





DUTYS v1.0 USER'S MANUAL

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Created by:

Status - Version: Date: Developed during: Kostas Siozios, George Koutroumpezis, Konstantinos Tatas and Dimitrios Soudris Shareware - 1.0 07/11/2006 AMDREL project, IST-2001-34379 DUTYS (**D**emocritus **U**niversity of Thrace Architecture file generator-synthesizer) creates the architecture file of the FPGA that is required by T-VPack and VPR. The architecture file contains a description of various parameters of the FPGA architecture, including size (array of CLBs), number of pins and their positions, number of BLEs per CLB, plus interconnection layout details such as relative channel widths, switch box type, etc. It has a GUI that helps the designer select the FPGA architecture features and then automatically creates the architecture file in the required format.

1. Function of DUTYS

DUTYS creates the FPGA architecture file in the required format. Each line in an architecture file consists of a keyword followed by one or more parameters. In the description below, strings between curly braces, {}, denote all the possible choices for an option. All of the following keywords must be specified in the architecture file.

1.1. Description of I/O pads

io_rat <int>: Sets the number of pads (inputs or outputs) that fit into the space occupied by one CLB. This is the number of pads in each row or column of the FPGA.

1.1.1. Description of Relative Channel Widths in the FPGA

The next three keywords are used to describe the relative widths of the various channels in the FPGA (Fig. 1). If global routing is to be performed, channels in different directions and in different parts of the FPGA can be set to different relative widths. *If detailed routing is to be performed, however, all the channels in the FPGA must have the same width.*

chan_width_io <**float>:** Width of the channels between the pads and core relative to the widest core channel.

chan_width_x {gaussian | uniform | pulse | delta} <peak> <width> <xpeak> <dc>: The italicized quantities are needed only for pulse, gaussian, and delta (which does not need width). Most values are from 0 to 1. Sets the distribution of tracks for the *x***-directed channels -- the channels that run horizontally.**

- *uniform*: if specified, you simply specify one argument, peak. This value (by convention between 0 and 1) sets the width of the *x*-directed core channels relative to the *y*-directed channels and the channels between the pads and core. Fig. 1 should make the specification of uniform (dashed line) and pulse (solid line) channel widths more clear.
- *gaussian*: This keyword takes the same four parameters as the pulse keyword, and they are all interpreted in exactly the same manner except that in the gaussian case width is the standard deviation of the function.
- *delta*: A function used to specify a channel width distribution in which all the channels have the same width except one. The syntax is **chan_width_x delta peak xpeak dc**.
- *peak:* The extra width of the single wide channel.
- *xpeak:* A parameter between 0 and 1 and specifies the location within the FPGA of the extra-wide channel -- it is the fractional distance across the FPGA at which this extra-wide channel lies.

dc: specifies the width of all the other channels. For example, the statement chan_width_x delta 3 0.5 1 specifies that the horizontal channel in the middle of the FPGA is four times as wide as the other channels.

chan_width_y {gaussian | uniform | pulse | delta} peak <width> <xpeak>
<dc>: Sets the distribution of tracks for the y-directed channels.



Fig. 1: Specification of relative channel widths.

1.2. Logic Block Description

inpin class: <int> [global] {top | bottom | left | right} {top | bottom | left | right} : Declares an input pin, determines the class to which this pin belongs, and sets the side(s) of CLBs on which the physical output pin connection(s) is (are). All pins with the same class number are logically equivalent -- such as all the inputs of a LUT. Class numbers must start at zero and be consecutive. The global keyword is optional; if specified, it comes after the class number. Global input pins can connect only to signals marked as global in the netlist (typically clocks). Global input pins are not connected into the normal routing; it is assumed they connect to a special, dedicate resource used for special nets like clocks.

outpin class: <int> {top | bottom | left | right} {top | bottom | left | right} : All parameters have the same meanings as their counterparts in the inpin statement. *NOTE: The order in which your inpin and outpin statements appear must be the same as the order in which your netlist (.net) file lists the connections to the CLBs.* For example, if the first pin on each CLB in the netlist file is the clock pin, your first pin statement in the architecture file must be an inpin statement defining the clock pin. Pads are always assumed to have only one pin (either an input or an output), and this pin is accessible from the one channel bordering that pad. Hence no inpin or outpin statements are given for pads.

subblocks_per_clb <int>: Specifies the maximum number of subblocks, or BLEs, in each logic block. This information is used only for timing analysis.

subblock_lut_size <int>: The number of LUT inputs to each of the subblock BLEs (i.e. K). Again, this information is only needed for timing analysis. Even if your logic block is not constructed from BLEs, it is possible to describe the timing relations between inputs and outputs in terms of BLEs, as one of the examples below illustrates.

The listing below is for an FPGA with all channels of the same width, and a CLB compatible with that produced by T-VPack with the **-no_clustering** option. This CLB contains a 4-input LUT and a flip flop; the input pins are listed first, followed by the CLB output pin, followed by the clock pin. Notice that the four inputs all have the same pin class, indicating that they are logically equivalent and the router may connect nets to any one of them. Notice also that pins can be physically accessible from several sides.

Uniform channel architecture, 4-input LUT and a FF (one BLE) per lb. io_rat 2 #2 Pads per row or column. chan width io 1 #Same as core channels. chan width x uniform 1 #All same width chan_width_y uniform 1 # 4-input LUT. LUT inputs first, then output, then clock. inpin class: 0 bottom top #Physical pins at both top and bottom of #clb. inpin class: 0 left right inpin class: 0 bottom top inpin class: 0 left right outpin class: 1 top bottom inpin class: 2 top # Clock pin # Class 0 is LUT inputs, class 1 is the output, class 2 is the clock # in this case. subblocks per clb 1 # One BLE in each logic block subblock_lut_size 4 # The LUT in a BLE has 4 inputs

As a second example of an architecture file, consider a logic block consisting of a cluster-based logic lock, where each logic block has 5 inputs for use by its BLEs, 2 outputs and one clock input. Each logic lock contains two separate BLEs, and each BLE consists of a 4-input LUT and a flip flop. If the .net file as created with T-VPack, the pin ordering we need to match the .net file is: inputs for use by BLEs outputs, clock.

```
Uniform
             channel
                       architecture,
                                      cluster-based
                                                       loqic
                                                              block
#
containing
# 2 BLEs.
io_rat 2 #2 Pads per row or column.
chan_width_io 1 #Same as core channels.
chan_width_x uniform 1 #All same width
chan_width_y uniform 1
# Logic block with 2 BLEs. 5 Inputs for use by BLEs first, then two
# outputs, then the clock.
inpin class: 0 bottom
inpin class: 0 left
```

inpin class: 0 right inpin class: 0 top inpin class: 0 bottom outpin class: 1 top bottom #Output 1 outpin class: 1 left right #Output 2 inpin class: 2 global top #Clock -> accessible only by global nets in # this case # Class 0 is LUT inputs, class 1 is the output, class 2 is the clock # in this case. subblocks_per_clb 2 # Two BLEs in each logic block subblock_lut_size 4 # The LUT in a BLE has 4 inputs

Notice that all the inputs are of the same class, indicating they are all logically equivalent, and all the outputs are of the same class, indicating they are also logically equivalent. This is true of all cluster-based logic blocks, as the local routing within the block provides full connectivity. However, for most logic blocks all the inputs and all the outputs are *not* logically equivalent. For example, consider the logic block in Fig. 2, which consists of a 3-input and gate and a 2-input or gate. In this case, the set {in1, in2, in3} is logically equivalent, and could all be made class 0. Similarly, the set {in4, in5} is logically equivalent, so each must be different class, say class 2 and class 3.

```
inpin class: 0 top #in1
inpin class: 0 left #in2
inpin class: 0 right #in3
inpin class: 1 bottom #in4
inpin class: 1 right #in5
outpin class: 2 left #out1
outpin class: 3 top #out2
```

If we want to perform timing analysis on the logic block of Fig. 7, we must describe the timing relationship between the inputs and outputs. Clearly out1 depends only on in1, in2 and in3, while out2 depends only on in4 and in5. Therefore we could model this logic block as consisting of two BLEs, with each BLE having 3 inputs.

subblocks_per_clb 2
subblock_lut_size 3

One line of a .net file of a circuit made out of such logic blocks might therefore be:

.clb block_1
pinlist: in1 in2 in3 in4 in5 out1 out2
subblock: and_gate 0 1 2 5 open # out1 depends on in1, in2 and
in3,
 # and is not registered.
 subblock: or_gate 3 4 open 6 open # out2 depends on in4 and in5
 # and is not registered.

1.3. Detailed Routing Architecture Description

The following information is only required to be in the architecture description file if combined global/detailed routing is to be performed. Note that currently combined global/detailed routing is possible only when all channels have been specified to have the same width.

switch_block_type {subset | wilton | universal}: All the switch blocks have $F_s = 3$. That is, whenever horizontal and vertical channels intersect, each wire segment can connect to three other wire segments. The exact topology of which wire segment connects to which can be one of three choices.

- *subset:* The switch box used in the AMDREL fine-grain architecture -- a wire segment in track 0 can only connect to other wire segments in track 0 and so on (Fig. 8a).
- *wilton:* The switch box illustrated in Fig. 8b.
- *universal:* The switch box illustrated in Fig. 8c. To see the topology of a switch box, simply hit the "Toggle RR" button when a completed routing is on screen in VPR. In general the wilton switch box is the best of these three topologies and leads to the most routable FPGAs.



Fig. 2: Example logic block where many pins are not logically equivalent.





b)



Fig. 3: Switch box types

Fc_type {absolute | fractional}: Indicates whether the three F_c values (see below) should be interpreted as the number of tracks to which each pin connects (*absolute*), or the fraction of tracks in a channel to which each pin connects (*fractional*).

Fc_input <float>: Sets the number of tracks to which each logic block input pin connects in each channel bordering the pin. The F_c value used is always the minimum of the specified F_c and the channel width, W, so you can set F_c to be huge if you want F_c to always be W.

Fc_output <float>: Sets the number of tracks to which each logic block output pin connects in each channel bordering the pin.

Fc_pad <float>: Sets the number of tracks to which each I/O pad connects in the channel bordering the pad.

segment frequency: <float> length: <int | longline> wire_switch: <int>
opin_switch: <int>

Frac_cb: <float> Frac_sb: <float> Rmetal: <float> Cmetal: <float>

Describes a type of segment. You can specify as many types of segments as you like -- just use one segment line for each. The meaning of each value is:

• *frequency:* The fraction (from 0 to 1) of routing tracks composed of this type of segment. The sum of the frequency values for all the segment lines must be 1 -- i. e. 100% of the tracks have been described.

- *length:* Either the number of logic blocks spanned by each segment, or the keyword *longline*. Longline means segments of this type span the entire FPGA array.
- *wire_switch:* The index of the switch type used by other wiring segments to drive this type of segment. That is, switches going *to* this segment from other pieces of wiring will use this type of switch.
- *opin_switch:* The index of the switch type used by CLB and pad output pins to drive this type of segment.
- *Frac_cb:* Describes the internal population of the segment for connection boxes (connections to logic blocks). This number gives the fraction (from 0 to 1) of logic blocks passed by this segment to which it will have a connection box. A switch exists from a segment to a logic block pin only if (1) the segment wants a connection box to that logic block and (2) the logic block connection box pattern for that pin wants a connection to that segment.
- *Frac_sb:* Describes the internal population of the segment for switch boxes (connections to other routing tracks). This number gives the fraction (from 0 to 1) of the length + 1 switch blocks which could exist along the segment that do in fact exist. So, a segment of length 9 that had a Frac_sb value of 0.5 would have 5 switch boxes along its length. Exactly which tracks a segment connects to at each switch box is determined by the *switch_box_type* parameter.
- *Rmetal:* Resistance per unit length (in terms of logic blocks) of this wiring track, in Ohms. For example, a segment of length 5 with Rmetal = 10 Ohms / logic block would have an end-to-end resistance of 50 Ohms.
- *Cmetal:* Capacitance per unit length (in terms of logic blocks) of this wiring track, in Farads. For example, a segment of length 5 with Cmetal = 2e-14 F / logic block would have a total metal capacitance of 10e-13F. For example, let's say an architecture file describes two types of segments.

segment frequency: 0.5 length: 2 wire_switch: 0 opin_switch: 0 Frac_cb: 1. \ Frac_sb: 0.666 Rmetal: 5 Cmetal: 5e-15 segment frequency: 0.5 length: 4 wire switch: 0 opin switch: 0 Frac cb: 0.5 \

Frac sb: 1. Rmetal: 3 Cmetal: 2e-15

If the FPGA you wish to route has a channel width of 4, one channel will look as shown in Fig. 4. Notice that 2 tracks (50% of the tracks) are segments of length 2, and 2 tracks are segments of length four.

Also notice that the number of switch boxes and connection boxes along each segment has been set in accordance with the Frac_sb and Frac_cb values for each segment type.

switch <int> buffered: { yes | no } R: <float> Cin: <float> Cout: <float> Tdel: <float>

Describes a a type of switch. This statement defines what a certain type of switch is -- segment statements refer to a switch types by their number (the number right after the switch keyword). The various values are:

• *buffered:* yes, if this switch is a tri-state buffer, no if this switch is a pass transistor.

- *R:* resistance of the switch.
- Cin: Input capacitance of the switch.
- Cout: Output capacitance of the switch.

• *Tdel:* Intrinsic delay through the switch. If this switch was driven by a zero resistance source, and

drove a zero capacitance load, its delay would be $T_{del} + R * C_{out}$.

R_minW_nmos <float>: The resistance of minimum-width nmos transistor. This data is used only

by the area model built into VPR.

R_minW_pmos <float>: The resistance of minimum-width pmos transistor. This data is used only

by the area model built into VPR.

The lines below give an example of a detailed routing description from a .arch file.



Fig. 4: Example of a segmented routing channel with four tracks per channel.

switch_block_type planar #Uses the fewest switches on a segmented architecture. Fc_type fractional #Fc values below are in terms of fraction of W.

Fc_output 1. #clb output pins connect to all W tracks in adjacent channels

(if each of those tracks wants a connection box there).

Fc_input 0.5 #clb input pins connect to half (0.5 * W) of adjacent tracks

(if each of those tracks wants a connection box there).

Fc_pad 0.7 # I/O pads connect to 70% (0.7 * W) of adjacent tracks

(if each of those tracks wants a connection box there).

50% of segments are length 2, 50% are length 4. Length two segments are driven by

type 1 switches when the connection is coming from another wire, and are driven # by type 0 switches when the connection comes from a clb output pin.

The length four segments are always driven by type 0 switches.

segment frequency: 0.5 length: 2 wire_switch: 1 opin_switch: 0 Frac_cb: 1. \ Frac_sb: 0.666 Rmetal: 5 Cmetal: 5e-15

segment frequency: 0.5 length: 4 Frac_cb: 0.5 Frac_sb: 1. Rmetal: 3 \ Cmetal: 2e-15

In this case, type 1 switches are pass transistors, while type 0 switches are # tri-state buffers.

switch 1 buffered: no R: 100 Cin: 2e-15 Cout: 2e-15 Tdel: 0 # Pass transistor switch 0 buffered: yes R: 50 Cin: .5e-15 Cout: 4e-15 Tdel: 1e-11 # Tri-state buffer

 R_minW_nmos 100 #Used by area model. Min-width transistor resistances. R_minW_pmos 200

1.4. Timing Analysis Parameters

The following parameters are required if timing analysis is to be performed on the placed and routed circuit, or the timing-driven router is to be used.

C_ipin_cblock <float>: Input capacitance of the buffer isolating a routing track from the connection boxes (multiplexers) that select the signal to be connected to a logic block input pin (Fig. 5). One of these buffers is inserted in the FPGA for each track at each location at which it connects to a connection box. For example, a routing segment that spans three logic blocks, and connects to logic blocks at two of these three possible locations would have two isolation buffers attached to it. If a routing track connects to the logic blocks both above and below it at some point, only one isolation buffer is inserted at that point. If your connection from routing track to connection block does not include a buffer, set this parameter to the capacitive loading a track would see at each point where it connects to a logic blocks.

T_ipin_cblock <float>: Delay to go from a routing track, through the isolation buffer (if your architecture contains these) and a connection block (typically a multiplexer) to a logic block input pin.

T_ipad <float>: Delay through an input pad.

T_opad <float>: Delay through an output pad.

T_clb_ipin_to_sblk_ipin <float>: Delay from an input pin of a CLB (logic block) to an input pin of a subblock within that CLB. For architectures without local routing (i.e. CLB input pins connect directly to some logic element, like a LUT or multiplexer) this delay is essentially zero.

T_sblk_opin_to_sblk_ipin <float>: Delay from the output of a subblock to the input of another subblock within the same CLB. For architectures without local routing (e.g. the output of one subblock is hardwired to the input of another) this delay is essentially zero.

T_sblk_opin_to_clb_opin <float>: Delay from the output of a subblock to a CLB(logic block) output pin. For architectures without local routing (e.g. the output of a LUT is hardwired to each logic block output), this delay is essentially zero.

T_subblock T_comb: <float> T_seq_in: <float> T_seq_out: <float>

Describes the delays within a subblock. There must be one T_subblock line for each subblock a logic block can contain -- i.e. there must be subblocks_per_clb of these T_subblock lines. The first line specifies the delays of subblock zero, which is the first subblock listed in each CLB in the circuit netlist file. The second T_subblock line specifies the delay of subblock one, (the second subblock in each CLB in the circuit netlist file), and so on. If the subblocks within a CLB have different delays then, you must list them in the same order in the architecture and netlist files.

The above delays are illustrated in Fig. 6.



Fig. 5: Routing track to logic block connection structure.



Fig. 6: Local routing delays within a logic block (CLB).

• *T_comb:* The delay from any subblock input to the subblock output when this subblock is used in combinational mode. A subblock is used in combinational mode when the netlist leaves its clock pin OPEN.

• *T_seq_in:* The delay from any subblock input pin to the FF storage element when this subblock is used in sequential mode. A subblock is used in sequential mode when the netlist hooks its clock pin to some signal. If this subblock was a simple flip flop, for example, then T_seq_in is the setup time. If this subblock corresponds to, say, a LUT feeding into a flip flop, then T_seq_in should be set to the LUT delay plus the setup time.

• *T_seq_out:* The delay from the subblock storage element (FF) to the subblock output pin when this block is used in sequential mode. A subblock is used in sequential mode when the netlist hooks its clock pin to some signal. If this

subblock had a flip flop hooked to its output pin, for example, then T_seq_out would be the clock-to-Q delay of the flip flop.

2. Running DUTYS

As with all tools in the AMDREL tool flow, DUTYS can be executed both from the command line and using the GUI.

2.1. Running DUTYS from the command line

DUTYS can be run from the command line by typing *dutys* at the command line. The tool then proceeds by asking the user to supply the name of the desired architecture file. After the name of the new architecture file is typed, the tool proceeds by asking the user to input the parameters mentioned in the previous section one by one.

2.2. DUTYS GUI

The DUTYS GUI (part from the MEANDER design framework) can be seen in the Fig. 7. The GUI has a number of fields that need to be filled and each one corresponds to the questions asked in the command line execution. Next to each field a short description of the corresponding architecture file parameter is included. The GUI is available from the URL http://vlsi.ee.duth.gr/amdrel.



Fig. 7: DUTYS GUI