modem to prepare to transmit. It typically turns the data carrier on. When you select the full duplex mode RTS is always asserted. When you select the half duplex mode RTS is turned on when the module has permission to transmit; otherwise it is off.
- CTS (clear to send) is a signal from the modem to the RS–232 connector that the carrier is stable and the modem is ready to transmit. The module will not transmit until CTS is true. If CTS is turned off during transmission, the module will stop sending until CTS is restored.
- DTR (data terminal ready) is a signal from the RS–232 connector to the modem to connect to the phone line (that is, "pick up the phone"). The module will assert DTR all the time except during the phone hangup

sequence. Some modem will not respond to DTR until the phone rings, while others will always pick up the phone whether it is ringing or not.

- DSR (data set ready) is a signal from the modem to the RS–232 connector that the phone is off–hook. (It is the modem's answer to DTR). The module will not transmit or receive unless DSR is true. If the modem does not properly control DSR, or if no modem is used, DSR must be jumpered to an RS–232 high signal at the RS–232 connector. (It can be jumpered to DTR).
- DCD (data clear ready) is a signal from the modem to the RS–232 connector that the carrier from another modem is being sensed on the phone line. It will not be asserted unless the phone is off–hook. Data will not be received at the RS– 232 connector unless DCD is true. In the full duplex mode the module will not transmit unless DCD is true. If the modem does not properly control DCD, or if a modem is not being used, DCD must be jumpered to DTR at the RS–232 connector.
- TXDRET (transmitted data return) is the return signal for TXD. It is connected to module logic ground through a resistor.
- RXDRET (received data return) is the return signal for RXD. It is connected to the isolated receiver, and is isolated from all other circuitry on the module.

Connections To The RS–232 Port

Connection to the RS–232 port of the 1775–KA can be one of two types:

- Short line (50 feet or less)
- Isolated long line (between 50 and 7,000 feet)

For short lines, the connection may be either direct or through modems.

You connect an intelligent, RS–232–C compatible device to an interface module by attaching a cable to both the device and to the module socket labeled RS– 232–C CHANNEL. The RS–232–C device may be another Allen–Bradley communication interface module or another manufacturer's device. For a standard RS–232–C connection, the cable should be no longer than 50 feet. If your RS–232–C device has an Allen–Bradley line driver/receiver, you may use a cable up to 7,000 feet long.

If you want to connect the 1775–KA module to a 1771–KG or 1771–KE/KF module through the RS–232–C channel, use the cabling pinout diagram (Figure 2.8) to construct your own cable.

1

Conductors 2 and 7, 3 and 25 must be twisted pairs for distances longer than 50 feet.

Set switch 3 (on the 1775–KA) OFF when the module is communicating **2** Set switch 3 (on the 1775–KA) OFF when the module is communicating
⊌with another Allen-Bradley device.

If you want to connect the 1775–KA module to a 1775–KA module through the RS– 232–C channel, use the cabling pinout diagram (Figure 2.9) to construct your own cable.

Figure 2.9 Connection to Allen-Bradley 1775-KA Module

Connect the Shield at One End Only

- Conductors 2 and 7, 3 and 25 must be twisted pairs for distances longer than 50 feet.
- Set switch 3 (on the 1775–KA) OFF when the module is communicating ■ Set switch 3 (on the 1775–KA) OFF when the module is communicating
— with another Allen-Bradley device.

If you want to connect the 1775–KA module to a modem or computer, use the cabling pinout diagram (Figure 2.10) to construct your own cable.

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RS–232–C **™** CHANNEL Connector of 1775–KA Module

1 Set Switch 3 ON to ground pin 25.

Private lines are permanently connected phone lines used with modems. Dialup is not needed. Usually the modem hold the handshake lines in the proper states.

The RS–232 port can be connected to standard American dial–up modems and some European modems. Other European standards specify that the DTR signal will cause the modem to answer the phone whether it is ringing or not, causing the phone to always be "busy". Since the modem port asserts DTR while waiting for a call, it cannot be used with such modems.

The types of dial–up network modems that can be used are classified into the following types:

- **Manual**: these are typically acoustically coupled modems. The connection is established by human operators at both ends, who then insert the handset into couplers to connect the computers.
- **DTE–controlled answer**: these unattended modems are directly connected to the phone lines. A module controls the modem via the

DTR, DSR, and DCD signals. It incorporates timeouts and tests to properly operate these types of modems.

Auto–answer: these modems have self–contained timeouts and tests, and can answer and hangup the phone automatically.

The modem port has no means to control an auto–dial modem, although it is possible that it can be used in conjunction with a separate auto–dialer.

Answering

The module continually asserts DTR when it is waiting for a call. Under this condition the modem will answer a call and assert DSR as soon as ringing is detected. The module does not monitor the RING indicator in the RS–232 interface. Once DSR is detected the module starts a timer (around 10 seconds) and waits for the DCD signal. When DCD is detected communication can start.

If DCD is not detected within the timeout, the module turns DTR off. This causes the modem to hangup and break the connection. When the hangup is complete the modem drops the DSR line. This causes the module to reassert the DTR line and wait for another call. This feature protects access to the phone if someone calling a wrong number reaches this station.

Once DCD is detected the module continues to monitor the DCD line. If DCD goes false the timeout is restarted. If DCD is not restored within the timeout, the hangup sequence is initiated. This feature allows the remote station to re–dial in the event the connection is lost by the phone network.

Note that this handshaking is necessary to guarantee access to the phone line. If this handshaking protocol is defeated by improper selection of modem options, or jumpers at the connectors, the modem may answer a call, but if the connection is lost the modem will not hangup. It will be impossible for the remote station to reestablish the connection because it will get a busy signal.

Character Transmission

Data is sent serially over the RS–232 interface, one eight–bit byte at a time. The transmission format conforms to ANSI X3.16, CCITT V.4, and ISO 1177, with the exception that the parity bit is retained while extending the data length to eight bits.

Figure 2.11 LIST Menu for 1775-KA Module

You access the LIST function by typing the word LIST and press the ENTER key. After accessing the LIST function, select option 6 MODULE STATUS from the SYSTEM–MODE MENU. LIST then presents you with a menu that describes the modules in your system. The menu varies according to the modules in your PLC–3. A typical menu might be:

MODULES: 1 01 1775–ME8 A/A 2 02 1775–ME8 A/A 3 03 1775–L3 A/A 4 04 1775–S4A B/A 5 05 1775–KA A/E 6 06 1775–LX A/A 7 07 1775–LX A/A 8 08 1775–S4B B/A 9 09 1775–S4B B/A ENTER NEXT<

Under MODULE STATUS, select the option for 1775–KA. At this point, LIST presents you with the following menu for the 1775–KA module:

DATA HWY COMM. ADAPTER–nn CHASSIS cc SLOT ss 1 MODULE OPTIONS 2 DATA HIGHWAY PORT 3 MODEM PORT ENTER NEXT>

In the above and all following menus, "nn" represents the thumbwheel setting of the 1775–KA module, "cc" represents the chassis number, and "ss" represents the number of the chassis slot containing the module. For more information about the LIST function, refer to Publication 1775–800, PLC–3 Installation and Operation Manual.

Module Options

Selecting option 1 MODULE OPTIONS from the above menu (section titled Programmable Configuration Parameters) causes LIST to present the following menu:

KA – nn MODULE OPTIONS 1 TIMEOUT 2 SEND UNPROTECTED

3 ACCEPT UPLOAD/DOWNLOAD 4 ACCEPT WRITES 5 BACKUP OPERATION 6 PLC–2 MASK ENTER NEXT>

This menu allows you to select options that apply equally to both the modem port and the Data Highway port of the 1775–KA module. These options are described below.

Timeout

The timeout is the maximum amount of time that the 1775–KA module will wait for another station to reply to one of its messages. The allowed entries are 0 to 9999, expressed in increments of 1/10 second. LIST displays the timeout as "xxxx/10 SEC". Thus, if you enter a timeout value of 100, the timeout period will be 10 seconds and will be displayed as 100/10 SEC.

The same timeout setting applies to both the Data Highway and the modem ports. The default timeout setting is 5 seconds, displayed as 50/10 SEC.

The timeout period applies to each individual transmission. Because of their size, some messages consist of several packets of data. Each message packet requires a separate transmission. Therefore, the timeout is restarted for each packet.

If the 1775–KA module waits longer than the timeout period for a reply to one of its messages, it generates an error code of 37 (Appendix B). The module then resumes executing the current message procedure at the line following the one in which the timeout occurred.

LIST keeps you at this timeout level and allows you to make repeated changes to the timeout value. To return to the preceding (next highest) level of LIST, press the ENTER key again without entering a new timeout value.

Send Unprotected

This option determines whether or not the 1775–KA module will be able to send unprotected command messages to other stations. If you select option 2 SEND UNPROTECTED, the 1775–KA module will be able to send both protected and unprotected commands.

You can use an unprotected command to read or write to any area of a PC data table. You can use a protected command, however, to write only to those areas of a PC data table specified by the PC that receives the command. For more information on protected and unprotected commands, see section titled Access Privileges, chapter 3.

If you do not select (enable) this option, the module will be able to transmit only protected commands. At initial power–up, the module enables the SEND UNPROTECTED option by default.

Accept Upload/Download

This option determines whether or not the 1775–KA module will be able to execute upload and download commands sent to it by a computer. If option 3 ACCEPT UPLOAD/DOWNLOAD is selected, the module will be able to execute both upload and download commands. You send a sequence of upload and download commands when you want to transfer the memory of the PLC–3 to another station, or to transfer the memory of another station to a PLC–3.

If this option is not selected (enabled) the module will not be able to execute either of these two types of commands. For a description of upload and download commands, refer to Appendix A. At initial power–up, the module enables the ACCEPT UPLOAD/DOWNLOAD option by default.

Accept Writes

This option determines whether or not the 1775–KA module will accept write–type command messages from a remote Data Highway station when the local PLC–3 processor's memory protect keyswitch is on. If option 4 ACCEPT WRITES is selected, the module will accept write commands regardless of the setting of the PLC–3's memory protect keyswitch. If this option is not selected (enabled) the module will accept write commands when the memory protect keyswitch is off but will reject the write commands and return an error code of 86 if the memory protect

keyswitch is on. At initial power–up, the module enables the ACCEPT WRITES option by default.

Backup Operation

This option determines whether or not a pair of 1775–KA modules will provide backup for each other. Enable option 5 (BACKUP OPERATION) for both the primary and backup 1775–KA modules to enable backup operation as described in section 2.3 (backup configurations). If you make no selection for option 5, backup operation is disabled by default. The revision C or earlier version of the module does not have the BACKUP OPERATION option.

PLC–2 Mask

This option determines whether or not the 1775–KA module will mask out the upper octal digit of the source address when receiving a PLC–2 type command from another station. If you enable option 6 (PLC–2–MASK) the module mask out the upper digit of the address for selecting the input file. This causes stations with common second and third digits of their address to access the same input file. For example, stations 023, 123, 223, and 323 would all access input file 023.

If you disable option 6, each station accesses a unique input file with the same number as the station number. For example, station 123 would access input file 123; station 223 would access input file 223. If you make no selection for option 6, the PLC–2 MASK option is disabled by default. The revision D or earlier version of the module does not have the PLC–2 MASK option.

Data Highway Port

Selecting option 2 DATA HIGHWAY PORT from the above menu (section titled Programmable Configuration Parameters) causes LIST to present the following menu:

KA – nn DATA HIGHWAY PORT 1 ENABLE/DISABLE PORT 2 STATION NUMBER 3 BAUD RATE ENTER NEXT>
This menu allows you to select options that apply to only the Data Highway port of the 1775–KA module. These options are described below.

Enable/Disable Port

This option determines whether or not the 1775–KA module can communicate over the Data Highway. you must select the ENABLE option in order to allow communication to take place. If you make no selection, the PLC–3 disables this port by default.

Note, however, that you cannot use LIST to change any other parameters of the Data Highway port unless you first DISABLE the port. After you are done entering parameters through LIST, don't forget to ENABLE the Data Highway port again.

Station Number

This option selects the number by which the PLC–3 station is identified on the Data Highway. Allowable station numbers are 1 to 376 octal. In particular, note that the number 377 is illegal. Entering 377 as the station number will automatically disable the 1775–KA module, and you will not be able to ENABLE it again through LIST until you select a different station number. If you make no selection, the PLC–3 assumes the illegal address 377 by default.

Baud Rate

This option specifies the communication rate over the Data Highway. A communication rate of 57,600 baud is recommended.

Modem Port

Selecting option 3 MODEM PORT from the above menu (section titled Programmable Configuration Parameters) causes LIST to present the following menu:

KA–NN modem port 1 ENABLE/DISABLE PORT 2 STATION NUMBER 3 BAUD RATE 4 COMMUNICATION MODE

5 EVEN PARITY 6 SEND EMBEDDED RESPONSES ENTER NEXT>

This menu allows you to select options that apply to only the modem port of the 1775–KA module. These options are described below.

Enable/Disable Port

This option determines whether or not the 1775–KA module can communicate through its RS–232–C port. You must select the ENABLE option in order to allow this communication to take place. If you make no selection, the PLC–3 disables the RS–232–C port by default.

Note, however, that you cannot use LIST to change any other parameters of the RS–232–C port unless you first disable the port. After you are done entering parameters through LIST, don't forget to enable the RS–232–C port again

Station Number

This option selects the number by which the PLC–3 station is identified on an RS–232–C communication link. In particular, note that the number 377 is illegal. Entering 377 as the station number will automatically disable the RS– 232–C port of the 1775–KA module, and you will not be able to enable it again through LIST until you select a different station number. If you make no selection, the station number will be 377 by default.

Baud Rate

This option specifies the communication rate over the RS–232–C port. The choices are:

- 110 baud 2400 baud
-
- -300 baud -4800 baud
	-
-
-
- 600 baud 9600 baud
- 1200 baud 19200 baud

For long–line communication, the maximum allowed rate is 4800 baud. The default baud rate is 1200 baud (see Table 2.D).

Communication Mode

This option determines whether the RS–232–C port of the 1775–KA module can operate in a half–duplex (polled) or full–duplex (unpolled) mode. Select full– duplex for point–to–point communication through the RS–232–C port. Select half–duplex if the 1775–KA module is installed as a slave station on a multipoint modem link. If you make no selection, the 1775–KA selects the full– duplex (unpolled) mode by default.

Even Parity

This option determines what kind of parity check is used for all communications through the RS–232–C port. If option 5 EVEN PARITY is selected, the 1775–KA module will test for even parity in all communications through its RS–232–C port. If this option is not selected, the module will not perform any parity checking. At power–up, the PLC–3 disables the even parity option by default.

Send Embedded Responses

This option determines whether or not the 1775–KA module will be able to send embedded responses through its RS–232–C port. Responses are acknowledgments (ACKs or NAKs) to messages received from other stations. An embedded response is one whose characters are transmitted between the bytes of a regular message. In this way, the response to a previously received message is transmitted along with a new message. At power–up, the PLC–3 disables the embedded responses option by default.

The 1775–KA module can combine with the PLC–3 processor to form a backup system. System backup is described in greater detail in the PLC–3 Programmable Controller Backup Concepts Manual (pub. no. 1775–6.3.1). The following discussion is an overview of system backup and the role of the 1775–KA in various backup procedures. Backup Configurations

There are two possible backup configurations for the 1775–KA module:

- Two 1775–KA modules in the same PLC–3 controller
- One 1775–KA module in a primary PLC–3 controller and another 1775–KA in a backup PLC–3 controller

The first configuration provides backup for the 1775–KA module itself. Here, both 1775–KA modules are always active, and both are independent stations on their communication network. Therefore, each 1775–KA

module must have its own unique station number. If you want to send the same message through both 1775– KA modules, you must program the two separate message instructions.

The second configuration provides system backup for the PLC–3 controller.

If the 1775–KA modules are Rev. C or earlier, you:

- **1.** Assign different station addresses to each communication adapter module.
- **2.** If the programs in the primary and backup processors are identical:
- you must be sure that all information sent to the primary processor is also sent to the backup processor.
- you must examine the run/backup bit (data table status section, file 0, word 3, bit 17) on every rung used to transmit data. This bit is set in the primary processor and reset in the backup processor. Examining it helps to prevent sending duplicate messages over the Data Highway.

If the 1775–KA modules are Rev. D or later, you can:

follow the two steps described above

or

■ select BACKUP OPERATION with the PLC–3 LIST function.

To implement backup operation, follow these steps:

1. Using the option switches on both 1775–KA modules, set switch 2 to the OPEN position. Recall (section titled Option Switches) that this causes the module to disable its Data Highway port if the PLC–3 becomes deactivated.

WARNING: If you do not set these switches OPEN on both 1775–KA modules, these modules will assume the same station address when the primary PLC–3 becomes deactive. This may shut down communication on the Data Highway, and unexpected machine motion may result.

- **2.** Use the LIST function to disable the modem and Data Highway ports.
- **3.** Use the LIST function to select BACKUP OPERATION for both the primary and backup 1775–KA modules. For more information, see the PLC–3 Installation and Operation manual (publication 1775–6.7.1). Thus, when you select BACKUP OPERATION, the condition appears like this:

5*BACKUP OPERATION

4. Use the LIST function to assign the same station address to the modules for the primary PLC–3 and the backup PLC–3 processor.

You can never give the 1775–KA module a station address of 3778, and when you select the BACKUP OPERATION, you can not give the module an address of 2778.

Because you have chosen BACKUP OPERATION, the module in the backup PLC–3 will assume an address other than the address you assign it with the LIST function (Figure 2.12).

You assign both modules an identical address.

At switchover, the address for the backup module returns to the station address you assigned to it with the LIST function.

5. Use the LIST function to enable whichever (Data Highway or Modem) port you are using to connect primary to backup.

Figure 2.12 How addresses of the primary and backup PLC-3 controllers change during switchover.

Using Manual Switchover

After you select the BACKUP OPERATION for a rev. D. (or later) 1775–KA module, you may choose to use your PLC–3 backup system for manual switchover. In manual switchover, you must initiate the switchover by changing the position of a switch in a backup cable. (Refer to the PLC–3 Programmable Controller Backup Concepts Manual, pub. 1775–803, for more details.) You must be sure to turn off the faulted PLC–3 processor before you begin the switchover, however.

You can program the MSG instruction to execute a message upon switchover or you can send commands to the backup PLC–3 processor (Figure 2.13). As long as you get responses from the backup processor, switchover has not yet occurred.

Using Automatic Switchover

After you select the BACKUP OPERATION for a rev. D. (or later) 1775–KA module, you may want to use automatic switchover for your PLC–3 backup system. During automatic switchover:

- the 1775–KA module for the primary PLC–3 processor disables its Data Highway port.
- the 1775–KA module for the backup PLC–3 processor becomes the address that you selected with the LIST function (rather than the corresponding address it received during the BACKUP OPERATION).

NOTE: You cannot select the BACKUP OPERATION for a multidrop modem applications because the modem port will not become disabled after a PLC–3 processor fault regardless of the switch settings on the module.

Run/Backup Bit

It is important to alert the proper personnel when a switchover occurs. One way you can provide such indication is by having your program monitor the run/backup bit (data table status section, file 0, word 3, bit 17) and turn on alarms or lights when the status changes from backup to run. This bit is set in the primary processor and reset in the backup processor.

It is also possible to link a single PLC–3 controller to more than one Data Highway by installing multiple 1775–KA modules in the same PLC–3. In this configuration, each 1775–KA module connects to a different Data Highway, and each has a unique station number on its associated highway. However, all the 1775–KA modules in the same PLC–3 controller can have either the same or different station numbers. Multiple 1775-KA Modules in One PLC-3

CAUTION: If such a PLC–3 station is communicating through a PLC/PLC–2 buffer file and all of the stations' 1775–KA modules have the same station number, then all of these modules will transfer data through the same buffer file. This can cause unpredictable results if several 1775–KA modules try to read or write to the buffer file at the same time.

When such a PLC–3 station transmits a command message to a remote Data Highway station, the thumbwheel number specified in the PLC–3 message instruction (section titled PLC–3 Stations) determines which 1775–KA module actually transmits the command.

Data Highway Communication

Each station consists of some type of processor and a **station interface module**. The station interface module enables the processor to communicate with other stations on the Data Highway. The 1775–KA module is the station interface module for the PLC–3 processor. Table 3.A lists all possible combinations of station interface modules and processors.

Chapter

Communication Terminology

Stations communicate with each other by sending **messages** over the Data Highway. There are two types of messages:

- Command messages
- Reply messages

A **command message** either gives (writes) data to, or requests (reads) data from, one station to another. A **reply message** is a station's response to a command message.

Command messages are generated by message procedures that you program into the 1775–KA module. Execution of a message procedure is controlled by the message (MSG) instruction in the PLC–3 ladder diagram program. When a 1775–KA module receives a command message from another station, the module automatically generates the appropriate reply message.

As points of reference, we can talk about local and remote stations. The local station is the one currently initiating some action, or the one we are currently doing something with. All other stations are then **remote**.

We can also describe stations in terms of their relationship to a message. The **transmitting** station is the one sending the message, and the **receiving** station is the one that gets the message. A station that transmits a command message is called a **command** station, and a station that transmits a reply message is called a **reply** station.

You can send either:

- a single message procedure command (Chapter 6) that may be up to 76 characters long.
- the name of a Data Highway message procedure which contains a group of commands and is stored in the 1775–KA module

You specify the station that will receive the command with a PLC–3 extended address. This address always takes the form:

E2.5.nn

where

E2 specifies that this command addresses the module status area of PLC–3 memory

5 specifies that you are sending the message instruction through the 1775–KA module

nn is replaced with the thumbwheel setting on your particular 1775–Ka module

To enter a message instruction, complete the steps below:

1. Enter a condition that, when true, will activate the message instruction. In Figure 3.1, we used an examine–on for input word 0012₈, bit 01.

Figure 3.1 Levels of Programming in Data Highway Communication

3) Data Highway

Command message transmitted to station 24

10014–I

- **2.** Press the message instruction key.
- **3.** Specify message type 1.
- **4.** Choose a control file word where status information about the message command can be stored. In our Figure 3.1, we used binary file 200, word 200.

Data transfers can be either solicited or unsolicited, depending on whether they are initiated by the local or a remote station, respectively. Either type of station initiates the data transfer by issuing a command message. If the local station issues the command message, the corresponding reply message is said to be **solicited** because the local station has solicited, or requested, the data contained in the reply message. If a remote station issues the command message, that message is said to be **unsolicited**.

For solicited messages, a local station receives data from a remote station during a **read** operation. The local station sends data to a remote station during a **write** operation.

For unsolicited messages, a local station receives data from a remote station during a **write** operation. A local station sends data to a remote station during a **read** operation.

In read operations, the command message requests the data transfer, but the corresponding reply message actually contains the data being transferred. In write operations, the command message contains the data being transferred, and the reply message merely reports the status (receipt or non–receipt) of the transfer.

The PLC–3 processor must be free to control its own processes at the same time that the 1775–KA module is communicating over the Data Highway. For this reason, both the processor and the module have their own programs and programming languages. Figure 3.1 illustrates how these two programming levels (processor and module) interrelate. Levels of Programming

PLC-3 Program

The first link in the communication process is your PLC–3 ladder diagram program. You send a Data Highway command message by means of the message (MSG) instruction. Figure 3.1 shows a typical MSG instruction.

When the rung becomes true, the message instruction begins sending command(s) across the Data Highway. At the same time, bits in a control file word change their state (Table 3.B) to reflect the status of the command. Even if the rung becomes false, the message command will continue to send commands across the highway.

Table 3.B The Status of Bits in a Control File Word

- **5.** Enter an extended address for the channel. In our Figure 3.1, we address the module status area of memory, specify the 1775–KA module, and a thumbwheel setting of 1.
- **6.** Enter either a command or a command procedure. In Figure 3.1, we entered the command procedure, PROC_A.

Data Highway Message Procedure

As already stated, the 1775–KA module has its own programming language that consists of commands (Chapter 8). A group of related commands make up a Data Highway message procedure. These commands and message procedures determine what messages are transmitted over the Data Highway.

Data Transfers

The whole purpose of Data Highway communication is to transfer data from one station processor memory location to another. To accomplish these data transfers, you can program the assignment command into the 1775–KA module.

Chapter 6 gives the details of the assignment command. For now, let's just look at the simple example in Figure 3.2. In this example, the assignment command copies a word (16 bits) of data from the source to the destination location. The **source** of the data is always specified on the **right** of the equals sign (=), and the **destination** is always on the **left.**

Figure 3.2 Example Assignment Command

10015-I

Note that an assignment command does not destroy the data at the source location; rather, it just makes a copy of the source data at the destination location. When the assignment is executed, both source and destination will contain the same data.

There are two ways to use a data transfer command with the 1775–KA module:

- as a single command within a PLC–3 message instruction
- as one of multiple commands within a message procedure

Figure 3.3 illustrates both of these methods for the same assignment command. Note that a message instruction in the PLC–3 ladder diagram program controls execution of the command in either case.

2) as part of a message procedure

Message Procedure PROC_A

(other commands)

\$B45:21 = \$112:33

(other commands)

PLC-3 Message Instruction to Control Execution of Procedure PROC_A

10016–I

Access privileges

Not every Data Highway station can read or write to every other station. In general, read and write access privileges depend on two factors:

- type of processor at the transmitting and receiving stations
- protections set at the receiving station

The rest of this section explains how these access privileges vary according to the above factors.

PLC–3 Stations

A PLC–3 station can always **read** data from any major area of another PLC–3's memory. However, one PLC–3 station can write only to the data table area of another PLC–3 station.

In addition, a local PLC–3 station can prevent remote PLC–3 stations from writing to the local station's data table by setting a memory protection switch. At the local station, the memory protect switch can be overridden by selecting option 4 in the Module Options Menu (section titled Accept Writes, chapter 2) at the local station.

PLC/PLC–2 Stations

For communication with a PLC or a PLC–2 station, read, and write access privileges depend on switch settings at that station. For an explanation of how to set the switches for read and write access, refer to the Communication Adapter Module User's Manual (publications 1771–6.5.1 and 1774–6.5.8).

Accessing a PLC/PLC–2 Station

Access to a PLC/PLC–2 station also depends on the type of command transmitted to that station. There are two types of commands:

- protected write commands
- unprotected read and write commands

Protected write commands can only write to specified sections of the data table in a PLC/PLC–2 processor. Memory access rungs in the PLC/PLC–2 ladder diagram program specify where in the data table the PLC–3 can write data.

Unprotected commands, on the other hand, can read or write to any section of the data table at a PLC/PLC–2 station. (Again, refer to publication 1771–801 or 1774–819 for an explanation of protected and unprotected commands and memory access rungs.)

A PLC–3 station can read from any part of a PLC/PLC–2 data table. However, A PLC– 3 station cannot write to a PLC/PLC–2 if the switch settings at the PLC/PLC–2 station forbid access. If the switches at the PLC/PLC–2 station are set to accept only protected write commands, then the ladder diagram program at the PLC/PLC–2 station must contain memory access rungs to define which areas of the PLC/PLC–2 station's data table are accessible. In such a case, a transmitting PLC–3 station can write to only those data table areas defined by the memory access rungs, and only by means of protected write commands. If the switches at the PLC/PLC–2 station are set to accept unprotected write commands, a PLC–3 station can then write to any area of the PLC/PLC–2's data table by transmitting an unprotected write command (section titled Command Message Type chapter 6).

Accessing a PLC–3 from a PLC/PLC–2 Processor

While a PLC–3 processor can address any area of a PLC/PLC–2 data table, a PLC/PLC–2 reads an input file that is a part of the PLC–3 data table. That file is the PLC–3 input file with a number that corresponds to the station number of the PLC/PLC–2 station. For example, the read/write files assigned to PLC/PLC–2 stations 1 to 100 (octal) would be as follows:

PLC/PLC–2 station numbers are octal, while PLC–3 input files have decimal addresses. This means that PLC–3 input files with an 8 or 9 in their address are not used for read/write access by a PLC/PLC–2 station.

The PLC/PLC–2 station can use either protected or unprotected commands to access its assigned PLC–3 file. Note, however, that the PLC/PLC–2 station cannot access its assigned file until that file is created and allocated at the PLC–3. To create a PLC–3 file, use the CREATE command described in the PLC–3 Programming Manual (publication 1775–801).

Note that it is possible to have two PLC–3 stations communicate with each other as if they were PLC/PLC–2 stations. To do this, simply allocate the appropriate PLC/PLC–2 buffer files in the PLC–3 stations and uses the PLC/PLC–2 addressing format (section titled PLC/PLC–2 Address Specifications, chapter 4) in the assignment commands. Similarly, a computer can sent PLC/PLC–2 commands to a PLC–3 station by using the appropriate message packet formats (Appendix A).

To allow as many as 4 remote stations to access the same PLC–3 input file:

- **1.** Enable the PLC–2 MASK option in the LIST function. PLC–2 MASK is option 6 on the Module Options menu.
- **2.** Select station numbers $100₈$ apart. For example, you could use stations 010, 110, 210 and 310.

The stations will have access to the input file which matches the lower two digits of these station numbers (input file 10 in this example). When the 1775– KA module receives a PLC–2 type command, it masks the upper octal digit in order to determine which input file to access. So commands sent from stations 010, 110, 210, and 310 would all access input file 10.

PLC–4 Stations

To read or write to a PLC–4 station, you can send either protected or unprotected commands.

Switches 2 and 3 (on the second row of switches) at the 1773–KA module specify whether the PLC–4 station will accept unprotected and protected commands (respectively) through the Data Highway port of the 1773–KA module.

Switches 1 and 3 (on the third row of switches) at the 1773–Ka module specify whether the PLC–4 station will accept protected and unprotected commands (respectively) through the RS–232–C port of the 1773–KA module. In all cases, if the switch is set to the closed position, the module will accept that type of command.

General

This chapter presents some general rules for specifying data addresses in message procedures. This chapter assumes that you are already familiar with the forms and meanings of addresses in the PLC–3 and other Allen–Bradley programmable controllers. For details on these subjects, refer to the appropriate documentation listed in Table 4.A.

Chapter

Table 4.A Memory Organization Documentation

In this chapter,the addressing formats are presented in shorthand notation. The notation used is as follows:

An expression can be used in place of any of the above fields in an address.

Within the above listed fields of an address specification, numbers are interpreted as decimal (base 10) unless you indicate that they are octal (base 8). You can specify an octal number by enclosing the number in parentheses and starting it with a leading zero. For example, 17 is interpreted as decimal 17, but (017) is interpreted as octal 17, or decimal 15. Number Systems

> An exception to the above rule occurs when addressing a word in the input or output sections of PLC–3 memory. In these cases, the word address <wordaddr> is normally interpreted as an octal number, regardless of leading zeros. To express an input or output word address as a decimal value, enclose the word address within parentheses and eliminate leading zeros.

> In addressing individual bits, parentheses have no affect on the address interpretation The bit address <bit> is interpreted as an octal number if it starts with a leading zero and as a decimal number if it does not start with a zero.

Figure 4.1 illustrates these addressing conventions.

Figure 4.1 Examples of Addressing Conventions

Interpretation (expressed in decimal)	
Input file 12, word 13	
Input file 12, word 13, bit 15	
Input file 12, word 13, bit 13	
Input file 12, word 15	
Input file 12, word 13	
Integer file 43, word 13	
Integer file 43, word 13	
Integer file 35, word 13	
Integer file 43, word 13	$10017 - I$

Addresses

Data is referenced by its address in memory. In a message procedure, you must precede an address with a dollar sign. The dollar sign acts as a delimiter to tell the 1775–KA module that it has encountered a data address. Figure 4.2 illustrates this addressing format in a simple assignment command.

Figure 4.2

Assignment Command

10018-I

For data locations at remote stations, the remote station number must precede the data address. Figure 4.3 illustrates this addressing format. For communication on the Data Highway, the characters #H are required to delimit the remote station number. For RS–232–C communication through the 1775–KA module's modem port, the characters #M delimit the remote station number.

10019-I

Symbols

You can also use symbols to represent data and data addresses in message procedures. A symbol can consist of numeric digits, alphabetic characters, and the underline character (_). No other special characters are allowed. The first character in a symbol must be a letter of the alphabet.

Both upper–case and lower–case letters are acceptable in a symbol, but they are distinguished. For example, ASYMBOL and Asymbol are two different symbols.

A symbol can be any length, but it must be unique in its first 8 characters. For example, SYMBOL_A and SYMBOL_B are distinguishable in a message procedure, but NEW_SYMBOL_A AND NEW_SYMBOL_B are not. Note that indistinguishable symbols are not flagged as programming errors. Rather, indistinguishable symbols are treated as equivalents.

Certain words and character combinations cannot be used as symbols because they are reserved for special uses in message procedures. The reserved words are:

Any abbreviated form of one of the above words is also invalid as a symbol. For example, the single letter C should not be used as a symbol because it is an abbreviation of the word CREATE. Similarly, PRO is an invalid symbol.

Figure 4.4 illustrates the classification of different types of symbols. The two major classifications are:

- system symbols
- user symbols

System Symbols

A system symbol is used as either a procedure name or a symbolic address. The characters in a system symbol must conform to the general rules given above for all symbols. System symbols are delimited by the character @, which distinguishes them from user symbols.

Procedure Names

A procedure name is a way of referring to a message procedure. You assign a procedure name at the time you generate, or edit, the message procedure (Chapter 6).

One procedure can execute a second procedure simply by stating the name of that second procedure. This allows for nesting of procedures up to three levels deep.

Symbolic Addresses

A symbolic address is another way of representing the logical address of data (section titled Addresses). You can generate a symbolic address by using the CREATE command (Chapter 8). A symbolic address can be used anywhere that a logical address can be used in a message procedure. The symbolic address is stored in the system symbols area of the PLC–3 memory.

Scope of System Symbols

System symbols can be either local or global in scope. A global system symbol is known in any context. A local system symbol is known only in the context in operation at the time the symbol was generated. Context is explained in the PLC–3 Programming Manual (publication 1775–801).

At the time you generate the system symbol, you can specify whether it is to be local or global. If you do not specify the scope of the system symbol, it is assumed to be local.

Note that the terms local and global symbols should not be confused with local and remote stations. Both local and global symbols having meaning only at the station in which they were generated.

User Symbols

A user symbol represents a numeric value. You can generate a user symbol and assign a value to it by means of the assignment command (Chapter 7).

User symbols are either procedural or interprocedural. Procedural user symbols are known only to the procedure in which they are generated. Interprocedural user symbols are known to the procedure in which they are generated and to any other procedure nested within that procedure.

User symbols can contain data that is up to 32 bits long. If the high–order bits are insignificant (that is, if they can be truncated without changing the value of the data), then the contents of the user symbol can be stored in a data field that is less than 32 bits long. Attempting to put a data value into a field that is too small for it will generate an error code of 189 (Appendix B).

PLC-3 Address Specifications

The PLC–3 processor uses logical addresses to reference data in memory. No PLC–3 address is valid unless it memory location has been allocated. You can allocate memory by using the CREATE command in PLC–3 programming. The PLC–3 Programming Manual (publication 1775–6.4.1) explains how to do this. Note that the CREATE command for memory allocation is different than the CREATE command for creating symbolic addresses (Chapter 6).

The following rules apply when specifying a PLC–3 logical address in a message procedure:

- **1.** Symbolic addresses must be defined to either the word level or the file level of specification.
- **2.** A word address may be either: (a) a block address specified to the word level, (b) a symbolic address of a word, or (c) a symbolic file address followed by a colon (:) and an offset.
- **3.** A size specification must be preceded by a word address and a comma (,).
- **4.** An offset specification must be preceded by a file address and a colon \therefore
- **5.** A bit number must be preceded by a word address and a slash (/).
- **6.** To access the pointer or floating point sections of memory, extended addressing must be used.

These rules are applied in the formats given below for addressing PLC–3 data locations.

Addressing a File

The format for addressing a PLC–3 file is one of the following:

<fileaddr> <flesym>

For assignment commands that copy data from one file to another, both the source and the destination file must be exactly the same size.

For PLC–3 timer and counter files, it is important to note that the data words are stored in the following order:

CTL PRE ACC

That is, the control, preset, and accumulated values for a given timer or counter are stored as consecutive words in the same file.

Addressing a Word Range

To address a range of words in PLC–3 memory, use one of the following formats:

<wordaddr><size> <filesym>:<offset>,<size> <wordsym>,<size>

Figure 4.5 is an example of addressing a range of PLC–3 words.

Figure 4.5 Example of Addressing a Range of PLC-3 Words

10022-I

You may use a word range only as the source field in an assignment command. The destination must be a file that is as large as, or larger than, the source range.

Addressing a Word

To address a single word in PLC–3 memory, use one of the following formats:

<wordaddr> <filesym>:<offset> <wordsym>

Note that <wordaddr>is interpreted as an octal value if the addressed word is in an input or output file. Otherwise, <wordaddr>is interpreted as a decimal value (section titled Number Systems).

To access words in the pointer of floating point sections of PLC–3 memory, use the PLC–3 extended addressing format. You can read about extended addressing in:

PLC–3 Programmable Controller Programming Manual (pub. no. 1775–6.4.1)

Addressing a Bit

To address a specific bit within a PLC–3 word, use one of the following formats:

<wordaddr>/

 <
 <
 <
 <

 <

 <

<t

 <t
 <t
 <t
 <t
 <t
 <t <filesym>:<offset>/
bit> <wordsym>/

bit>

Figure 4.6 gives some examples of addressing individual bits in PLC–3 memory.

PLC/PLC-2 Address **Specifications**

The PLC and PLC–2 processors use logical data addresses. These addresses are usually specified as octal numbers. However, the 1775–KA module interprets these addresses as decimal numbers unless they contain leading zeros (section titled Number Systems). Therefore, if you want to specify a PLC or PLC–2 word address as an octal number, begin the number with a 0(zero).

Addressing a Word Range

To address a range of words in PLC or PLC–2 memory, use this format:

<offset>,<size>

Figure 4.7 illustrates this addressing format

You may use a word range only as the source field in an assignment command.

Addressing a Word

To address a single word in PLC or PLC–2 memory, use this format:

<offset>

Addressing a Bit

To address an individual bit in PLC or PLC–2 memory, use this format:

<offset>/

 \leq

Figure 4.8 illustrates this addressing format.

Remote Station Address Specifications

To specify the address of data at a remote station, use the format shown in Figure 4.9. This format applies to both PLC–3 and non–PLC–3 remote stations. The characters #H delimit a remote Date Highway station, and the characters #M delimit a remote modem station.

Figure 4.9 Example of Addressing a Word in a Remote PLC-3 Station

Remote station addresses are subject to the following restrictions:

- **1.** A remote address can be used only with the single equals sign $(=)$ type of assignment command.
- **2.** In the assignment command, either the source or the destination, but not both, may be a remote address.
- **3.** A remote address may contain an embedded expression, but a remote address may not be embedded in an expression.

Expression

Expressions use operators to combine two or more numeric values into a single value. Table 4.B lists the operators that can be used in an expression. These operators are listed from highest priority (1) to lowest priority (10). Expressions may be nested within other expressions by enclosing the inner expression within parentheses.

Table 4.B Expression Operators

The result of an expression depends on the order in which the operators are executed. The order of execution depends on the type of operator and on left– to–right placement within the expression. Table 4.B gives the order of execution for the different operators. For example, the command

 $$B67:45 = 6+3*2$

would store the value 12 in word 45 of binary file 67. This is because multiplication is performed before addition.

If an expression contains several operators with the same order of execution, those operators will be executed in the left–to–right order in which they appear within the expression. Extra set of parentheses can be nested within each other to change the order of execution. In such cases, the expression within the inner–most set of parentheses is evaluated first. For example, the command

 $$B67:45=36\%((6+3)*2)$

would store the value 2 in word 45 of binary file 67 (% is the operator for division).

Expressions can be used anywhere that direct numeric values can be used within a message procedure, including within an address field. For example, in the statement

\$B67:(WORD+3)=5

the expression (WORD+3) specifies the address of a word within binary file 67. Note that the parentheses are necessary to indicate that $+3$ is part of the word address in this case.

Number Systems

Within an expression, direct values are always interpreted as decimal (base 10) numbers unless you indicate that they are octal (base 8). You can specify an octal value by starting the number with a leading zero. For example, 17 in an expression is interpreted as decimal 17, but 017 is interpreted as octal 17 (or decimal 15).

Operators

This section describes the operators listed in Table 4.B.

Bit Operator

The bit operator allows you to address a specific bit of a value stored under a user symbol. For example, the statement

\$I12:24/7=US_3/4

puts the value (0 or 1) of bit number 4 of user symbol US_3 into input file 12, word 24, bit 7.

The bit address itself can also be a user symbol or an expression. For example, in the statement

\$I12:24/7=US_3/(4+US_1)

the expression (4+US_1) specifies a particular bit within user symbol US_3 .

Note that the value appearing after the bit operator must be within the range of values allowed for bit addresses. Since user symbols are 32–bit values, a bit address for a user symbol must be in the range of 0 to 31 (decimal). Bit addresses for data table words must fall in the range of 0 to 15 (decimal).

Logical Operators

The logical operations are complement, AND, and OR. These operations are used to construct logically true or false conditions. They are generally used in decision statements such as the IF command (see section titled IF Command, chapter 6).

The result of a logical complement is 1 (true) if the expression following the .NOT. is a value of 0 (zero). Otherwise, the result is 0 (false). For example, consider the command

\$I12:24=.NOT.SYMBOL_A

If the value of SYMBOL_A is 0 (zero), then a 1 is stored in word 24 of input file 12. If the value of SYMBOL_A IS anything other than 0, then a 0 (zero) is stored in word 24 of input file 12.

The result of a logical AND is 1 (true) if the expression preceding the .AND. and the expression following the .AND. are both non–zero. Otherwise, the result is 0 (false).

The result of a logical OR is 1 (true) if either the expression preceding the .OR., the expression following the .OR., or both expressions are non–zero. Otherwise, the result is 0 (false).

Bitwise 32–Bit Operators

Bitwise 32–bit operators manipulate the individual bits in a 32–bit operand.

The bitwise 32–bit complement (.BNOT.) inverts the state of each bit in the 32– bit expression. That is, bits set to 1 are inverted to 0, and bits set to 0 are inverted to 1.

The bitwise 32–bit AND (.BAND.) forms a bit–by–bit logical AND of two 32–bit operands. There is no carry from one bit position to the next within the operand. For example, if

A contains the bit pattern 10101010010011110010101010101011

B contains the bit pattern 01110101011100100010101110001010

then the assignment C=A.BAND.B yields

C contains the bit pattern 00100000010000100010101010001010
The bitwise 32–bit EXCLUSIVE OR (.BXOR.) forms a bit–by–bit logical EXCLUSIVE OR of two 32–bit operands. There is not carry from one bit position to the next within the operand.

The bitwise 32–bit OR (.BOR.) forms the bit–by–bit logical OR of two 32–bit operands. There is not carry from one bit position to the next within the operand.

Arithmetic Operators

The arithmetic operations are addition, subtracting, multiplication, and division. These are binary (not BCD) operations that produce 32–bit signed integer results.

A result from these arithmetic operations should normally be assigned to a 32– bit destination. The result can be assigned to a 16–bit destination only if the result is small enough in absolute value (less than 65,535) to fit into 16 bits. If the result is assigned to a 16–bit destination but is too large to fit into 16 bits, then an error code of 215 results.

There is no indication of overflow or underflow conditions with arithmetic operations.

Shift Operators

When a left arithmetic shift $\left\langle \langle \cdot \rangle \right\rangle$ is executed, zeros are shifted into the rightmost bits of the expression. The leftmost bit are shifted out of the expression and are lost.

When a right arithmetic shift (\gg) is executed, the leftmost bit of the expression does not change. If the leftmost bits is a 1, then 1's are shifted in from the left. If the leftmost bit is a 0 (zero), then 0's are shifted in from the left. Since the leftmost bit of an expression is the sign bit, this means that the right arithmetic shift does not change the sign of a numeric value. The rightmost bits are shifted out of the expression and are lost.

Comparison Operators

Comparison operators result in a value of 1 if the comparison is true and 0 (zero) if the comparison is false. For example, consider the command

 $$I12:23 = ($CACC:1.GE.SCACC:2)$

If the accumulated value of counter 1 is greater than or equal to the accumulated value of counter 2, then the number 1 is stored in word 23 of input file 12. If the accumulated value of counter is less than the accumulated value of counter 2, a value of 0 (zero) is stored in word 23 of input file 12.

Resulting Values

The result of an expression is a 32–bit value. If the high–order bits are not significant (that is, if they can be truncated without changing the value of the expression), then the result can be stored in a data field that is less than 32 bits long. Attempting to put a value into a field that is too small for it results in an error code of 215 (Appendix B).

Editing

Table 5.A Example of Message instruction Editing

System Prompt	Action	Key Strokes
	Start edits.	SED [ENT]
	Insert rung.	IR [ENT]
	Enter the energize bit for the message rung. In this case, binary file 0, word 0, bit 0.	$-$] [$-$ B0:0/0 [ENT]
	Enter the message instruction.	MSG [ENT]
ENTER FILE ADDRESS	Enter the address of the file where the message instruction will reside in memory. In this case, binary file 1.	FB1 [ENT]
ENTER SYSTEM ADDRESS OR SYMBOL	Enter the channel designation for the 1775-KA module. In this case, 2 is the module status, 5 is the 1775-KA module type, and 1 is the thumbwheel number of the module.	E2.5.1 [ENT]
ENTER MESSAGE TYPE	Enter the message type. This is always 1 for the 1775-KA module.	1 [ENT]
	Enter a single 1775-KA assignment command or the name of a message procedure. In this case, the name of the message procedure is PROC 1.	@PROC-1 [ENT]
	End edits.	EE [ENT]

Before the 1775–KA module can transfer data to or from any file in PLC–3 memory, that file must exist and it must have enough memory allocated to it to accommodate the data transfer. You can create and allocate a file using the PLC–3 memory management commands. Refer to the PLC–3 Programming Manual (publication 1775–801) for a description of memory management. Allocating Memory

Table 5.B shows an example of editing a message procedure through an Industrial Terminal connected to a 1775–S4A module. Table 5.C shows how to edit the same message procedure through a data terminal connected to a 1775–S4B module. Editing Message Procedures

Table 5.C Example of Editing a Message Procedure Through a Data Terminal

Note that it is not always necessary to create a message procedure. If you want to execute just a single assignment command that is no more than 76 characters long, then you can enter that command as part of the ladder diagram message instruction (Table 5.A). If you want to execute more than one 1775–KA command, or if a single assignment command is more than 76 characters long, then you must create a message procedure to contain those commands.

Also note that every message procedure must end with either an EXIT or a STOP command. The EXIT command is normally preferred because the STOP command is a more extreme measure that results in error 179 (Appendix B).

Message Procedure Commands

General

The 1775–KA module has its own command language that you can use in programming message procedures. This chapter describes the available commands and gives some examples on how to use them. Table 6.A summarizes the commands.

Table 6.A Message Procedure Commands

Command	Format and Explanation	
Ξ (assignment)	<destination>3=<source/></destination>	
	Assign a numeric value to a user symbol or copy data from the source to the destination	
CREATE	C@ <system symbol=""><logical address=""> Create a symbolic address and equate it to a logical address.</logical></system>	
DELETE	D @ <system symbol=""> Delete a symbolic address or an entire message procedure from PLC-3 memory.</system>	
(execute)	@ <system symbol=""> Execute the named message procedure.</system>	
EXIT	E Terminate execution of the current message procedure.	
GOTO	G <label> Continue executing the current procedure from the point specified by the label.</label>	
IF	I <expression><embedded command=""> Execute the embedded command only if the specified expression is true.</embedded></expression>	
ON ERROR	O <embedded command="">3 Execute the embedded command only if an error occurs after this statement in the procedure.</embedded>	
STOP	S Terminate execution of the message (MSG) instruction in the PLC-3 ladder diagram program.	

Each command can be abbreviated to the letters shown in the format column of Table 6.A. In general, it is best to abbreviate a command to the shortest possible form. This not only makes the commands easier to program, but it also saves memory space and reduces execution time.

Blanks may be inserted anywhere to improve the readability of a message procedure. However, blanks should be kept to a minimum because they use memory space and slow execution of the message procedure.

The assignment command is the most fundamental yet versatile of all the commands. Its primary purpose is to copy data from the source location to the destination location. Table 6.B lists the various types of sources and destinations. Any type of source in Table 6.B may be used with any type of destination listed. Assignment Command

Of special interest is the case where a user symbol is the destination of the assignment. In such a case, if the user symbol was not previously defined in the message procedure, a new symbol is generated. If the symbol has already been defined, using it again as a destination causes its value to be changed to the value given it by the latest assignment command.

Note that you can not transfer data from another station and place it into a user symbol defined at your local PLC–3.

Format

The equals sign $(=)$ is the assignment command. As Table 6.A shows, the destination for the assignment is on the left of the equal sign, and the source or the numeric value is on the right. In all cases, the source value is assigned to (or copied to) the destination location. Thus, the assignment is from right to left on the command line. For example, the statement

\$I12:024–US_5

copies the value of user symbol US_5 into word 24 (octal) of input file 12.

Modifiers

Several modifiers may be added to the basic assignment command. These modifiers affect three aspects of the assignment:

- Scope of assignment
- Priority level of Data Highway message
- Type of command message transmitted

Scope of Assignment

A double equals sign (==) can also be used for the assignment command. The extra equals sign modifies the scope of an assignment involving a user symbol. If the destination of the assignment is a user symbol, the double equals sign defines the destination to be an interprocedural user symbol. With the single equals sign, the destination becomes a procedural user symbol.

For example, the statement

 $US_2 == 6$

defines US_2 to be an interprocedural user symbol and assigns to it the value 6.

Do not use the double equals sign $(==)$ with anything other than a user symbol as the destination.

Message Priority

Data Highway messages may be either one of the following priority levels:

- Normal
- **Priority**

If you use the less–than sign $\left(\leq \right)$ with the assignment command, the command will generate a priority Data Highway message. Without the less–than sign, the assignment command will generate a normal Data Highway message.

For example, the statement

#H027\$I15:4<=\$I12:24

transmits a priority message to Data Highway station 27 (octal).

The priority modifier can be used with either type of assignment $(=o r==).$

Important: Stations with high priority messages are given priority over stations with normal priority messages throughout the command/reply cycle. For this reason, a command should be given a high priority designation only when special handling of specific data is required. Using an excessive number of high priority commands defeats the purpose of this feature and could delay or inhibit the transmission of normal priority messages.

Command Message Type

Command messages are of two types:

- **protected**
- unprotected

As explained in section titled Data Transfers, (chapter 3), protected commands can access only specified areas of data table memory at a PLC/PLC–2 station. You will need to send a protected write command only if a switch at the remote PLC/PLC–2 prohibits other stations from sending unprotected write commands. Unprotected commands can access any area of the data table.

By default, command messages generated by the assignment command in PLC–3 message procedures are of the protected type. To generate an unprotected command message, use a blank space and the modifier U after the assignment command.

the value of the expression is false (0), the embedded command is not executed. The embedded command may be any of the available commands except another IF or an ON_ERROR.

Figure 6.1 demonstrates the combination of a label, a GOTO command, and an IF command to construct a simple loop that assigns the integers 0 through 7 to successive words in binary file 50.

ON_ERROR Command

The ON_ERROR command specifies what action should be taken if an error is encountered during execution of the message procedure. The ON_ERROR command is not executed sequentially in the procedure; it is executed only when an error occurs.

Table 6.A illustrates the format of the ON_ERROR command. The ON_ERROR command contains an embedded command that is executed when an error occurs.

The ON_ERROR command applies to all other commands between itself and the next ON_ERROR command. For example, consider the following sequence:

There are two functions:

- TO_BCD
- FROM_BCD

Figure 6.2 illustrates the format of these functions as they might appear in an assignment command.

Figure 6.2 Examples of TO-BCD and FROM-BCD Functions

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The parameter of the function must be enclosed in parentheses. The parameter may be any one of the following:

- A direct numeric value (either decimal or octal)
- An expression
- A user symbol
- A logical address
- A symbolic address

TO_BCD Function

The TO_BCD function converts its parameter into a binary coded decimal value that is 32 bits long. For example, the TO_BCD function in Figure 6.2 stores the number 27 in binary–coded–decimal format in word 12 of the decimal section of PLC–3 memory. After this function is executed, word 12 will contain the following bit pattern:

0000 0000 0010 0111

FROM_BCD Function

The FROM_BCD function converts its parameter from binary–coded–decimal format to binary format. The resulting value is 32 bits long. For example, the FROM_BCD function in Figure 6.2 converts the contents of decimal word 12 from binary coded decimal to a regular decimal value of 27. From the above example (section titled TO_BCD Function), the FROM_BCD function stores the following bit pattern in user symbol COUNT:

0000 0000 0000 0000 0000 0000 0001 1011

As you can see from these examples, TO_BCD and FROM_BCD perform opposite functions.

You can add your own explanatory comments to any command line in a message procedure. To do this, enter a semicolon (;) after the command. Then enter your comment after the semicolon. Figure 6.3 illustrates the format for comments. **Comments**

Anything that appears between a semicolon and the end of the command line is considered to be a comment. Comments may be any length. The end of the command line, and therefore the end of your comment, is delimited by the carriage– return and line–feed pair of characters.

A comment can be the only thing on a line. Do not use comments on the same line as label. Doing so will cause errors in the message procedure.

Chapter

Error Reporting

Figure 7.1 illustrates how the error block works. In this figure, an addressing error (invalid destination address) occurs in procedure SUB2, which is nested 3 levels deep. Word 5 of the error block gives the line number where the error occurred in procedure SUB2. Word 4 gives the number of the line in procedure SUB1 that executed procedure SUB2.

And word 3 gives the number of the line in procedure MAIN that executed procedure SUB1.

Figure 7.1 Examples of Error Block Operation

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Note that an ON_ERROR or an IF command may contain an embedded command to execute another procedure. In these cases, the embedded execute command is treated just like a nesting level. Figure 7.2 illustrates this point for an ON_ERROR command. In this figure, an addressing error in line 10 of procedure MAIN causes activation of the ON_ERROR command, which calls for execution of procedure SAM. But SAM also contains an error. The error in SAM is the last one detected, so it is the one finally reported in the error block. Since procedure SAM is called by the ON_ERROR command in procedure MAIN, the nesting for SAM is 2 levels deep.

Figure 7.2 Examples of ON_ERROR Nesting

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Access to Error Block

The error block retains its data even after the message procedures are done executing. It is re–initialized with each execution of a MSG instruction in the PLC–3 ladder diagram program.

The extended address for the beginning of the error block file is

 $\text{SE2.5.m.4.0}.$

where "nn" is the thumbwheel number of the 1775–KA module. You can access this error block by any one of the following means:

- Displaying it through the front panel of the PLC–3 controller
- Using the data monitor mode of the Industrial Terminal (cat. no. 1770–T4)
- Using the move status (MVS) command in the PLC–3 ladder diagram program
- Using the I/O Scanner–Message Handling Module (cat. no. 1775–S4B)
- Using the 1775–KA module
- Using the 1775–GA

Chapter **6**

General	This chapter presents some detailed examples of 1775–KA module commands and message procedures.
Individual Commands	The first set of examples shows individual commands that could be programmed directly into a PLC–3 message (MSG) instruction. Figure 8.1 illustrates the differences in reading and writing data between two PLC–3 stations. Figure 8.2 shows how to write different types of data to a remote PLC–3 station. Figure 8.3 shows how to write different types of data to a remote PLC or PLC–2 station.

Figure 8.1 Reading and Writing PLC-3 Data Word

Reading word 5 of binary file 3 at data highway station 045 into word 17 of integer file 4.

Figure 8.2 Writing Data to a Remote PLC-3 Station

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Message Procedure

Figure 8.4 presents a printed listing of a Data Highway message procedure. As the listing indicates, the purpose of the procedure is to monitor the state of a status bit in a remote Data Highway station.

Figure 8.4 Example Data Highway Message Procedure

```
; PROCEDURE –– @REM_TURNON
; This procedure will monitor the state of a bit in a remote
; station and, when that bit goes true, turn on a bit
; locally for either 300 seconds or until the remote bit
; goes false.
;
   ON_ERROR @LOG_ERROR
   A = = 0 \lambda = = 0
   CREATE @TIM–START $B0:0
   CREATE @TIM_CTL $TCTL:1
   CREATE @TIM_PRE $TPRE:1
   T\_ON\_BIT = 0T_DONE_BIT = 017
   CREATE @PROCESS $N3:7
   P ON BIT = 5
   ON = 1OFF = 0LOOP 1:
         B0:0/1 = $H023$B5:3/2 IF ($B0:0/1 .EQ. OF) GOTO LOOP1
   @TIM PRE = 300
   @TIM_START/T_ON_BIT = ON
   @PROCESS/P_ON_BIT = ON
   LOOP 2:
         $B0:0/1 = $H023$B:3/2 IF (($B0:0/1 .EQ. ON) .AND. (@TIM_CTL/T_DONE_BIT .EQ. OFF)) GOTO LOOP2
   @PROCESS/P_ON_BIT = OFF
   EXIT
                                            ;initialize error pointer
                                             ;timer start word
                                             ;timer control word
                                            ;timer preset word
                                            ;timer on bit
                                            ;timer done bit
                                             ;process word
                                             ;process on bit
                                            ;check remote bit in loop
                                            ;fetch and save remote bit
                                             ;set timer for 300 sec
                                             ;turn timer on
                                             ;turn process on
                                             ;check timer and remote bit in loop
                                             ;fetch and save remote bit
; PROCEDURE –– @LOG_ERROR
;This procedure will fetch the error block out of the
; Module Status Area and record it along with the time
; of day in status file 5.
;
    CREATE @STATUS $S5
    CREATE @ERR_BLK $E2.5.1.4.0
 CREATE @TOD $S1:3
@STATUS: (A) = @ERR_BLK, 6@STATUS: (A + 6) = @TOD, 2 IF ((ERROR .GE. 81) .AND. (ERROR .LE. 92)) GOTO NO_STN :no station – fatal error
    IF (A .GE. 72) GOTO TIMEOUT
   A = = A + 8 EXIT
                                            ;copy error block (6 words)
                                            ;copy time of day (hrs, mins)
                                            iafter ten errors, tell operator
;
    NO_STN:
        SS4:3/5 = 1 STOP
;
   TIMEOUT:
        $54:3/4 = 1 EXIT
                                             ;energize 1775–S4B report generation rung
                                             ;exit procedure with an error
                                            ;energize 1775–S4B report generation rung
                                             ;return to @PREM_TURNON
```
Some of the statements in the sample procedure are not necessary to accomplish the bit monitoring. However, they were included to illustrate more of the functions and programming techniques available with the 1775–KA module.

Note that the 300 second timer used in this example is not an accurate, real– time clock. This is because the time between successive executions of the bit/timer check depends on Data Highway activity and on the activity of the local PLC–3 processor. For example, if the 300 second timer times out immediately after its done bit is checked, the 1775–KA module will not detect this condition until its next pass through LOOP2. If the Data Highway is busy with other activity, it will take a while for LOOP2 to check the remote bit. PLC–3 ladder diagram programming provides better timer updates and responses.

The example procedure also assumes that the referenced memory areas have been created. Specifically:

- **1.** Status file S5 must be big enough to hold a reasonable number of timeout errors (error #37).
- **2.** Timer T1 is a one–second timebase timer. Bit B0:0/0 controls the ladder diagram rung that activates the timer. Figure 8.4 refers to this bit as TIM_START/T_ON–BIT.
- **3.** Bit S4:3/4 activates a message instruction that executes a report generation procedure. In this way, the 1775–KA module can indirectly cause execution of a report generation procedure to display a message on the operator's terminal.

Computer to PC Communication

Introduction to Layered Communication

This chapter and the chapters that follow (10,11, and 12) described how to write a software driver that enables your computer to communicate through the RS–232–C port of the PLC–3 Communication Adapter Module. Therefore, you do not need to read these chapters if you are only using PC's. The interface modules contain software drivers for PC to PC communication.

Chapter

In this chapter and the chapters that follow (10,11, and 12) we describe a layered approach to writing a software driver for your computer. According to the standard for network architecture developed by the International Standards Organization (ISO), communication networks should be divided into layers. Each layer performs specific functions. By separating the communication network into independent layers, it is easier to make changes to one of the network's functions without having to redesign the entire network. Ideally, the layers of a network should be as independent of one another and interact with one another in the same way as the organs of the human body. Because the organs of the human body are independent of one another, it's possible for a surgeon to operate on the lungs or heart without losing the life of the patient. Yet at the same time the organs of the body interact when we run or walk or type on a word processor.

You should use a layered approach to developing communication software for your computer. You don't have to design your communication software in this layered fashion, but your software must perform all the functions described for the layers in this manual. In most cases, it will be easier for you to implement and debug the communication software if you follow this layered approach.

The Data Highway uses these four layers of the ISO model for communication between stations:

- application layer provides the Data Highway commands that you use to transfer data and manage the network.
- network layer determines how you address a Data Highway command. It also provides less visible functions, such as controlling the flow of information, establishing a path between stations, and routing messages from your station to another station.
- \blacksquare data link checks the path between stations for errors to ensure that data is transmitted in a proper sequence, frames messages sent by a station and checks the integrity of messages received by a station. This layer is not visible to the person placing Data Highway commands in a program.
- physical link sets up, maintains, and disconnects a physical link between two stations. This layer consists largely of hardware (Data Highway modules and cable). Like the data link layer, this layer is not visible to the person placing Data Highway commands in a program.

A Data Highway command consists of many fields, each of which originates from one of the above layers at the sending station. When a station receives a Data Highway command, it separates these fields so that a single layer uses only those fields it needs to perform its specific function.

The application layer uses these fields of a Data Highway command:

- a command (CMD) field and function (FNC) field identify the type of command that is sent.
- a status field (STS) contains a code that indicates if the command was successfully sent from one station to another.
- an address (ADDR) field specifies the address in the remote station's memory
- a data (DATA) field contains the data that is sent from one station to the other.

These bytes and fields are discussed in greater detail later. For now, notice (Figure 9.1) that the FNC, ADDR, and DATA fields from the application layer are treated as data by the network layer. We say that the data is framed by the fields of the network layer.

Legend: $x =$ low hex digit of CMD byte supplied by application layer

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The network layer uses these fields of a Data Highway command:

- \bullet the destination (DST) and source (SRC) byte specify the address of the station that is receiving the command and the address of the station that is sending the command.
- a command (CMD) byte at the network layer indicates whether the message is a command or reply.
- a status (STS) byte at the network layer indicates whether the message was successfully executed by the sending station.
- a transaction (TNS) field identifies the particular command and reply cycle the message belongs to.

Again, these bytes and fields are described in much greater detail later in this manual. Notice (Figure 9.2) that the network layer (DST, CMD, STS, TNS,and the application layer) is treated as data by the data link layer. We say that this data is framed by the data link layer.

The data link layer uses these fields of a Data Highway command:

- start of text (STX)
- \blacksquare end of text (ETX)
- \blacksquare enquiry (ENQ)
- acknowledge (ACK)
- data link escape (DLE)
- negative acknowledge (NAK)
- start of header (SOH)
- end of transmission (EOT)
- **block check character (BCC)**

These control characters are described in greater detail later in this manual. The type of **protocol** you choose for the data link layer determines the meaning of these control characters.

Full Duplex Link Packet and Half–Duplex Slave Packet

Half–Duplex Master Message Link Packet

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Full-Duplex vs Half-Duplex Protocol for the Data Link Layer

We use the term protocol to describe the relationship between two similar layers at two different stations. The protocol could, for example, be the relationship between the data link layer at station A and the data link layer at station B.

When you write a software driver for your computer, you may implement either a data link layer that uses a half–duplex protocol or a data link layer that uses a full duplex protocol. You select the half–duplex or full–duplex protocol with the LIST function (section titled Modem Port, chapter 2).

In general, the full–duplex protocol provides higher data throughput, but you can use it only for communication between two peer stations.

Half–duplex protocol is for one master and one or more slaves. Half–duplex protocol provides lower data throughput but is easier to implement than the full–duplex protocol. You should use half–duplex protocol if:

- You are using multidrop baseband MODEMS to connect multiple slave stations to a single master computer. (You must use MODEMS for this type of link unless there is only one slave).
- You are using MODEMS that have only half–duplex capability
- You are willing to sacrifice data throughput in exchange for ease of implementation

Half–duplex protocol does not allow embedded responses. The 1775–KA module has slave mode capability only; you must provide the master function from your computer.

Chapter 10

Full-Duplex Protocol

In general, full–duplex protocol gives higher data throughput, but it can handle communication between only two peer stations. Half–duplex protocol provides master–slave polling capability and can handle communication with as many as 255 slave stations, but it gives lower data throughput.

This chapter describes the data link layer for full–duplex protocol. Chapter10 describes the data link layer for half–duplex protocol. Chapter 12 describes the network layer and the application layer for both protocols.

The full–duplex protocol resembles ANSI X3.28–1976 specification, combining features of subcategories D1 (data transparency) and F1 (two–way simultaneous transmission with embedded responses). Full-Duplex Protocol

You can use full–duplex protocol for a point–to–point link or a multidrop broadband MODEM link that allows two–way simultaneous transmission. It is more difficult to implement than half–duplex because it requires you to use interrupts and multi–tasking programming techniques. It is intended for high– performance applications where you need to get the highest possible throughput from the available communication medium.

At the 1775–KA module, select the unpolled mode with the LIST function (chapter 2).

Transmission Codes

Full–duplex protocol is a character oriented protocol that uses the ASCII control characters extended to eight bits by adding a zero for bit 7. See ANSI X3.4, CCITT V.3, or ISO 646 for the standard definition of these characters.

The particular ASCII control characters used are listed below.

Additionally, a block check character (BCC) is used at the end of each packet for error checking. These bytes can be any value from 00 to FF hex.

In the following paragraphs we use the term code to mean an indivisible sequence of ASCII characters or values having specific meaning to the protocol. Indivisible means that the component bytes of a code must be sent one after another with no other bytes between them. It does not refer to the timing of the bytes.

Full–duplex protocol uses these codes:

Control codes:

- DLE STX
- DLE ETX BCC
- DLE ACK
- DLE NAK
- DLE ENQ

Link–layer data codes:

- Data (single bytes having values 00–0F and 10–FF hex)
- DLE DLE (to represent the value 10 hex)

We can also group codes into two classes according to their use:

- **1.** Message codes issued from a station sending a message.
- **2.** Response codes issued from a station receiving a message.

The full–duplex codes sent by the station transmitting a message are:

- DLE STX indicates the start of a message packet.
- Link–layer data (00–OF and 10–FF hex) encodes the bytes of the network packet.
- \blacksquare DLE DLE encodes the value 10 hex in the network packet. This is necessary to distinguish a text code of 10 hex from a DLE control code of 10 hex.
- DLE ETX BCC terminates a message packet.
- DLE ENQ requests the retransmission of the last received transmission.

The full–duplex response codes sent by a station receiving a message are:

- \blacksquare DLE ACK signals that the receiver has successfully received the last message sent.
- DLE NAK signals that the receiver did not successfully receive the last message sent.

Link-Layer Message Packets

A link–layer message packet starts with a DLE STX, ends with a DLE ETX BCC, and includes all link–layer data codes in between. Data codes can occur only inside a message packet.

Response codes occur inside a message packet. If you select the embedded responses option with LIST (chapter 2) the response codes can also occur between a DLE STX and a DLE ETX BCC, but these response codes are not part of the message packet: they are referred to as embedded responses.

Figure 10.1 shows the format of a link–layer message packet for full–duplex protocol, and the layer at which each portion should be implemented. At the end of each message packet is the one–byte BCC field.

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Figure 10.1 Link Packet Format for Full-Duplex Protocol
Block Check

The block check character (BCC) is a means of checking the accuracy of each message packet transmission. It is the 2's complement of the 8–bit sum (modulo–256 arithmetic sum) of all data bytes between the DLE STX and the DLE ETX BCC. It does not include any other message packet codes or response codes.

For example,if message packet contained the data codes 8, 9, 6, 0, 2, 4, and 3, the message packet codes would be (in hex):

The sum of the data bytes in this message packet is 20 hex. The BCC is the 2's complement of this sum, or EO hex. This is shown in the following binary calculation:

> 0010 0000 20 hex 11011011 1s compliment +1 1110 0000 2s compliment (E0 hex)

To transmit the data value 10 hex, you must use the data code DLE DLE. However, only one of these DLE data bytes is included in the BCC sum. For example, to transmit the values 8, 9, 6, 0, 10, 4, and 3 hex, you would use the following message codes:

Represents single data byte value of 10

DLE STX Data DLE ETX BCC

In this case, the sum of the data bytes is 2E hex because only one DLE text code is included in the BCC. So the BCC is D2 hex.

The BCC algorithm provides a medium level of data security. It cannot detect transposition of bytes during transmission of a packet. It also cannot detect the insertion or deletion of data values of zero within a packet.

Two-Way Simultaneous Operation

On a two–way simultaneous link, two physical circuits connect four distinct and independent software routines. Figure 10.2 shows these software routines as transmitters (XMTR) A and B and receivers (RCVR) A and B.

Figure 10.2 Data Paths for Two-Way Simultaneous Operation

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There are also four independent data paths involved. Paths 1 and 3 carry message codes between A and B; paths 2 and 4 carry response codes between A and B.

A software multiplexer combines those message codes and response codes going in the same direction. At the other end of the link, a software separator separates those message codes from the response codes. Internal software directs the message codes to the receiver and the response codes to the transmitter. On each physical circuit, you can intermingle response codes from a receiver to a transmitter with message codes sent from a transmitter to a receiver (unless you do not choose the embedded response option in LIST). Figure 10.3 shows this implementation.

Figure 10.4 shows path 1 with unrelated parts of Figure 10.3 removed.

We could show paths 2, 3, and 4 in a similar way.

The full–duplex protocol is symmetrical; that is, anything that we can say about transmitter A, receiver B, and paths 1 and 2 applies equally to transmitter B, receiver A, and paths 3 and 4. There are actually two independent instances of the protocol operating simultaneously. For simplicity, we define the link protocol on the subsystem that carries messages from A to B, with reference to figure 10.5

Figure 10.5 Message Transmission from A to B

Although the protocols on each subsystem operate independently, there is a slight delay when you transmit a response code in the middle of a stream of message codes. Also, any non–transient hardware problem that affects message codes traveling over a hardware circuit affects response codes on the same circuit.

Message Characteristics

In the network layer (chapter 12) the message source provides the transmitter with the messages it sends. The message sink tells the receiver what to do with the messages it receives.

Upon request from the transmitter, the message source supplies one network packet at a time. It must be notified about the success or failure of the transfer to the receiver before supplying the next message. When the message source is empty, the transmitter waits in an inactive state until a message is available.

Whenever the receiver has received a link packet successfully, it attempts to give the network packet portion (link level data) to the message sink. If the message sink is full, it must notify the receiver.

Figure 10.6 represents the protocol environment.

Figure 10.6 Protocol Environment

Full–duplex protocol places the following restrictions on the network packet that is submitted to the link layer for transfer:

- The size of a valid network packet is 6 bytes minimum, and 250 bytes maximum.
- The first byte of a network packet must be the station number of the receiver station. The receiver ignores messages that do not contain the correct station number.
- As part of the duplicate message detection algorithm, the receiver checks the second, third, fifth, and sixth bytes of each network packet. (Recall that a network packet consists of the source, command, and transaction fields.) At least one of these bytes of the current network packet must differ from the corresponding byte of the previous network packet in order for the receiver to accept the current network packet. Otherwise, the receiver assumes that the current packet is a retransmission of the previous packet, so it discards the current packet.

Transmitter Actions

Whenever the message source can supply a packet and the transmitter is not busy, transmitter A sends a link packet on path 1. It then starts a

timeout, and waits for a response on path 2. You can use the diagnostic set timeout command to set this timeout period for the 1775–KA module. The default setting is 3 seconds.

When transmitter A gets a DLE ACK, the message transfer is complete. After signaling the message source that the message has been sent successfully, transmitter A proceeds with the next message.

If transmitter A gets a DLE NAK, it retransmits the same message. The transmitter restarts the timeout and waits again for a response. By using the diagnostic set NAKs command, you can specify how many times the 1775–KA module will attempt to retransmit a given message. The default setting is 3. Once the number of retransmissions exceeds this limit, the transmitter should notify the message source that the transmission has failed. The transmitter can then proceed with the next message.

If the timeout expires before transmitter A gets a response, it sends a DLE ENQ on path 1 to request a retransmission of the last response sent on path 2. Transmitter A restarts the timeout and waits for a response. By using the diagnostic set ENQs command, you can specify how many timeout periods the 1775– KA module will allow per message it transmits. The default setting is 10. If this ENQ limit is exceeded, the transmitter notifies the message source that the transmission has failed. The transmitter can then proceed with the next message.

DLE ACK and DLE NAK are the only response codes defined. If the receiver gets an invalid response code, it ignores it.

Note that the transmitter must encode a text value of 10 hex as two consecutive (indivisible) bytes, each of value 10 hex. This is necessary to distinguish the text value of 10 hex from the DLE control code of 10 hex. This technique is known as DLE stuffing. The receiver must be able to reverse this process and extract the original text value of 10 hex.

Figure 10.7 is a flowchart which gives a simplified view of an example of software logic for implementing the transmitter. Table 6.A gives a detailed description of an example of software logic for implementing the transmitter in structured English procedures. In appendix D are flowcharts which give a detailed view of an example of software logic for implementing the transmitter.

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Table 10.A Transmitter for Full-Duplex Protocol

TRANSMITTER is defined as
loop
Message=GET-MESSAGE-TO-SEND
Status=TRANSFER (Message)
SIGNAL-RESULTS (Status)
end
TRANSFER (Message) is defined as
initialize nak-limit and enq-limit
SEND (Message)
start timeout
loop
WAIT for response on path 2 or timeout.
if received DLE ACK then return SUCCESS
else if received DLE NAK then
begin
if nak-limit is exceeded then return FAILURE
else
begin count NAK retries;
SEND-MESSAGE (message);
start timeout
end
end
else if timeout
begin
if enq-limit is exceeded then return FAILURE
else
begin
count ENQ retries;
send DLE ENQ on path 1;
start timeout
end
end
end loop
SEND (Message) is defined as
begin
$BCC = O$
send DLE STX on path 1
for every byte in the message do
begin
add the byte to the BCC;
send the corresponding data code on path 1
end send DLE ETX BCC on path 1
end
GET-MESSAGE-TO-SEND
This is an operating-system-dependent interface routine that waits and allows the rest of the system to run until the message source
has supplied a message to be sent.
SIGNAL-RESULTS
This is an implementation-dependent routine that tells the message source of the results of the attempted message transfer.
WAIT
This is an operating-system-dependent routine that waits for any of several events to occur while allowing other parts of the system
to run.

Receiver Actions

Since the receiver gets "dirty" input from the physical world, it is more complex and must be capable of responding to many adverse situations. Some of the things that can conceivably happen are listed here:

- The message sink can be full, leaving the receiver with nowhere to put a message.
- A message can contain a parity error.
- The BCC can be invalid.
- The DLE STX or DLE ETX BCC may be missing.
- The message can be too long or too short.
- A spurious control or text code can occur outside a message.
- A spurious control code can occur inside a message.
- Any combination of the above can occur.
- The DLE ACK response can be lost, causing the transmitter to send a duplicate copy of a message that has already passed to the message sink.

Receiver B must keep a record of the last response code (DLE ACK or DLE NAK) sent on path 2 (Figure 10.5). If it receives a DLE ENQ, the receiver sends this recorded response code again.

The receiver also keeps a record of the first six link–level data bytes of the last message received. If the SRC, CMD, and both TSN bytes of a new message are identical to the corresponding bytes of this record, the receiver responds with a DLE ACK but ignores the new message. This process is known as duplicate message detection, and is part of the link–level data security. It guards against re–execution of a message that has already been received successfully, but for which the response code (DLE ACK) has been lost.

Until it receives a DLE STX or a DLE ENQ, the receiver ignores all input from path 1 except to set the last response variable to NAK. With the last response variable set to NAK, the receiver responds with DLE NAK to a DLE ENQ input. Otherwise, the receiver responds to a DLE ENQ input by sending it last response on path 2 and continues waiting for input. If the receiver gets a DLE STX, it resets its BCC accumulator and data buffer to zero and starts storing the link– level data in the data buffer so that it can later pass the link–level data to the network layer.

While the receiver stores all link–level data codes in the data buffer, it adds the link–level data code values to the BCC. If the data buffer

overflows, the receiver continues summing the BCC, but it discards the data.

The receiver also sets an error flag to indicate the occurrence of a parity, buffer overrun, message framing,or modem handshaking error. If the receiver receives any control code other than DLE ETX during this time, it aborts the message and sends a DLE NAK on path 2. When the receiver gets a DLE ETX BCC, it checks the error flag, the BCC, the message size, and the destination station number. If any of the tests fail, the receiver sends a DLE NAK on path 2.

If the current message packet passes the above tests, the receiver next begins the duplicate message detection process. In this process, the receiver compares the SRC, CMD, and both TNS bytes of the current message with the corresponding bytes of the previous message received. If these bytes are the same, the receiver discards the current message and sends a DLE ACK.

If the current message differs from the previous one, the receiver next tests the state of the message sink. If the message sink is full, the receiver sends a DLE NAK. Otherwise, the receiver:

- forwards the current link–level data to the message sink
- keeps a copy of the first six bytes of the current link–level data for purposes of duplicate message detection
- sends a DLE ACK

Figure 10.8 is a flowchart which gives a simplified view of an example of software logic for implementing the receiver. Table 10.B gives a detailed description of an example of software logic for implementing the receiver in structured English procedures.

In appendix D are flowcharts which give a detailed view of an example of software logic for implementing the transmitter.

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Table 10.B Receiver for Full-Duplex Protocol

RECEIVER is defined as
variables
LAST-HEADER is 4 bytes copied out of the last good message
RESPONSE is the value of the last ACK or NAK sent
BCC is an 8-bit block check accumulator
$LAST-HEADER = invalid$
LAST RESPONSE = NAK
loop
reset parity error flag
GET-CODE
if DLE STX then
begin
$BCC=0$
GET-CODE
while it is a data code
begin
if buffer is not overflowed put data in buffer
GET-CODE
end
if the control code is not a DLE ETX then send DLE NAK
else if error flag is set then send DLE NAK
else if BCC is not zero then send DLE NAK
else if message is too small then send DLE NAK
else if message is too large then send DLE NAK
else if header is same as last message send a DLE ACK
else if message sink is full send DLE NAK
else
begin
send message to message sink
send a DLE ACK
save last header
end
end
else if DLE ENQ then send LAST-RESPONSE
else LAST-RESPONSE = NAK
end
GET-CODE is defined as
loop
variable
GET-CHAR
if char is not a DLE
begin
add char to BCC
return the char and data flag
end
else

Full-Duplex Protocol Diagrams

The following figures show some events that can occur on the various interfaces. Control characters are shown in bold type. Link–level data is represented by xxxx. Line noise is represented by ???. BCC is shown at the end of each message packet. Time is represented as increasing from the top of the figure to the bottom. Figure 10.9 shows normal message transfer.

Figure 10.9 Normal Message Transfer

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Figure 10.10 shows a DLE NAK response to the initial message transmission. After the message is retransmitted, a DLE ACK response is given.

Figure 10.11 shows the transmitting station sending a DLE ENQ sequence after a timeout because it did not receive the initial DLE ACK response.

In Figure 10.12, retransmission occurs when noise hits both sides of the line. This type of noise destroys the DLE ACK while also producing invalid characters at the receiver. The result is that the receiver changes its last response to NAK and the transmitter retransmits the original message packet.

Figure 10.12 Message Transfer with Retransmission

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Figure 10.13 shows a DLE NAK response to the initial message transmission because the message sink is full. After the message sink is no longer full, a retransmission of the message causes a DLE ACK response.

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If you were to connect a line monitor to the wires between station A and B, and only the A to B subsystem were active, you could observe the following:

Examples

Normal Message Path 1: DLE STX xxxDLE ETX BCC BLE STXxxxxDLE ETX BCC Path 2: DLE ACK DER ROUGH DER ACK **Message with parity or BCC error and recovery** Path 1: DLE STXxx???xxDLE ETX BCC BLE STXxxxxDLE ETX BCC Path 2: DLE NAK DER NAK DER DER ACK **Message with ETX destroyed** Path 1: DLE STXxxxxx????[timeout] DLE ENQ DLE STXxxxxDLE ETX BCC Path 2: DLE NAK DER NAK DER DER ACK **Good message but ACK destroyed** Path 1: DLE STXxxxDLE ETX BCC [timeout] DLE ENQ DLE STXxxx etc. Path 2: DL???CK DLE ACK **Messages being sent in both directions** Path 1: DLE STXxxxDLE ETX BCC DLE STXxxxxDLE ETX BCC DLE STX Path 2: Path 2: DLE ACK DLE ACK DLE ACK Path 3: DLE STXxxx xxxxDLE ETX BCC DLE STX Path 4: DLE ACK Combined – Circuit AB:DLE STXxxxDLE ETX BCC DLE STXxxxxDLE ETX BCC DLE ACK DLE STX Circuit BA: DLE STXxxxDLE ACKxxxxDLE ETX BCC BLOCK DUE ACK DLE STX embedded response ACK on AB delayed slightly because ETX BCC are indivisible

Embedded Response Option

To simplify the design of the receiver in some cases, you can disable transmission of embedded responses by turning off the embedded response switch. If you turn this switch off, the 1775–KA module's multiplexer cannot embed response codes while sending a message. Instead,it delays sending response codes until after it sends the next DLE ETX BCC sequence.

Chapter

Half-Duplex Protocol

You must program a computer to serve as a master that controls which station has access to the link. All other stations are slaves and must wait for permission from the master before transmitting. Each slave station has a unique station number from 0 to 376 octal. The number 377 is a broadcast address. When the master sends a message addressed to 377, all slaves receive it.

The master can send and receive messages to and from each station on the multidrop link. If the master is programmed to relay messages, then slave stations on the multidrop link can engage in peer–to–peer communication.

Your multidrop link may be either a two–circuit system (master sends and slaves receive on one circuit, slaves send and master receives on the other), or a one–circuit system (master and slaves send and receive on the same circuit).

You may use a half–duplex, dial–up modem to connect the 1775–KA module to the multidrop link. The modem must signal data–carrier–detect at least once every 8 seconds. If it does not, the module will hang up. On a dedicated line, you can jumper lines 6,8, and 11 at the 1775–KA module to prevent the module from hanging up.

You cannot use multiple masters unless one master is limited to acting as a backup to the other, and does not communicate until the primary is shut down.

Half–duplex protocol is a character oriented protocol that uses the following ASCII control characters: Transmission Codes

These ASCII control characters are extended to 8 bits by adding a zero for bit 7. See ANSI X3.4, CCITT V.3, or ISO 646 for the standard definition of these characters.

Additionally, a block check character (BCC) is used at the end of each transmission packet for error checking. This byte can be any value from 00 to FF hex.

The term code means (in the following paragraphs) an indivisible sequence of one or more bytes having a specific meaning to the protocol. Indivisible means that the component bytes of a code must be sent one after another with no other bytes inserted between them. It does not refer to the timing of the bytes. (This definition has less significance than for full–duplex protocol, since there is no multiplexing of transmission codes in half–duplex protocol).

Half–duplex protocol uses the following control codes:

- DLE SOH
- DLE STX
- DLE ETX BCC/CRC
- DLE ACK
- DLE NAK
- DLE ENQ
- DLE EOT

Half–duplex protocol also uses the following link–layer data codes:

- Data (single bytes having values 00–0F and 11–FF hex)
- DLE DLE (to represent the value 10 hex)
- Link–layer address code
- STN (station identifier)

We can group these codes into two classes according to their use:

- **1.** message codes issued from a station sending a message (or poll)
- **2.** response codes issued from a station receiving a message (or poll).

These codes are issued by a station transmitting a message (or poll):

- DLE SOH indicates the start of a message packet.
- STN helps to designate the station number. When the 1775–KA is communicating with another station as a peer, the $STN = DST$. If the

1775–KA is just one of several stations on a Data Highway, the STN together with the DST identifies the 1775–KA station

- DLE STX separates the data link protocol information from the network packet.
- Link–layer data: (00–0F and 11–FF hex) encodes the bytes of the network packet.
- \blacksquare DLE DLE encodes the value 10 hex in the network packet. This is necessary to distinguish a text code of 10 hex from a DLE control code of 10 hex.
- DLE ETX BCC terminates a message or polling packet.
- \blacksquare DLE ENQ indicates the start of a polling packet.

Response codes from station receiving a message (or poll):

- DLE ACK signals that the receiver has successfully received the last message sent.
- DLE NAK serves as a global link reset command. It causes all slaves to cancel all messages they have ready to transmit to the master. The 1775–KA module responds to this by writing error code 84 into its error word in the PC data table.
- \blacksquare DLE EOT is the response that a slave sends to a poll from the master when the salve has no messages to send.

Half–duplex protocol uses three types of transmissions: Link-Layer Packets

- Polling packet
- Master message packet
- Slave message packet

The master station transmits both polling packets and master message packets, while slave stations transmit slave message packets.

Figure 11.1 illustrates the formats of these packets. Note that the slave message packet has the same format as the full–duplex message packet. The master message packet is the same as the slave message packet except that it is prefixed with DLE SOH and an address code to specify a slave station number.

At the end of each polling packet is a BCC byte. At the end of each message packet is a one–byte BCC field.

Figure 11.1 Formats for Half-Duplex Protocol

a) Polling Packet

b) Slave Message Link Packet

c) Master Message Link Packet

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Block Check

The block check character (BCC) is a means of checking the accuracy of each packet transmission. It is the 2's complement of the 8–bit sum (modulo–256 arithmetic sum) of the slave station number (STN) and all the data bytes in the packet. For polling packets, the BCC is simply the 2's complement of STN. The BCC does not include any other message packets codes or response codes.

For example, if the master station wanted to send the data codes 8, 9, 6, 0, 2, 4, and 3 to slave station 20 hex (40 octal), the master message codes would be (in hex):

The sum of the STN and data bytes in this message packet is 40 hex. The BCC is the 2's complement of this sum, or C0 hex. This is shown in the following binary calculation:

> 0100 0000 40 hex 1011 1011 ls complement $+1$ 1010 0000 2s complement (E0 hex)

To transmit the STN or data value 10 hex, you must use the data code DLE DLE. However, only one of these DLE text characters is included in the BCC sum. For example, to transmit the values 8, 9, 6, 0, 10, 4, and 3 hex, a slave station would use the following message codes:

In this case, the sum of the data bytes is 2E hex because only one DLE text code is included in the BCC. So the BCC is D2 hex.

Protocol Environment Definition

Each station on the multidrop link must contain a software routine, known as a transceiver, that can both transmit and receive message packets. The 1775–KA module already contains a slave transceiver routine, so it will function as a slave station if you select Polled–Subscriber Mode with LIST (chapter 2). To establish master station, you have to program a transceiver routine at a computer. In addition to transmitting and receiving message packets, the master transceiver must also be able to transmit polling packets.

Note that you can program separate transmitter and receiver routines instead of a single transceiver. For purposes of the discussion here, however, we assume that the transceiver is a single software routine.

Figure 11.2 illustrates the operation of master and slave transceivers. To fully define the protocol environment, you must tell the master transceiver where to get the messages it sends and how to dispose of messages it receives. These are implementation–dependent functions that we call the message source and the message sink respectively.

We assume that the message source supplies one network packet at a time upon request from the transceiver, and that the source has to be notified about the success or failure of transfer before supplying the next. Whenever the transceiver has received a link packet successfully, it attempts to give the network packet portion to the message sink. The message sink may be full. The message sink must notify the transceiver when it is full.

Figure 11.2 Slave Transceiver

Message Characteristics

Half–duplex protocol places the following restrictions on the network packet that is submitted to the link layer for transfer.

- The size of a valid network packet is 6 bytes minimum, and 250 bytes maximum.
- The first byte of a network packet must be the station number of the receiver station (see DST in chapter 12). The receiver ignores messages that do not contain the correct station number.
- As part of the duplicate message detection algorithm, the transceiver checks the second, third, fifth, and sixth bytes of each network packet. At least one of these bytes of the current network packet must differ from the corresponding byte of the previous network packet in order for the transceiver to act upon the current network packet. Otherwise, the transceiver assumes that the current network packet is a retransmission of the previous network packet, so it discards the current network packet.

Master Polling Responsibilities

You may vary the master polling algorithm, depending on how much activity you expect on your network.

The master should poll each slave repeatedly until that slave has transmitted all of its messages. The master should then send any messages it has for that slave. Then the master can poll the next slave in the same way.

If a slave station fails to respond to a poll, the master should remove that slave from the list of active slaves. To save time, the master should poll only the active slaves on a regular basis. The master should poll the inactive slaves occasionally to see whether they will respond.

It is best not to allow the master station's transceiver to relay messages directly from one slave station to another. Instead, the transceiver should funnel all received messages to the message sink (network layer). The network layer can then analyze the messages and retransmit any that are addressed to a slave station.

Figure 11.3 is a flowchart which gives a simplified view of an example of software logic for implementing half–duplex protocol from the master station's point of view.

Transceiver Actions

Since the transceiver receives "dirty" input from the physical world, it must be capable of responding to many adverse situations. Some of the things that can conceivably happen are listed here:

- The message sink can be full, leaving the transceiver with nowhere to put message.
- A message can contain a parity error.
- The BCC can be invalid.
- The DLE SOH, DLE STX, or DLE ETX BCC may be missing.
- The message can be too long or too short.
- A spurious control or data code can occur outside a message.
- A spurious control code can occur inside a message.
- Any combination of the above can occur.
- The DLE ACK response can be lost, causing the transceiver to send a duplicate copy of a message that has already been passed to the message sink.

Each slave station is in a passive mode until it receives a DLE ENQ or DLE SOH code. While in a passive mode a slave ignores any transmission code that is not DLE ENQ or DLE SOH.

When a slave receives a DLE SOH, it resets its BCC accumulator and message receiving buffer. The next code it receives must be its specific station number of the global station number 377 (octal). If the packet does not contain the appropriate station number, the slave ignores it and waits for the start of a new transmission.

If a slave receives a message packet with the appropriate station number, it adds the value of that station number to its accumulated BCC. If the next characters after the station number are DLE STX, then the slave transceiver starts storing the incoming link–layer data in a buffer. The transceiver stores all data codes in the buffer and adds these code values to the accumulated BCC. Even if the storage buffer overflows, the transceiver continues summing the BCC, while discarding the data.

The slave also sets an error flag to indicate the occurrence of a parity, buffer overrun, message framing, or MODEM handshaking error. When the slave gets a DLE ETX BCC, it checks this error flag, the BCC, and the message size. If any of these tests fail, the slave ignores the message.

If the current message packet passes the above tests, the slave next begins the duplicate message detection process. In this process, the slave

compares the SRC, CMD, and both TNS bytes with the corresponding bytes of the previous message received. If these bytes are the same, the slave discards the current message and sends a DLE ACK.

If the current message differs from the previous one, the slave next tests the state of the message sink. If the message sink is full, the transceiver discards the current message and does not respond. Otherwise, the transceiver:

- Forwards the current link–level data to the message sink
- Keeps a copy of the first six bytes of the current link–level data for purposes of duplicate message detection
- Sends a DLE ACK

While waiting to receive a message, a slave station could receive a polling packet that begins with a DLE ENQ sequence. The slave will ignore the poll if the polling packet does not contain the slave's station number or if the BCC is the polling block is incorrect. If the poll is valid, then one of three conditions can exist:

- The slave is still holding a message that it had transmitted previously but had not been acknowledged by the master station. There is a limit on the number of times the slave will attempt to transmit a message. If this limit has been exceeded, the slave responds to this by writing an error code into its error word in the PC data table, and then tries to transmit the next message from the message source. If the NAK limit is not exceeded, the slave tries to retransmit the current message.
- If the slave does not currently have a message to send, it tries to get one from the message source. If a message is available, the transceiver initializes its retry counter and transmits the message in response to the poll.
- If not message is available, the transceiver responds to a poll by transmitting a DLE EOT.

To transmit a message, the slave transceiver uses the same message block format as the full–duplex format (section 10.3.2). After sending a message, the transceiver keeps a copy of that message until it receives a DLE ACK from the master station, or until its retry limit is exceeded.

When the slave transceiver receives a DLE ACK, it discards the current message. The next time the slave is polled, it will send the next message available from the message source. If no message is available in the message source, the slave responds to a poll with DLE EOT.

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When the slave transceiver receives a DLE NAK, it takes messages from the source until the source is empty. It discards each message while sending an error code back to the source. The master can use this to clear the message source buffer of each slave after the master has been down.

The following figures show the events that occur on various interfaces. Control characters are shown in bold type. Link–level data is represented by xxxx. Line noise is represented by ???. Each message packet is shown ending in BCC. Time is represented as increasing from the top of the figure to the bottom. Figure 11.4 shows normal message transfer from the master to a slave. Half-Duplex Protocol Diagrams

Figure 11.5 shows a message transfer in which the BCC was invalid. After a timeout, the message is retransmitted. After the retransmission the response is DLE ACK.

Figure 11.5 Message Transfer with Invalid BCC

Figure 11.6 shows a message transfer in which the acknowledgment was destroyed by noise. After a timeout, the message is retransmitted and the DLE ACK response is detected.

Figure 11.6 Message Transfer with ACK Destroyed

Figure 11.7 shows a slave being polled, and responding with DLE EOT because it has no messages to transfer.

Figure 11.7 Poll with No Message Available

Figure 11.8 shows a slave being polled, and answering with a message. Because a block check error is found, the master does not acknowledge; instead, it sends the poll to the slave again. Since the slave did not receive an acknowledgement to its first message transmission, it retransmits the same message in answer to the second poll. The master receives the second transmission of the message with no error and responds with DLE ACK.

Figure 11.8 Poll with Message Returned

Figure 11.9 shows a slave unable to receive the acknowledgement from the master after the master successfully received the message from the slave. Sometime later when the master polls that same slave again, the slave sends the same message again. The master responds with DLE ACK, but discards the received transmission block because it detects it to be a duplicate message received from that slave. With multiple slaves, to implement this duplicate message detection, the master must do either of the following:

• Poll a station repeatedly (without polling any other station until it receives a DLE EOT to be sure it has detected any retransmissions. If each station is polled only once per cycle, the master must keep a record of the first 6 link–level data bytes of the last transmission from each station, since other stations may transfer messages between retransmissions from a given station.

Figure 11.9 Duplicate Message Transmission

When a slave station fails to respond to a message from the master, you should poll the slave to see if it is there. If it answers the poll with a DLE EOT but consistently fails to ACK the master's message, the slave's message sink is probably full. If the slave answers with DLE EOT to a poll, you should wait for the slave's receiver buffers to clear. This situation is illustrated in figure 11.10.
Figure 11.10 Message Sink Full, Case 1

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When a slave station's message source and sink share a common memory pool (as in the 1775–KA module) it may be that the message sink full indication results from an abundance of messages in the message source, which uses up all free pool memory. In this case, the memory can be freed up by receiving messages from that slave station. Waiting for the memory to clear by the action of the slave station alone may not work, since it could be that the only way to fee up space is for the slave to send a message to the master. This situation is illustrated in Figure 11.11.

Figure 11.11 Message Sink Full, Case 2

Line Monitoring

When monitoring half–duplex protocol on a two–wire link, you need to monitor only one line. the example below shows a message sent by the master and a reply sent by the slave in answer to a poll. Slave responses are in bold.

Message from master to slave:

DLE SOH STN DLE STX xxxx DLE ETX BCC **DLE ACK**

Message sent from slave to master in answer to poll:

DLE ENQ STN BCC **DLE STX xxxx DLE ETX BCC** DLE ACK

Poll with a DLE EOT answer:

DLE ENQ STN BCC **DLE EOT**

Chapter 12

Network Layer

The network protocol defines a network packet format for interaction between application programs. The link protocol merely serves to carry data blocks between two applications, regardless of which data link protocol (half or full– duplex) you use. The application programs may be located at opposite ends of a point–to–point full duplex link, or at different points on a multidrop half– duplex link. The network protocol can even handle the transfer of messages between application programs in the same device.

The network layer ignores the internal functioning of data link protocols. It requires that the data link driver accepts a message for delivery, tries to send it, and indicates whether it was delivered.

Program And Message Types

The network protocol was designed on the assumption that application programs are of two types: command initiators and command executors. Corresponding to this division there are two message types:

- Command messages initiated by command initiators and carried over the network to a command executor.
- Reply messages the replies that command executors send to command initiators.

For each command message there is normally one and only one reply. (A rare exception occurs when the data link delivers a message but receives no acknowledgement to verify delivery. At the command initiator, the network layer sends a reply command executor receives and executes the command packet, and sends a single reply message to the command initiator. In any case, the command executor generates only one reply message for each command it receives.)

If the network layer of the command initiator station cannot deliver a command to another station, it generates a reply message with an error code in its own application layer. If a reply cannot be delivered, the network layer destroys it.

Network Model

To implement your Data Highway network layer software, use a routing subroutine and a queue. Messages created by the application are sent to the router for transmission over the network. Messages that are delivered by the network are placed on an incoming message queue that is unique for each application. Figure 12.1 illustrates this model.

Reply messages are not necessarily sent in the same order that their corresponding command messages were received. It is impossible for the network to guarantee delivery, and in some cases it may not be possible to provide notification of non–delivery. Therefore, the command initiator should maintain a timer for each outstanding command message. Non–deliverable reply messages are not returned to the command executor.

The application task is notified via the operating system when a message arrives on the queue. Messages do not necessarily have to be removed from the queue in order of arrival.

Network Packet Fields

As we discussed the communication protocol used on the data link, we described control characters framing the network packet. Here at the network level, you must generate the network packet. In this protocol the network packet characters are generated directly from binary coded bytes of data. This provides faster throughput on the link than if this data were coded into ASCII characters.

Figure 12.2 shows the general format of the network packet for a command message. Figure 12.3 shows the general format of the network packet for a reply message. Note that bytes are shown from left to right in the order in which they are transmitted on the link.

Note that the only difference between the network packet for a command message and the network packet for a reply message is in the high nibble of the CMD byte.

Figure 12.2 Command Message Packet Format

10063–I

Legend: $x =$ low hex digit of CMD byte supplied by application layer

10064–I

DST and SRC

The DST (destination) byte is the number of the station that receives the network packet. The SRC (source) byte is the number of the station that sent the packet. There are 255 possible station numbers from 0 to 254 decimal. You reverse the DST and SRC of the command message to form the DST and SRC of the corresponding reply message.

CMD (High Nibble)

The high nibble of the CMD (command) byte is supplied by the network layer. Bit 6 (26 value) of the CMD byte is the command/reply indicator. It is 0 for command messages and 1 for reply messages. Therefore the high hex digit of the command byte is 0 for command messages and 4 for reply messages. (The low nibble comes from the application layer).

STS (Low Nibble)

The low nibble of the STS (status) byte is supplied by the network layer. In a command message, this field is set to zero. In a reply message reporting no error or a remote error, this field is also set to zero. (The high nibble comes from the application layer.)

If the network layer of your computer cannot deliver a command to another station, it writes a local error code into this field to generate a reply message which it returns to the command indicator in your application layer. All error codes are listed in appendix B.

TNS

The two TNS (transaction) bytes contain a unique 16–bit transaction identifier field. A complete transaction consists of a command message and its corresponding reply message. The TNS value in the reply must be the same as the TNS value in its associated command. This enables the command initiator to associate an incoming reply message with one of the command messages it transmitted previously.

For command messages transmitted by a PC station, the 1775–KA module assigns the TNS values. For each command message transmitted by your computer station, your application programs must assign a unique 16–bit transaction number. A simple way to generate the transaction number is to maintain a 16–bit counter in your application program. Increment the counter every time your command initiator (application program) creates a new message, and store the counter value in the two TNS bytes of the new message.

When your computer program receives a reply to one of its command messages, it can use the TNS value to tie the reply message to its corresponding command. If the TNS value of a reply message matches the TNS value of a command message, then that reply is the appropriate one for that command.

Whenever your computer network layer receives a command from another station, it should copy the TNS bytes of the command message into the same bytes of the corresponding reply message. Do not change the TNS value in a reply message. If you do, the command initiator will not be able to match its command to the corresponding reply message.

Note that the low byte (least significant bits) of your TNS value will be transmitted across the link before the high byte (most significant bits).

At any instant, the combination of SRC, CMD, and TNS values are sufficient to uniquely identify every message packet in transit for duplicate message detection. At least one of these fields in the current message must be different, the command executor ignores the current

message. During an upload or download, the TNS value is the only way to distinguish between the physical read or write reply messages.

Recall from chapter nine that the application layer provides the Data Highway commands that you use to transfer data and manage the network. This function is provided by the command initiators and command executors. Application Layer

At the application layer, the command initiators are responsible for:

- creating a message packet and submitting that packet to the network layer
- maintaining the sequence number and the timeout
- **accepting the reply**
- canceling the timeout and sequence number
- destroying the reply message packet when it is no longer needed

At the application layer, the command executors must:

- create the reply message packet
- copy over certain information from the command
- fill in any reply information
- submit the packet to the network
- destroy the command packet

Application programs communicate by sending information back and forth in the command, status, and data fields of network packets. Application protocols may vary depending on the types of application programs that are communicating.

Application Message Fields

Figure 12.2 shows the general format of the application fields for a command message. Not all command messages have FNC, ADDR, or DATA bytes.

Figure 12.3 shows the general format of the application fields for a reply message. Not all reply messages have DATA bytes.

In addition to the application layer fields shown in these figures, some of the PLC–3 command messages also contain these application layer fields:

- EXT STS (extended status)
- Packet Offset
- **Word Offset**
- TOTAL TRANS (total transaction size)

Appendix A lists the message formats (command and reply) of every command the PLC–3 can send or receive.

CMD and FNC (command and function)

For these message formats which include an FNC byte, the low nibble of the CMD (command) byte together with the FNC (function) byte define what action the command executor at the destination station will perform. For those message formats which do not include an FNC byte, the CMD byte alone defines what action the command executor at the destination station will perform. The hexidecimal values of the CMD and FND bytes are listed in Table 12.A for each of the types of messages that can be transmitted across this link.

Bits 0 through 3 of the CMD byte must be the same in the reply message as it is in the corresponding command message. In current implementations this is either a command code or a command executor selector. The application program must always copy this field from the command to the reply message.

The commands that the PLC-3 can send and/or receive, and the hexadecimal values for the CMD and FNC bytes

[1] RS-232-C device means a computer or intelligent terminal.

[2] The "Basis" commands can be sent to any Allen-Bradley PC. These commands are sometimes called PLC/PLC-2 commands, but all Allen-Bradley PC's can receive them. The PLC, PLC-2 family, and PLC-3 processors can also send these commands as well as receive them. [3] Allen-Bradley recommends using these commands for uploading or downloading only. To write or read specific words or bits use the "basic"

commands or the PLC-3 commands.

STS (status)

The high nibble of the STS (status) byte is supplied by the application layer. In command messages the STS byte is set to zero.

In reply messages the STS is used for reporting either application or network error codes. A value of zero should be interpreted as no error (that is, the message was delivered and executed successfully). Non–zero status can be divided into two categories: remote errors and local errors.

Remote errors mean that a command was successfully delivered by the network, but the remote station was unable to execute the command. The remote station then placed an error code in the high nibble of the STS byte.

Local errors mean that your network layer was unable to deliver the message to the remote station. Your network layer then turns the command around, stuffs the low nibble of the STS byte with the appropriate error code, and returns it to your application.

All error codes are listed in appendix B.

When you receive a reply message from a PC station, check the STS byte at the application layer. If the STS byte is non–zero, refer to appendix B for the type of error that has occurred.

If your application layer receives a command message and detects an error, it should format a reply message with a remote error code in the high nibble of the STS byte.

ETX STS (extended status)

If the PLC–3 receives PLC/PLC–2 commands, error codes for those commands are returned in the STS byte only.

The PLC–3 can also create a second layer of error codes, however, relative to PLC–3 type commands (CMD byte–15). If the command is a PLC–3 level command addressed to a remote PLC–3, then the remote error returned from the 1775–KA will have an additional status byte stuffed into the data area, called an ETX STS.

If the STS byte is zero, then the ETX STS will also be zero, indicating no error. If the STS bytes contain the value F0 hex, this is a flag to indicate that the ETX STS contains the non–zero code.

To decode the contents of the STS byte and the ETX STS byte relating to the application programs of specific processors, refer to Appendix B error codes (80–88) for remote errors and error codes (90–97) for local errors.

ADDR (address)

The address field in command messages can be in one of the following formats:

- PLC/PLC–2 addressing format
- Symbolic addressing format
- Logical addressing format
- Physical addressing format

The **PLC/PLC–2 addressing forma**t applies to PLC/PLC–2 type commands transmitted to the 1775–KA module. Use this addressing format whenever you have established a PLC–3 input file to imitate PLC/PLC–2 memory (section titled PLC/PLC–2 Stations, chapter 3).

The ADDR (address) field is a 2–byte address field sent low byte first. PC programs use logical addressing to specify octal bytes. To generate a protected/unprotected read/write command, use that same octal addressing to format the ADDR field as shown in Figure 12.4. For a block read/write command, always set the least significant bit to 0 to select the low byte of the word. To allow PC read/write commands to the computer, set up a file in the computer to be addressed in this same way as if it were a PC data table.

Since ADDR specifies an address as the number of bytes from the beginning of PC memory, its value is double the corresponding PC word address.

The **PLC–3 logical addressing format** also applies to PLC–3 type commands. You can use this format to specify up to 6 levels of PLC–3 extended addressing. Figure 12.5 shows an example of the logical addressing format for addressing a word in the PLC–3 data table.

The first field in the format contains a set of bit flags. Each flag is associated with one of the levels of a PLC–3 extended address. If a flag bit is set to 1, there must be an address specification for the corresponding level in the address fields that follow. If a flag bit is zero, the address fields that follow should not contain an address specification for that level; instead, a default value is assumed.

You must always specify the value for the lowest level of the desired extended address, even if it is the default value. Since the present PLC–3 recognizes a maximum of 7 levels of extended addressing, you cannot specify more than 7 levels with the logical addressing format.

If the address fields can be specified in one byte each, then you can code the values directly. If it takes two bytes to specify an address field, then you must use a delimiter byte of value FF hex before each 2–byte field. Any 2–byte fields should be coded low–byte–first.

In Figure 12.5, the first byte contains the bit flags to indicate which addressing levels are specified. In this example, only levels 3, 4, and 6 are specified; default values are used for the other levels. This format reduces the total number of bytes needed to specify a PLC–3 logical address in a command message.

In Figure 12.5, the level–4 address is 260 (decimal), which is too large to fit in one byte. Therefore, a byte of all 1's is used to delimit the 2–byte address field for this level. The value 260 is then coded low–byte–first. Note that the last level (level 6 in this case) must be specified in the address field even though it is equal to the default value of zero.

PLC–3 Extended Address

Logical Addressing Format

Flag bit for level:

- Byte 1 –– is the flag byte. In this case it indicates that the addresses for levels 3, 4, and 6 are specified in the bytes that follow. Default values are used for the levels 1, 2, and 5.
- Byte 2 –– is the value of the level–3 address.
- Byte 3 is a delimeter that says the next two bytes are one address.
- Byte 4 –– is the low byte of the level–4 address.
- Byte 5 –– is the high byte of the level–4 address. Note that bytes 4 and 5 together give a value of 260 for the level–4 address.
- Byte 6 is the value of the level-6 address. Even though it is the default value, it must be specified because it is the last level in the desired extended address.

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The **symbolic addressing format** applies to PLC–3 type commands (Table 12.A, appendix A) transmitted to the 1775–KA module. You can use this addressing format whenever you have defined a system symbol to represent a symbolic address at the PLC–3 station that is to receive the command message. Figure 12.6 shows this format for PLC–3 symbolic addresses. Always enter zeros for the first and last bytes of the symbolic address field. Between these zero delimiter bytes, enter the ASCII codes for the 1 to 8 characters of the symbol name. If the symbol name is more than 8 characters long, enter only the first 8 characters.

Figure 12.6 Format for PLC-3 Symbolic Address

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The **physical addressing format** applies only to PLC–3 physical read and physical write commands. Physical word addresses run in sequence, starting with 0 (zero) for the first word of PLC–3 memory. Physical addresses occupy 2 words (4 bytes), in the following bit format:

In this format, A1 through A24 represents the 1 to 24 bits of the physical address. For example, to address a command message to physical word address 12,200 decimal (002FA8 hex) you would use the following binary code in the address field:

Packet Offset

A single message packet cannot transmit more than 244 bytes of user data. To allow for data transfers of more than 244 bytes, the 1775–KA module automatically transmits as many packets as necessary to complete the message. Thus, a message may consist of several packets strung together.

The packet offset is the difference (in 2 byte words) between the starting address for the current packet and the starting address for the first packet in the message. The packet offset for the first packet in the message is always 0 (zero). Thus, adding the packet offset to the address field specified in the message gives the destination address where the current packet begins in the destination station.

Appendix

Message Formats

Introduction

This appendix presents the detailed message formats for each type of command and reply message that the PLC–3 can send and/ore receive. We discuss the message format in the following order:

Basic Command Set

Protected bit write Protected write Unprotected bit write Unprotected read Unprotected write

PLC–3 Commands

Bit write Word range read Word range write File read File write

Privileged Commands

Privileged read Privileged write Shutdown request Download request Upload request Restart request Diagnostic counters reset Diagnostic loop Diagnostic read Diagnostic status *Set ENQs *Set NAKs *Set timeout *Set variables

*Use these commands to affect only the RS–232–C port of the 1775–KA module.

Important: In the formats shown in this section, CMD and FNC values are expressed in hexadecimal notation. All other values are given in decimal form. Network layer fields are shaded in blue; data link layer fields are shaded in grey; application layer fields are not shaded.

Building a Line Monitor

This appendix presents the formats for each of the Data Highway commands that the 1775–KA module can send and/or receive. You can build a line monitor that allows you to see these commands as they are sent or received by the 1775–KA module. This can be useful during installation or when you are troubleshooting. By viewing these commands on a line monitor, you can verify the types of messages the 1775–KA is sending or receiving and the data that is contained in the messages.

We do not recommend building a line monitor for Data Highway communication. Messages on the Data Highway use a different protocol. But you can monitor communication to the RS–232–C link. To do this, you must create a cable like the one shown in Figure A.1. This cable links the 1775–KA module, an RS–232–C device, and a 1770–T4 Industrial Terminal (Figure A.2). The cable you build connects the receiving data, transmit data, and ground in the RS–232–C link to the receive data and ground of the Industrial Terminal. A switch on the cable allows you to choose the monitor either messages received or messages sent by the 1775–KA module.

 $A-3$

To program the industrial terminal to act as a line monitor, complete the following steps:

- **1.** Attach the 1770–KDA programming keyboard to the industrial terminal.
- **2.** Press SHIFT MODE and type 2 to select the alphanumeric mode.

In the steps that follow, you type a letter key repeatedly until the selection you want appears on screen.

- **3.** Type A to select the baud rate you are currently using on the line.
- **4.** Type B to select no parity.
- **5.** Type C to select 1 stop bit.
- **6.** Type E to select cursor on.
- **7.** Type F to select full–duplex.

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- **8.** Type G to select auto linefeed off.
- **9.** Type I to select control code display on.
- **10.** Press ENTER to put these selections into effect.

Now as you send a command from the 1775–KA or the RS–232–C device, ASCII characters will appear on the industrial terminal. Use Table A.A to find the command bytes (DLE, FNC, etc.) that these characters represent. You may not care what the actual numeric value is for control characters such as DLE STX; on the other hand, you will probably want to know the numeric value of DATA sent with a command. Figure A.3 shows an example of a command as it might appear on the line monitor, and how it is translated.

Figure A.3 Typical monitor display and how it is interpreted

Monitoring a diagnosticc status command and reply acknowledgement.

Line monitor displays:

Examining the ASCII table shows that for the above message:

 $CMD = 6$ $STS = 0$ $FNC = 3$

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Basic Command Set

The basic command set includes those commands that can generally be executed by any PC station on the communication link, regardless of the type of PC controller at the station. In some cases, switch settings on the station interface module can disable execution of a particular type of command at that station. For more details, refer to the user's manual for the station interface module.

The basic commands are:

- protected bit write
- protected write
- unprotected bit write
- unprotected read
- unprotected write

Protected Bit Write

Use this command to set or reset individual bits within limited areas of the PC data table memory. Your access can be limited by memory access rungs in the communication zone of the PC's ladder diagram program.

The data field in this packet consists of 4–byte blocks, each of which contains a 16–bit address field, a set mask, and a reset mask. Use the ADDR field to specify the address of the byte to be modified in the PC data table memory. Put the low byte (lest significant bits) of the PC address value into the first byte of the ADDr field.

Use the SET mask to specify which bits to set to 1 in the addressed PC byte. A1 in a bit position of the SET mask means to set the corresponding bit in the addressed PC byte to 1; a 0 in a bit position of the SET mask means to leave the corresponding bit in the PC byte unchanged.

Use the RESET mask to specify which bits to reset to 0 in the addressed PC byte. A1 in a bit position of the RESET mask means to reset the corresponding bit in the addressed PC byte to 0; a 0 in a bit position of the reset mask means to leave the corresponding bit in the PC byte unchanged.

Note that the interface module at the receiving PC station executes this command by first making a copy of the addressed PC byte. It then sets or resets the appropriate bits and writes the byte back into PC memory. At the same time, the PC processor can be changing the states of the original bits in memory. Because of this, some data bits may unintentionally be overwritten.

Command Format:

Reply Format:

Protected Write

Use this command to write words of data into limited areas of the PC data table memory. Your access can be limited by memory access rungs in the communication zone of the PC's ladder diagram program.

Command Format:

Reply Format:

Unprotected Bit Write

Use this command to set or reset individual bits in any area of PC data table memory. The data field in this packet consists of 4–byte blocks, each of which contains a 16–bit address field, a set mask, and a reset mask. Use the ADDR field to specify the address of the byte to be modified in the PC data table memory. Put the low byte (least significant bits) of the PC address value into the first byte of the ADDR field.

Use the SET mask to specify which bits to set to 1 in the addressed PC byte. A1 in a bit position of the SET mask means to set the corresponding bit in the addressed PC byte to 1; a 0 in a bit position of the SET mask means to leave the corresponding bit in the PC byte unchanged.

Use the RESET mask to specify which bits to reset to 0 in the addressed PC byte. A 1 in a bit position of the RESET mask means to reset the corresponding bit in the addressed PC byte to 0; a 0 in a bit position of the RESET mask means to leave the corresponding bit in the PC byte unchanged.

Note that the interface module at the receiving PC station executes this command by first making a copy of the addressed PC byte. It then sets or resets the appropriate bits and writes the byte back into PC memory. At the same time, the PC processor can be changing the states of the original bits in memory. Because of this, some data bits may unintentionally be overwritten.

Command Format:

Message Formats Appendix A

Reply Format:

Unprotected Read

Use this command to read words of data from any area of PC data table memory. Use the SIZE field to specify the number of bytes to be read. To specify a number of PC words, SIZE should be an even value because PC words are two bytes long.

Command Format:

Reply Format:

Unprotected Write

Use this command to write words of data into any area of PC data table memory.

Command Format:

Reply Format:

PLC-3 Commands

PLC–3 stations can receive any of the commands in the basic command set and execute them within a specified file in the PLC–3 memory. They can also execute the following commands, which apply only to PLC–3 controllers:

- bit write
- word range read
- word range write
- file read
- file write

Only a computer can send privileged commands. Their primary use is for uploading and downloading PLC–3 memory.

Only a computer or another PLC–3 station can initiate the non–privileged PLC–3 commands listed above. Their primary use is for transferring data between two PLC–3 files. Those files may be located in the same PLC–3 processor or in two different PLC–3's

Bit Write

Use this bit write command to modify the bits at the address specified by either a word symbol, a file symbol plus a word offset, or a logical address. This write command can write a block of data. This address must point to a word within a file. The function code is 2. Unlike the unprotected and protected bit writes in the basic command set, this command can be used to change the bits in a single word only.

Command Format:

A. Word symbol address

B. File symbol address plus word offset

C. Logical address

Reply Format:

This the same as the reply packet format for all bit writes

A. Format when successful execution

B. Format when reporting an error

Where the extended status byte is optional

Word Range Read

Use this read command with a word symbol, a file symbol plus a word offset, or a block address as a starting address. This starting address must point to a word in a file. This read command can read a block of data. The function code is 1. A special case of this command is the single–word read, where the number of bytes to read is only two bytes (one word).

Command Format:

A. Word symbol address

B. Format when reporting an error

Where the extended status byte is optional.

Word Range Write

Use this write command with a word symbol, a file symbol plus a word offset, or a logical address as a starting address. This starting address must point to a word in a file. This write command can write a block of data. The function code is 0 (zero). A special case of this command is the single word write, where the data field is only one word long.

Command Format:

A. Word symbol address

B. File symbol address plus word offset

C. Logical address

Reply Format:

This is the same as the reply packet format for all writes.

A. Format when successful execution.

B. Format when reporting an error

Where the extended status byte is optional.

File Read

Use this read command with either a file symbol or a block address for a starting address. This starting address points to a file of words. This read command reads a block of data. You must read the entire file. The file size must equal the exact size of the file or an error will be returned. The function code is 4.

Command Format:

A. File symbol address

B. Logical address

Reply Format:

This is the same as the reply packet format for all reads.

A. Format when the command was successfully executed

B. Format when reporting an error

Where the extended status byte is optional.
File Write

Use this write command with either a file symbol or a block address as a starting address. This starting address points to a file of words. This write command can write a block of data. You must read the entire file. The file size must equal the exact size of the file or an error will be returned. The function code is 3.

Command Packet Format:

A. File symbol address

B. Logical address

Reply Format:

This is the same as the reply packet format for all writes.

A. Format when the command was successfully executed

B. Format when reporting an error

Where the extended status byte is optional.

Privileged Commands

A PLC–3 receives privileged commands from an RS–232–C device (such as a computer); the PLC–3 does not send these commands. The privileged commands are:

- privileged read **download request**
- privileged write **upload request**
- diagnostic counters reset restart request
- diagnostic loop set ENQs
- diagnostic read set NAKs
- diagnostic status set timeout
-
-
-
-
- -

Privileged Read

Use this read command with a PLC–3 physical address as a starting address. You use this command to upload from a PLC–3 to a computer. The destination 1775–KA module will accept this command only after the source station has successfully transmitted a shutdown request. The function code for this command is 9.

Command Format:

Reply Format:

TDATA–Max of 238 bytes or 119 words DLE ETX BCChis is the same as the reply packet format for all reads.

A. Format when the command was successfully executed

B. Format when reporting an error

Where the extended status byte is optional.

Privileged Write

Use this write command with a PLC–3 physical address as a starting address. You use this command to download to a PLC–3 from a computer. The destination 1775– KA module will accept this command only after the source station has successfully transmitted a shutdown request. The function code for this command is 8.

Command Format:

Reply Format:

This is the same as the reply packet format for all writes.

A. Format when the command was successfully executed

B. Format when reporting an error

Where the extended status byte is optional.

From a computer you use this command to ask the 1775–KA module to initiate either a PLC–3 shutdown (if the computer has download privileges) or a freeze on file allocations (if the computer has upload privileges). This command halts program and I/O scanning. You cannot issue this command until you have successfully transmitted an upload or download request to the 1775–KA module. This command has a function code of 7.

Command Format:

Reply Format:

A. Format when the command was successfully executed

B. Format when reporting an error

Where the extended status byte is optional.

Download Request

A computer can use this command to inform the 1775–KA module that it wants to do a download. If the 1775–KA module grants the download privilege, the computer may begin issuing privileged writes. You should, however, issue a shutdown command first. If a different station already has the download privilege, the second station is denied the privilege. The function code is 5.

Command Format:

Reply Format:

A. Format when the command was successfully executed

B. Format when reporting an error

Where the extended status byte is optional.

Upload Request

From a computer you use this command to inform the 1775–KA module that it wants to do an upload. If the module grants the upload privilege, you may begin issuing privileged reads. (You should, however, issue a shutdown request first.) If a different station already has the upload privilege, the second station is denied the privilege. The function code is 6.

Command Format:

Reply Format:

A. Format when the command was successfully executed

B. Format when reporting an error

Where the extended status byte is optional.

Restart Request

From a computer you use this command to terminate an upload or a download. You cannot issue this command until after you have successfully completed an upload or download operation with the destination station. This command causes the 1775–KA module to revoke the upload and download privileges for the source computer station and to initialize a PLC–3 restart. The function code for this command is 0A.

Command Format:

Reply Format:

A. Format when the command was successfully executed

B. Format when reporting an error

Where the extended status byte is optional.

Diagnostic Counters Reset

Use this command to reset to zero all the diagnostic timers and counters in the station interface module. The diagnostic status command gives the starting address for this block of counters and timers.

Command Format:

Reply Format:

Diagnostic Loop

You can use this command to check the integrity of transmissions over the communication link. The command message transmits up to 243 bytes of data to a station interface module. The receiving module should reply to this command by transmitting the same data back to the originating station.

Command Format:

Diagnostic Read

You use this command to read up to 244 bytes of data from the PROM or RAM of the station interface module. You can use it to read the module's diagnostic timers and counters. Use the diagnostic status command to obtain the starting address of the diagnostic counters.

Command Format:

Reply Format:

Diagnostic Status

You use this command to read a block of status information from the station interface module. The reply to this command contains the status information in its DATA field.

Command Format:

Reply Format:

The status information varies with the type of station interface module. Table A.B describes this status DATA for 1775–KA modules.

Table A.B Contents of Status DATA for 1775-KA Modules

Set ENQs

Use this command to set the maximum number of ENQs that the station interface module will issue per message transmission. Put the number in the DATA field. The default setting for the KE/KF module is 10 ENQs per transmission.

Command Format:

Set NAKs

Use this command to set the maximum number of NAKs that the station interface module will accept per message transmission. Put the number in the DATA field. The default setting for the KE/KF module is 3 NAKs per transmission.

Command Format:

Reply Format:

Set Timeout

Use this command to set the maximum amount of time that the station interface module will wait for an acknowledgment to its message transmission. The setting is expressed as the number of cycles of an internal clock, where 40 cycles equals 1 second. Put the number of desired cycles in the DATA field. The default setting for the KE/KF module is 128 cycles, or about 3 seconds.

Command Format:

Set Variables

Use this command to set the timeout and maximum NAKs, and ENQs all at once. Put the timeout in the first byte of the DATA field, the NAK setting in the second byte, and the ENQ setting in the third byte. If you do not specify a data value for any one the variables in this command, that variable is automatically reset to zero.

Command Format:

Error Codes

Note that a value of F0 (hex) in the STS byte indicates that the EXT STS byte actually contains the error code for the reply message. Currently, only the 1775–KA module is capable of generating and accepting reply messages with error codes reported in this way. In particular, 1771–KA and 1774–KA modules cannot interpret these error codes.

The meaning of each error code depends on the command message that the local PLC–3 station receives from a remote station. Section titled Local and Reply Error Codes describes the error conditions that the various commands can generate. The error codes are listed according to the decimal value that would be stored at the command initiating station. When a remote station transmits a command, the local 1775–KA module might issue a reply message that contains one of the error codes listed in section titled Local and Reply Error Codes. Error codes 81 to 87 appear in the STS byte of the reply message, and codes 231 to 241 appear in the EXT STS byte (Appendix A).

The local PLC–3 station receives remote error codes in a reply to a command it has sent to a remote station. These error codes are stored under user symbol ERROR in the local PLC–3 station. Remote Error Codes

The extended address for the beginning of the error block file is

\$E2.5.nn.4.0

where nn is the thumbwheel number of the 1775–KA module. You can access this error block by any one of the following means:

- displaying it through the front panel of the PLC–3 controller
- using the data monitor mode of the Industrial Terminal (cat. no. 1770–T4)
- using the move status (MVS) command in the PLC–3 ladder diagram program
- using the I/O Scanner–Message Handling Module (cat. no. 1775–S4B)
- using the 1775–KA module

The meaning of a particular remote error code will vary, depending on the type of communication interface module at the remote station. For example, if the remote station is a PLC–3 processor with a 1775–KA interface module, the remote error codes will have the meanings listed in section tilted Local and Reply Error Codes. For the meanings of other remote error codes, refer to the appropriate user's manual for the communication interface module at the remote station.

Local and Reply Error Codes

Local and Reply Error Code Listing for the PLC–3 Processor

Remote Error codes received from the 1771-KE/KF, 1771-KG, 1771-KA, and 1774-KA Modules

If the remote station has a 1771–KG, 1771–KA, or 1774–KA module, the remote error codes will have the meanings listed in the table below:

Remote Error Codes Received from the 1773-KA Module

If the remote station has a 1773–KA module, the remote error codes will have the meanings listed in the table below:

Diagnostic Counter Block

Data Highway Port Counters

NOTES:

[1] The address of the first byte of the counter block can be determined by using the diagnostic status command. Addresses of the other bytes listed here can be derived by adding the appropriate number of bytes to the staring address of the counter block. [2] An acknowledgment is part of the data highway protocol.

Modem Port Counters

Detailed Flowcharts

Overview

The flowcharts in chapter 10 and 11 give a simplified view of an example of software logic for implementing full–duplex protocol. In this appendix, we present flowcharts which give a detailed view of an example of software logic for implementing full–duplex protocol.

We have not shown any error checking or recovery relating to interaction with the modem handshake driver, a third process. To do this would overly complicate the flow charts, and in many cases, such error checking and recovery are not needed.

Figure D.1 Data Flow Program for Full-Duplex Protocol

10071–I

Figure D.3 Receiver Routine for Full-Duplex Protocol

Figure D.5 SENDM Subroutine

10075–I

10076–I

Figure D.7 STOPTIME Subroutine

Detailed Flowcharts Appendix D

Figure D.8 TIMEOUT Subroutine

Scheduled By:

• STARTTIME

Aborted By:

• STOPTIME

Figure D.9 GETMSG Subroutine

UART Sharing

Figure D.11 shows the transmitter (XMIT) and receiver (RCVE) routines sharing the transmit side of the UART. XMIT transmits messages and inquiries. RCVE transmits responses. Therefore, each must wait until the other is thru before it can transmit thru the UART. The TXALLOC and TXFREE subroutines work together to ensure that XMIT and RVCE do not try to use the UART at the same time. TXALLOC and TXFREE are called in the SENDCTRL, SENDETX, and SENDDATA subroutines (Figure D.12, Figure D.13, and Figure D.14).

Figure D.11 Sharing the Transmit Side of the UART

Figure D.12 SENDCTL Subroutine

Figure D.13 SENDTX Subroutine

Input:

• BCC

Common:

• RESP: The Response Code Variable

Figure D.15 SENDDATA Subroutine

11667

Figure D.17 TXFREE Subroutine

Common:

- IN–USE Flag
- TXALWT: A Sleep Variable

Figure D.18 TRANSMIT INTERRUPT Subroutine

NOTE: This figure assumes the use of a Z80 S10.

NOTE: UART transmit interrupt must be enabled and disabled without affecting the current state of the **receive** and **status interrupt**enable flags.

10088-I

10089-I

Figure D.19 SLEEP and WAKEUP Subroutines

Figure D.20 SLEEP and WAKEUP Interaction

SLEEP and WAKEUP

The SLEEP and WAKEUP subroutines are always used in connection with some type of indivisible process interlock. Indivisibility is achieved on many processors (as on the Z80) by disabling processor interrupts. For this reason, SLEEP and WAKEUP assume that interrupts are off when they are called. They will always return with interrupts off.

When one process calls SLEEP, the results is a return from a call to WAKEUP by another process. When a process calls WAKEUP, the result is a return from a call to SLEEP by another process. An interrupt subroutine that calls WAKEUP is viewed as a subroutine of the interrupted process.

Figure D.20 shows an example of interaction between SLEEP and WAKEUP. In this example, process B work up process A some time ago. Now at 1, when A goes to sleep, actual execution resumes after the wakeup call in B at 3 and 4. Some time later, process C (at an interrupt, for example) calls WAKEUP at 5. Execution flow proceeds to the instructions at 8 following the call to sleep in process A. The next time A calls SLEEP, the WAKEUP call in C will terminate.

This is not the only possible implementation of SLEEP and WAKEUP. A second alternative implementation would allow a process to call

WAKEUP without losing immediate control of the processor. If B wakes up A, context switching would be deferred until B itself has executed a SLEEP.

A third alternative would cause a context switch if a process performed a WAKEUP on a higher priority process. If a WAKEUP had been performed on a lower priority process, the context switch would be deferred until the first process has gone to sleep.

The first alternative is the result of implementing the driver totally at the interrupt level where scheduling is dictated by the interrupt daisy chain hardware.

The third alternative would be used if the driver were implemented as tasks under a multitasking operating system. Such an implementation might be easier, but would probably be limited to lower communication rates.

POWERUP

At powerup, the Z80 processor starts executing code at location 0. The POWERUP subroutine starts at XMIT and RCVE processes by executing a SPAWN. A SPAWN is very similar to a WAKEUP except that the corresponding SLEEP is imaginary and is located prior to the first instruction of the spawned process.

Figure D.21 POWERUP Routine

Figure D.23 UNLINK Subroutine

Input:

Address of Queue •

Output:

 Message Control Block •

Figure D.24 LINK Subroutine

Figure D.26 GETCODE Subroutine

Figure D.28 Receive Interrupt Subroutine

NOTE: This figure assumes the use of a Z80 S10

10098-I

Figure D.29 SENDNET Subroutine

Input:

• Message Buffer

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