



Course and Laboratory on Electronic Design Automation

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FSM Design for Digital Signal Processing Using VHDL

Prof. Dr. Martin J. W. Schubert Electronics Laboratory Regensburg University of Applied Sciences Regensburg **Abstract.** This tutorial is intended to teach digital finite state machine (FSM) design using VHDL with emphasis on digital signal processing (DSP). Some VHDL knowledge is presumed. Understanding of the theoretical DSP background is helpful but not required.

1 Introduction

The goal of this exercise is to teach finite-state-machine (FSM) design using VHDL.

Nowadays, synthesizable description of digital hardware is either done with VHDL [1], Verilog [2]. Higher level design tools like SystemC/C++ [3] or HDL coder [4] typically generate optionally VHDL or Verilog code for further synthesis.

Required software tools

VHDL simulator ModelSim simulator [5] and VHDL synthesizer Quartus II 8.1 [6] are available at the CIP pools of Regensburg Univ. of Appl. Sciences or downloadable free from Altera [6]. The Altera Edition (ModelSim AE) is by a factor 2...3 slower than the paid versions. The ModelSim and Quartus II software used within this tutorial can be obtained free from Altera [6] and other sources.

The particular tools used here are most probably no more available from Altera [6] and their latest versions of are significantly larger without any advantage for this tutorial. Therefore, ModelSim and Quartus tools used here are available for members of Regensburg Univ. of Appl. Sciences also at the internal network, drive k:\[15].

Required knowledge

For this tutorial some VHDL knowledge is presumed. It can be obtained or deepened e.g. with [7], [8]. Understanding of the theoretical background of DSP with respect to digital filter design is not required but helpful an can be obtained from [9]. This script is accompanied with the scripts "FSM Design for DSP Using Fixed-Point Numbers" [10] and "FSM Design for DSP Using Matlab" [11].

The organization of this document is as follows:

- Chapter 2 summarizes some fundamental FSM design rules and
- Chapter 3 some Fundamental Rules for Synchronous Design.
- Chapter 4 simulates a simple VHDL counter model using ModelSim.
- Chapter 5 uses Quartus II to synthesizes the VHDL counter model and download it to the DE2 board [12] with Altera Cyclone II FPGA.
- Chapter 6 is a short answer to some frequently asked questions on the VHDL statements used in this course. Significantly more details can be obtained from [7], [8].
- Chapter 7 draws relevant conclusion and
- Chapter 8 offers references.

2 Some Fundamental FSM Desing Rules2.1 The FSM Model



Fig. 2.1-1: Finite-State Machine model. (Signal *enable* depends in the implementation.)

Moore outputs depend on the state vector only. **Mealy** outputs can change value independently of the clock (difficult to test).

The **state memory** is typically made of n D-flipflops. It can represent a finite number of 2^n states. An infinite state memory is for example a capacitor, because it can memorize an infinite number of voltages.

Fig. 2.1-2 compares the different ways to draw a FSM.



Fig. 2.1-2: Finite-State Machine models (a) Loop model, (b) combinational NextState logic described as table, (c) bubble model.

2.2 An Unequivocal Logic Representation for Synthesis



Fig. 2.2: Process flow of synthesis of digital circuits.

A circuit can be described in many ways. For synthesis we need an unequivocal description as the bit-map shown in table 2.2 for combinational logic. Challenging is hierarchical design including blocks as e.g. multipliers as "black boxes" into the optimization process.

- 1. Distinguish combinational from sequential logic.
- 2. Sequential logic is typically realized with D-flipflops.
- 3. Combinational logic: no memory, no tri-state outputs; unequivocal description: bitmap.
- 4. Synthesize logic with respect to target technology and goals

Table 2.2: A bit map unequivocally describes the static behavior of combinational logic

Comb	inatio	nal L	ogic I	nputs		Combi	inatio	nal Lo	ogic C	Dutputs
<u>i1</u>	i2	i3	i4	•••	ix	o1	o2	03	•••	oy
0	0	0	0		0	1	0	1		0
1	0	0	0		0	0	1	1		1
1	1	0	0		0	1	0	1		1
1	1	1	1		1	1	1	0		1

2.3 Some Fundamental Rules for Synchronous Design

Experience has shown that the safest method for the design and test of digital circuits and systems is synchronous circuit design.

A design is synchronous, if

- all data storage elements are clocked and in normal operation change state only in response to the active edge of the clock signal,
- the same active edge of the clock signal is applied at precisely the same time point at every clocked cell in the design.

Not recommended: Non-synchronous changes of flipflop states:





Do not gate clocks !



Fig. 2.3-2: (a) Bad enable-flipflop (eff) with gated clock, (b) good eff, (c) transistor-level eff.

Asynchronous designs are typically faster and less expensive but difficult to test. Use enableflipflops optimized on transistor-level when available. The maximum speed of the ripple counter below does not depend on its length, which is a significant advantage over the synchronous counter. However, synchronous circuits are easier to test.



Fig. 2.3-3: (a) asynchronous 4-bit ripple counter, (b) synchronous 4-bit counter.

A toggle flipflop (tff) and an enabled toggle flipflop (etff) uses d=NOT(q).

3 VHDL Design for Synthesis

3.1 Fundamental Guidelines for Synthesizable VHDL Code

- Synthesizable code is typically termed **rtl** (register transfer level) Code.
- Non-synthesizable code is typically termed **behavioral** (testbenches).

3.1.1 Do not Reinvent the Wheel!

According to Wikipedia [13] "*Reinventing the wheel* is a phrase that means to duplicate a basic method that has already previously been created or optimized by others" which normally doesn't make sense. In most situations a simple rule delivers the best results:

Describe your circuit as detailed as necessary and as general as possible.

- What is synthesizable or not depends on your libraries. An adder or multiplier from your library is in the majority of cases significantly better than a self-made one.
- Do not design primitives at gate level! Particularly flipflops (FFs) are optimized by the technology foundry. Use behavioral FF descriptions as detailed below.
- For information on synthesis exceeding the recommendations below the reader is referred the literature, in particular to the work of Keating and Bricaud [14], which results from a joint design-for-reuse initiative of Cadence Systems and Mentor Graphics.

3.1.2 Design for Portability and Reusability

Preserving portability of VHDL models requires to design without taking advantage of features offered by particular soft- and hardware manufacturers.

ASYNCHRONOUS RESET: Most FPGAs automatically reset all FFs when they are loaded. Therefore, some FPGA designers code FFs without reset. However, it is up-to date practice to test a design in a FPGA before copying the models into an ASIC. When HDL models are ported to an ASIC, memory without reset will have a problem, as FFs fall randomly into '0' or '1' state.

Design for reusability is an important topic [14] that would exceed the scope of this tutorial.

3.1.3 Designing Edge-Triggered Memory

Among the many possibilities of modeling sequential logic VHDL synthesizers are typically sensitive to the following constructs

```
a) ps_dff:PROCESS(clk,reset)
                                           b) ps_state_memory:PROCESS(clk,reset)
   BEGIN
                                              BEGIN
     IF reset='0' THEN
                                                IF reset = '0' THEN
       q<='0' [AFTER delay];</pre>
                                                  state<=reset_state [AFTER delay];</pre>
     ELSIF clk'EVENT AND clk='1' THEN
                                                ELSIF clk'EVENT AND clk='1' THEN
       q<=d [AFTER delay];</pre>
                                                  state<=next_state [AFTER delay];</pre>
     END IF;
                                                END IF;
   END PROCESS ps_dff;
                                              END PROCESS ps_state_memory;
```

Listing 3.1.3-1: Modeling memory: (a) for a scalar state and (b) for a state vector.

The AFTER clause is optional, improves simulation and is ignored by the synthesizer. For FFs with inverting output a model as illustrated in Listing 3.1.3-2(a) is recommended, because it guarantees qb=NOT(q) in all situations.

Not recommended are constructs assigning both q and qb within the PROCESS statement as shown in Listing 3.1.3-2(b). During initialization all processes run through and the forbidden situation q=qb='0' or q=qb='1' may occur, when there is neither a rising clock edge nor reset='0'.

a)	ps_dff_good:PROCESS(clk,reset)	<pre>b) ps_dff_err:PROCESS(clk,reset)</pre>
,	BEGIN	BEGIN
	IF reset='0' THEN	IF reset ='0' THEN
	q<='0' AFTER delay;	q<='0' AFTER delay;
	ELSIF clk'EVENT AND clk='1' THEN	qb<='1' AFTER delay;
	q<=d AFTER delay;	ELSIF clk'EVENT AND clk='1' THEN
	END IF;	q<=d AFTER delay;
	END PROCESS ps_dff_good;	qb<=NOT(d) AFTER delay;
	qb<=NOT(q); this is always o.k.	END IF;
		END PROCESS ps_dff_err;

Listing 3.1.3-2: FF model with inverting output: (a) recommended (b) not recommended.

Many silicon foundries offer optimized enable flipflops (EFFs) and the construction shown in Fig. 2.3-2 can be realized as:

```
ps_state_memory:PROCESS(clk,reset)
BEGIN

IF reset ='0' THEN
state<=reset_state;
ELSIF clk'EVENT AND clk='1' AND enable='1' THEN
state<=next_state;
END IF;
END FROCESS ps_state_memory;</pre>
```

Listing 3.1.3-3: Design of an enable flipflop (EFF).

3.1.4 Designing Combinational Logic

Modeling combinational is well done with concurrent statements. Example for a 1-bit half adder:

sum <= a xor b; carry_out <= a AND b;</pre>

In many situations designers use a **PROCESS** statement of the form

```
ps_combinational:PROCESS(sensitive_signals)
BEGIN
driven_output_signals <= f(sensitive_signals, other_signals);
END PROCESS ps_combinational;</pre>
```

The process generates (forbidden) latches, if one of the following two basic rules for combinational logic is violated:

- 1. All input signals of the process must be listed in its sensitivity list.
- 2. All output signals must be driven all time.

Violation example for rule 1:

In the process below labeled $ps_partial_sensitivity_list$ the assignment y <=b; can be executed if and only if an event on signal a occurs. To realize that a double-edge triggered FF using signal a as clock signal is required, which is not combinational.

```
ps_partial_sensitivity_list:PROCESS(a)
BEGIN
y<=b; -- latch generation: b is not in the sesitivity list!
END PROCESS ps_partial_sensitivity_list;</pre>
```

Violation example for rule 2:

In the process below labeled ps_wrong the assignment $next_y <=b$; can be executed if and only if a='1'. When a='0' then $next_y$ must not change. To guarantee this the synthesizer generates a latch, so that the code is no more purely combinational. This problem is removed in process ps_good . Don't forget that y has now to appear in the sensitivity list because it is an input signal.

```
ps_wrong:PROCESS(a,b)
                                              ps_good:PROCESS(a,b,y)
a)
                                         b)
     BEGIN
                                               BEGIN
       IF a='1' THEN
                                                 IF a='1' THEN
         next_y <= b;</pre>
                                                   next_y <= b;</pre>
       -- what happens if a='0'?
                                                 ELSE
       END IF;
                                                   next_y <= y;</pre>
     END PROCESS ps_wrong;
                                                 END IF;
                                               END PROCESS ps_good;
```

CASE | **SELECT** versus IF | "<= ... WHEN" statement. VHDL synthesizer manuals recommend to prefer CASE to IF statements wherever possible. This is because CASE uses one conditions for the evaluation of all situations as symbolized in Fig. 3.3.1(a) while the priority bit coding of the IF statement evaluates a new condition for every branch. The concurrent counterparts are SELECTed and conditional signal assignment.



Fig. 3.1.4: Conditional VHDL assignments (a) using the sequential CASE <condition> IS or concurrent WITH <condition> SELECT statement, (b) using the sequential IF <condition> THEN or concurrent <signal> <= <expression> WHEN<condition> statement

End conditional statements with an OTHERS or ELSE clause. Be aware of state values other than '0' and '1', e.g. logic 'X'. Even if it is sure at the moment that all cases are covered by a CASE statement, the designer should keep in mind that the code may be maintained or reused by other persons. (Exceptions to this rule hold for FSM design using ONE process for next-state logic and memory -> see below.)

Examples: Let us assume the following VHDL signal declaration: SIGNAL condition, a, b, y:std_logic ;

Selected signal assignment:

Concurrent	Sequential
	<pre>p_case: PROCESS(condition,a,b) BEGIN</pre>
WITH condition SELECT y <= a WHEN '0', b WHEN '1', 'X' WHEN OTHERS;	CASE condition IS WHEN '0' => y<=a; WHEN '1' => y<=b; WHEN OTHERS => y<='X'; END CASE END PROCESS p_case;

Conditional signal assignment:

Concurrent:	Sequential
y <= a WHEN condition='0' ELSE b WHEN condition='1' ELSE 'X';	<pre>p_if: PROCESS(condition,a,b) BEGIN IF condition='0' THEN y<=a; ELSIF condition='1' THEN y<=b; ELSE</pre>

3.2 VHDL-Design of Synchronous Finite State Machines

3.2.1 General Guidelines for Event-Driven FSM Design



Fig. 3.2.1-1: Clocked state machine structure: Mealy/Moore machine. Moore outputs depend only on *state* and not on *stimuli*, the latter may be asynchronous. Signal *enable* depends on the implementation (can be avoided by not changing *NextState*).

Avoid Mealy outputs if possible. Mealy outputs depend directly on input signals and are difficult to test because they are subject to change any time. The delay of a chain of modules with mealy outputs is difficult to predict.

There are two possibilities to code the feedback-loop of a finite state machine:

- 1. Using a single process for both next-state logic and memory.
- 2. Using two processes: one for the next-state logic and an other for the memory.

Example counter: We will now build the counter known from the Matlab model above.

Fig. 3.2.1-2: The counter entity corresponds to the schematics symbol. The prefix 'c' denotes cPeriod as a constant (Constants are passed to a VHDL entity as "generics".)



A bitwidth cM can be declared directly using e.g. std_logic_vector(cM-1 DOWNTO 0) or INTEGER RANGE 0 TO 2**cM, or indirectly using e.g. INTEGER RANGE 0 TO 2**cM. In this case the bitwidth can be computed from the C-code

cM=(int)ceil(log(cMax)/log(2)), where the logarithms dualis can be computed from $ld(x)=log_a(x)/log_a(2)$ for any base a.

Entity: Assume the following Entity for the counter that counts from 0...cPeriod-1.

```
ENTITY counter IS
   GENERIC(cPeriod:POSITIVE:=10);
   PORT(reset,clock,enable:IN std_logic;
        maxflg:BUFFER std_logic;
        count:BUFFER NATURAL RANGE 0 TO cPeriod-1
      );
END ENTITY counter;
```

There are two design techniques to model the feedback-loop a FSM: Two PROCESS statements for next-state logic and memory or a single PROCESS statement for both.

3.2.2 **Two PROCESS Statements for Next-State Logic and State Memory**

```
ARCHITECTURE rtl_counter2 OF counter IS
  SIGNAL NextCount:NATURAL RANGE 0 TO cPeriod-1;
BEGIN
  -- Begin NextState Logic:
  p ns:PROCESS(enable,count) BEGIN
    IF enable='1' THEN
      IF count = cPeriod-1 THEN
        NextCount<=0;</pre>
      ELSE
        NextCount<=count+1;</pre>
      END IF;
    ELSE
      -- what happens at enable ='0' ?!
      NextCount<=count;</pre>
    END IF;
  END PROCESS; -- End NextState Logic
  -- Begin State Memory
  p_mem:PROCESS(reset,clock)
  BEGIN
    IF reset='0' THEN
      count <= 0;
    ELSIF clock'EVENT AND clock='1' THEN
      count <= NextCount;</pre>
    END IF;
  END PROCESS p_mem;
  _ _
  -- output logic:
  maxflg <= '1' WHEN count=cPeriod-1 ELSE '0';</pre>
END ARCHITECTURE rtl_counter2;
CONFIGURATION con_counter2 OF counter IS
  FOR rtl_counter2
  END FOR;
END CONFIGURATION con_counter2;
```

The above example reflects the state machine composed of next-state logic and memory. When choosing this solution, you should keep in mind the 2 rules for combinational logic:

- 1. Have all input signals in the sensitivity list of the combinational-logic process and
- 2. Drive every output-bit of the next-state logic any time in any situation. (Use ELSE and OTHERS keywords in IF and CASE statements, respectively!)

One method to guarantee point two above is to begin the combinational next-state process with the following code line:

```
State <= NextState;</pre>
```

Within a process the last of several signal assignments is valid. So we have driven any bit of the state vector and may now drive some of its bits if desired.

3.2.3 One **PROCESS** Statement for Next-State Logic and State Memory

```
ARCHITECTURE rtl_counter1 OF counter IS
  SIGNAL NextCount:NATURAL RANGE 0 TO cPeriod-1;
BEGIN
  -- Counter
  p_fsm:PROCESS(reset,clock)
  BEGIN
    IF reset='0' THEN
      count <= 0;
    ELSIF clock'EVENT AND clock='1' AND enable='1' THEN
      IF count = cPeriod-1 THEN -- Begin NextState Logic
        count <= 0;</pre>
      ELSE
        count <= count+1;</pre>
                                  -- End NextState Logic
      END IF;
    END IF;
  END PROCESS p fsm;
  -- output logic:
  maxflg <= '1' WHEN count=cPeriod-1 ELSE '0';</pre>
END ARCHITECTURE rtl_counter1;
```

This realization uses the same memory model as the example above, but replaces the next-state assignment count<=NextCount; by the combinational logic. Advantages:

- There are no other signals in the sensitivity list than reset and clock,
- There is no ELSE or OTHERS clause to define what happens when enable='0',
- Such code is often easier to read, shorter and less error prone to maintain.

To obtain the solution in example 2 from the solution in example 1 ...

- 1. Replace the statement count<=NextCount; (or state<=NextState;) inside the memory by the next-state logic.
- 2. Rename NextCount to count (or NextState to state) and remove the signal declaration statement for the NextCount (or NextState) vector.
- 3. Remove the process for the NextState logic.

4 Simulate the VHDL Model

4.1 VHDL Testbench

```
LIBRARY ieee; USE ieee.std_logic_1164.ALL;
ENTITY tb_counter IS END ENTITY tb_counter;
LIBRARY ieee;
USE ieee.std_logic_1164.ALL;
ARCHITECTURE beh_tb_counter OF tb_counter IS
  CONSTANT cPeriod:NATURAL:=10;
  SIGNAL count:NATURAL RANGE 0 TO cPeriod-1;
  SIGNAL reset,clock,maxflg:std_logic:='0';
  SIGNAL enable:std logic;
  COMPONENT counter IS
    GENERIC(cPeriod:POSITIVE:=10);
    PORT(reset,clock,enable:IN std_logic;
         maxflq:BUFFER std logic;
         count:BUFFER NATURAL RANGE 0 TO cPeriod-1
        );
  END COMPONENT counter;
  CONSTANT fclk:REAL:=50.0E6;
BEGIN
  clock <= NOT clock AFTER sec/(2.0*fclk);</pre>
  reset <= '0', '1' AFTER 12 ns;
  enable <= '1', '0' AFTER 355 ns, '1' AFTER 545 ns;
  i:counter GENERIC MAP(cPeriod)
            PORT MAP(reset,clock,enable,maxflg,count);
END ARCHITECTURE beh_tb_counter;
CONFIGURATION con_tb_counter OF tb_counter IS
  FOR beh_tb_counter
    FOR i:counter USE ENTITY WORK.counter(rtl_counter2); -- config #
    END FOR;
  END FOR;
END CONFIGURATION con_tb_counter;
```

📕 wave - default	
<u>File E</u> dit <u>V</u> iew <u>A</u> dd F <u>o</u> rmat <u>T</u> ools <u>W</u>	ndow
D 6886 18 06 20	MA ฿ ฿ 🕸 🚟 🚑 🕺 ↑ ← 🛶 EF 100 ps 🕈 EL EL EL ?) (?) 🛐 🌇 🛃
<u>│┟┎┍┰╶╶╶</u> ┊│	
Messages	
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0 ps to 1050 ns No	v: 1 us Delta: 1

Fig. 4.1: Simulation of the code above obtained with the ModelSim simulator

Watch the VHDL code line above with comment "config #". It is configured to use architecture rtl_counter2 for entity counter. If this explicit configuration is omitted the default configuration is: Use the last compiled architecture for entity counter.

4.2 Using the ModelSim Simulator

The ModelSim and Quartus II installation files can be started from the CIP-Pools in the HS.R selecting Start \rightarrow Fachbereiche \rightarrow Elektrotechnik \rightarrow Altera \rightarrow {ModelSim | Quartus II}.

The installation files

- 81_modelsim_6.3g_p1_ae_free.exe
- 81_quartus_free.exe

and can also be obtained for self-installation from drive k:\Sb\Software\Altera\...

```
Start the ModelSim simulator at Regensburg Univ. of Appl. Sciences (HS.R):
File -> Change Directory -> ...\counter -> ok
Tools -> Tcl -> Execute Macro -> work.do
```

The file *work.do* Tcl-contains the commands illustrated in the listing below. You can also type these commands into the transcript window.

vlib work	# Create working library named work
vmap work work	# Map logical name work to working lib. work
vcom -work work counter.vhd	# Compile file <i>counter.vhd</i> to working lib. <i>work</i>
vcom -work work rtl_ounter1.vhd	# Compile file <i>counter1_rtl.vhd</i> to lib. <i>work</i>
vcom -work work rtl_counter2.vhd	# Compile file <i>counter2_rtl.vhd</i> to lib. <i>work</i>
vcom -work work counter_tb.vhd	# Compile file <i>counter_tb.vhd</i> to lib. <i>work</i>
vsim work.con_counter_tb	# Simul. config. <i>con_counter_tb</i> located in <i>work</i>
do wave.do	# Run file wave.do defining the wave-win
run 1 us	# Simulate a time span of 1 µs

Library creation and usage in ModelSim:

```
vlib <physical_name> # Create physical library <physical_name> 
wmap <logical_name> <physical_name> # Map <logical_name> to <physical_name> 
# Map <logical_name> to <physical_name> 
# '.' stands for this and '..' for parent directory: 
wmap adac_lib ../adac_bin 
wmap adac_lib ../adac_bin 
wcom -work adac_lib ../filter.vhd 
# '.' stands for this and '..' for parent directory: 
# Create physical library ../adac_bin 
# Map logical name adac_lib to lib. adac_bin 
# Create physical library ../adac_bin 
# Map logical name adac_lib to lib. adac_bin 
# Compile filter.vhd to library adac_lib
```

Listing 4.2: How to create a physical library and a logical name for it within ModelSim

Exercise: Writing VHDL Code

Take the architecture listings rtl_counter# (#=1,2) above. Delete some of the code lines describing the FSM as shown below. Rewrite the model by yourself and test it.

ARCHITECTURE rtl_counter2 OF counter IS SIGNAL NextCount:NATURAL RANGE 0 TO cPeriod-1; BEGIN -- Begin NextState Logic: p_ns:PROCESS(....) BEGIN END PROCESS; -- End NextState Logic -- Begin State Memory BEGIN IF reset='0' THEN count <= 0; ELSIF clock'EVENT AND clock='1' THEN END IF; END PROCESS p_mem; -- output logic: maxflg <= '1' WHEN count=cPeriod-1 ELSE '0';</pre> END ARCHITECTURE rtl_counter2;

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```
-- end of encapsulated combinational Next-State logic
END IF;
END PROCESS p_fsm;
--
-- output logic:
maxflg <= '1' WHEN count=cPeriod-1 ELSE '0';
END ARCHITECTURE rtl_counter1;
```

Hint: You will find the solution in "ARCHITECTURE rtl_counter# OF counter" above (#=1,2).

5 Synthesize the VHDL Model



VHDL Testbench for Synthesis & Download into the DE2 FPGA Board

Fig. 5.1: The Terasic DE2 board with Altera FPGA, copied from the DE2 User Manual [12].

Using Altera's Quartus II [6] to perform the synthesis



(a) Interface compent *de2_counter* (b) Synthesized component *counter*

Fig. 5.2: Quartus II 8.1 RTL Viewer: synthesized models

Assignments between the port signals of the entity and corresponding on-board wires are defined within the file *DE2_pin_assignments.csv* supplied by Altera and can be

modified there. It is imported to Quartus2 using Assignments \rightarrow Import Assignments... \rightarrow Filename. Short summary of the most important signals.

Listing 5: VHDL Testbench for the counter module

```
-- For Board: Altera DE2 with FPGA cyclone II EP2C35F672C6
LIBRARY ieee;
USE ieee.std_logic_1164.ALL, ieee.std_logic_signed.ALL;
ENTITY de2_counter IS
    PORT(CLOCK_50,CLOCK_27:IN std_logic;
      key:IN std_logic_vector(3 DOWNTO 0); -- low when pressed
       sw:IN std_logic_vector(17 DOWNTO 0); -- low when pulled down
       ledg:BUFFER std_logic_vector(8 DOWNTO 0); -- high active
       ledr: BUFFER std logic vector(17 DOWNTO 0); -- high active
      hex0,hex1,hex2,hex3,hex4,hex5,hex6,hex7:OUT std_logic_vector(0 TO 6);
      gpio_0:BUFFER std_logic_vector(35 DOWNTO 0);
      gpio_1:INOUT std_logic_vector(35 DOWNTO 0)
     );
END ENTITY de2_counter;
ARCHITECTURE rtl_de2_counter OF de2_counter IS
  CONSTANT cPeriod:NATURAL:=2_147_483_647; -- 2**31-1
  CONSTANT cCountWidth:NATURAL:=32;
  SIGNAL count:NATURAL RANGE 0 TO cPeriod-1;
  CONSTANT cZeroVector:std_logic_vector(cCountWidth-1 DOWNTO 0):=(OTHERS=>'0');
  SIGNAL count_vector:std_logic_vector(cCountWidth-1 DOWNTO 0);
  TYPE t_7seg IS ARRAY(0 TO 15) OF std_logic_vector(0 TO 6);
  CONSTANT c7seg:t_7seg:=("1111110", "0110000", "1101101", "1111001",
"0110011", "1011011", "1011111", "1110000", "1111111", "1110011",
"1110111", "0011111", "1001110", "0111101", "1001111", "1000111");
  COMPONENT counter IS
    GENERIC(cPeriod:POSITIVE:=10);
    PORT(reset,clock,enable:IN std_logic;
          maxflq:BUFFER std_logic;
          count:BUFFER NATURAL RANGE 0 TO cPeriod-1
         );
  END COMPONENT counter;
  SIGNAL reset, enable, maxflg:std_logic;
BEGIN
  reset <= key(0);</pre>
  enable <= key(1);</pre>
  ledg(6) <= NOT ledg(6) WHEN clock_50'EVENT AND clock_50='1' AND maxflg='1';</pre>
  i:counter GENERIC MAP(cPeriod) PORT MAP(reset=>key(0),clock=>clock_50
                                    enable=>key(1),count=>count,maxflg=>maxflg);
  count_vector<=cZeroVector+count;
  ledr <= count_vector(cCountWidth-2 DOWNTO cCountWidth-19); -- MSB='0'</pre>
  ledg(8) <= NOT ledg(8) WHEN ledr(11)'EVENT AND ledr(11)='1';</pre>
  ledg(7)<=sw(7); ledg(5 DOWNTO 4)<=sw(5 DOWNTO 4);</pre>
  ledg(3 DOWNTO 0) <= key;</pre>
  -- use concatenated leading '0' as data type is std_logic_signed:
  hex7<=NOT c7seg(CONV_INTEGER('0'&count_vector(31 DOWNTO 28)));</pre>
  hex6<=NOT c7seg(CONV_INTEGER('0'&count_vector(27 DOWNTO 24)));
hex5<=NOT c7seg(CONV_INTEGER('0'&count_vector(23 DOWNTO 20)));</pre>
  hex4<=NOT c7seq(CONV_INTEGER('0'&count_vector(19 DOWNTO 16)));
  hex3<=NOT c7seg(CONV_INTEGER('0'&count_vector(15 DOWNTO 12)));</pre>
  hex2<=NOT c7seg(CONV_INTEGER('0'&count_vector(11 DOWNTO 8)));</pre>
  hex1<=NOT c7seg(CONV_INTEGER('0'&count_vector( 7 DOWNTO 4)));</pre>
  hex0<=NOT c7seg(CONV_INTEGER('0'&count_vector( 3 DOWNTO 0)));</pre>
END ARCHITECTURE rtl_de2_counter;
```

Signal in the PORT of entity de2	Connected to element of DE2 board
CLOCK_50,CLOCK_27:IN std_logic	External 50MHz and 27MHz clocks
<pre>key:IN std_logic_vector(3 DOWNTO 0);</pre>	the 4 push buttons below green diodes
<pre>sw:IN std_logic_vector(17 DOWNTO 0);</pre>	low when pulled down
<pre>ledg:BUFFER std_logic_vector(7 DOWNTO 0);</pre>	Green LEDs, high active
<pre>ledr:BUFFER std_logic_vector(17 DOWNTO 0);</pre>	Red LEDs, high active
hex0,hex1,hex2,hex3,hex4,hex5,hex6,hex7	7-segment displays, low-active

Table 5.1: Mapping of FPGA-internal VHDL signals to external devices on the DE2 board

Table 5.2: Functionality of the testbench de2_counter(rtl_de2_counter):

Signal name	mode	features	functionality	comments
key(0)	IN	'0' when key pressed	reset	resets all components when pressed
key(1)	IN	'0' when key pressed	enable for counter	
key(2)	IN	'0' when key pressed		
key(3)	IN	'0' when key pressed		
sw(17:0)	IN	'1' when pushed to board		
ledg(0)	OUT	LED is on @ '1'	= key(0) $=$ reset	
ledg(1)	OUT	LED is on @ '1'	= key(1) $=$ enable	
ledg(2)	OUT	LED is on @ '1'	= key(2)	
ledg(3)	OUT	LED is on @ '1'	= key(3)	
ledg(4)	OUT	LED is on @ '1'	= sw(4)	
ledg(5)	OUT	LED is on @ '1'	= sw(5)	
ledg(6)	OUT	LED is on @ '1'	NOT ledg(6) @ maxflg	toggles when maxflg='1'
ledg(7)	OUT	LED is on @ '1'	= sw(7)	
ledg(8)	OUT	LED is on @ '1'	toggles with ledr(11)	should flash with ca. 1Hz
ledr(17:0)	OUT	LED is on @ '1'	=count_vector(31:13)	makes toggling msb's visible
hex7	OUT	LEDs on @ '0'	=>count_vector(31:28)	makes counter content visible
hex7	OUT	LEDs on @ '0'	=>count_vector(27:24)	makes counter content visible
hex7	OUT	LEDs on @ '0'	=>count_vector(23:20)	makes counter content visible
hex7	OUT	LEDs on @ '0'	$=>$ count_vector(19:16)	makes counter content visible
hex7	OUT	LEDs on @ '0'	=>count vector(15:12)	makes counter content visible
hex7	OUT	LEDs on @ '0'	=>count_vector(11:08)	makes counter content visible
hex7	OUT	LEDs on @ '0'	$=>$ count_vector(07:04)	makes counter content visible
hex7	OUT	LEDs on @ '0'	$=>$ count_vector(03:00)	makes counter content visible

The Quartus II 8.1 files to save (except from the VHDL files) are

<project_name>.qpf</project_name>	ASCII	Quartus Project file
<project_name>.qsf</project_name>	ASCII	Quartus Specification file
<project_name>.sof</project_name>	Binary	downloadable image

The complete required information to rebuild the project is contained in the qpf and qsf files. (Linking of previously compiled files is not possible with the free version of Quartus II.)

- •.qpf: Double-click left mouse button on the *.qpf file to open Quartus II with the respective project loaded. Use an ASCII editor and look into this file.
- •.qsf: Use an ASCII editor to look into this file: It contains settings as the Signal–Pin assignments and the files are to compile. It is always dangerous to change such files with an ASCII editor, but adding or removing VHDL files to be compiled here might be easier and faster than using the menu "Project → Add/Remove File in Project...".
- •.sof: This final "result" file can be downloaded into the FPGA using the programmer.

6 VHDL

This chapter points out some important issues of VHDL, it is no comprehensive tutorial. It is based on the author's experience of frequently asked questions during circuit design courses with students having some previous knowledge on that topic.

6.1 Design Units

Table 6.1: Building blocks in VHDL

	COMPONENT	LIBRARY	Comments
To the outer world	ENTITY	PACKAGE	Corresponds to a symbol
Realization inside	ARCITECTURE	BPACKAGE BODY	Corresponds to a schematics
Combination	CONFIGURATION		

6.2 Compilation Order Dependence

VHDL is compilation order dependent. It has no linker but only a loader. Sub-modules are loaded immediately during the compilation process. We can organize our VHDL code arbitrarily in different files, but it is important that code which is used by other code is compiled before the code using it. The compilation sequence is

 $\texttt{PACKAGE} \rightarrow \texttt{ENTITY} \rightarrow \texttt{ARCHITECTURE} \rightarrow \texttt{PACKAGE} \ \texttt{BODY}$

The package body may be compiled directly after the package but also as lastly.

6.3 Kinds of Code: Concurrent – Sequential – Structural

VHDL code can be written in three different modes: Concurrent is the default mode. In concurrent code the sequence of statements is irrelevant. Using components is structural and within a PROCESS statement or subprograms code is sequential, i.e. it is processed top-down.

6.4 Data

6.4.1 Data Objects

VHDL handles data by one of the following data objects:

1.SIGNAL:	intero	connection wires
2.VARIABLE:	local	storage of temporary data
3.CONSTANT:	named	constant values

Table 6.4.1 shows in which environment a data object may be declared and where it may be used, i.e. assigned or read.

data object	environment to be declared in	environment to be used in
SIGNAL	concurrent	everywhere
VARIABLE	sequential	sequential
CONSTANT	everywhere	everywhere

Table 6.1: The three data objects and the environments where they may be declared and used.



Fig. 6.4.1: VHDL data objects.

6.4.2 Data Types

Data object have to be declared with a data type. Available are scalar the types INTEGER, REAL, enumerated and BIT and the composite data types such ARRAY and RECORD.

Scheme of data type declaration statements:

TYPE type_name IS type_mark; SUBTYPE subtype_name IS <range_specification> OF type_name;

Examples for type and subtype declarations:

```
TYPE BOOLEAN IS (false,true);
SUBTYPE t_address IS RANGE 0 TO 15 OF INTEGER;
SUBTYPE NATURAL IS RANGE 0 TO INTEGER'HIGH OF INTEGER;
```



Fig. 6.4.2: VHDL data types.

The assignment of a value to a data object is guarded such, that only values of the same data type or a subtype can be assigned to a data object. The user can specifiy own data types. See e.g. [Schubert: VHDL Skript] Chapter 2: Data for more details.

6.5 Libraries and Packages

6.5.1 Using Existing Libraries and Packages

Reinventing the wheel is inefficient. Better reuse existing code if possible. It is typically organized in libraries which are composed of packages. A typical library retrieval is

```
(1) LIBRARY ieee;
```

- (2) USE ieee.std_logic_1164.ALL;
- (3) USE ieee.std_logic_unsigned."+";
- (4) LIBRARY adac_lib;
- (5) USE adac_lib.pk_adac.ALL;
- (6) USE work.pk_mypack.ALL;

The code above has the following meanings:

- Line (1): "Retrieve library with name ieee".
- Line (2): Use ALL from package std_logic_1164 found within library ieee.
- Line (3): Use only the declaration of the "+" operator found in package std_logic_unsigned within library ieee. Applied on bit-vectors as operands it will be synthesized as arithmetic summation and the most significant bit will not be interpreted as sign bit. (To treat the first bit as sign bit use the "+" operator from package std_logic_signed.)
- Line (4): "Retrieve library with name adac_lib". As it is not a standard library it must be introduced to the tool, the respective commands are tool dependent.
- Line (5): Use ALL declarations from package pk_adac found within library adac_lib.
- Line (6): Use ALL declarations from package pk_mypack found within library work. The working library work needs no LIBRARY statement, at is always linked.

If the LIBRARY / USE statements are written above ...

- an ENTITY, then they are valid for this entity and its architectures,
- an ARCHITECTURE, then they are valid for this architecture,
- a PACKAGE, then they are valid for this package and its package body,
- a PACKAGE BODY, then they are valid for this PACKAGE BODY.

6.5.2 Creating an Own Package

Package and package body are like entity and architecture. Packages may contain nonexecutable declarations only while package bodies may contain executable code also. Declarations made within a package are available whenever loading this package. Declarations made in a package body are known only in this package body below the declaration. Listing 4.2 shows Tcl-commands to create a library and symbol for it that can be used within the VHDL code when working with the *ModelSim* simulator. M. Schubert

```
LIBRARY ieee; USE ieee.std logic 1164.ALL;
PACKAGE pk_example IS
  -- Declaration of an externally visible constant:
  CONSTANT cExtern: INTEGER:=10;
  -- gate delay as "deferred" constant: no value assigned:
  CONSTANT delay:TIME;
  -- multiplexer: interface declaration only, no executable code:
  FUNCTION mux(sel:INTEGER;vec:std_logic_vector) RETURN std_logic;
  SIGNAL big array:std logic vector(1 TO 40 000);
END PACKAGE pk_example;
PACKAGE BODY pk_example IS
  -- Declaration of an only internally visible constant:
  CONSTANT cIntern: INTEGER:=20;
  _ _
  -- here the deferred constant has to get its value:
  CONSTANT delay:TIME:=2 ns;
  -- here the function mux has to get its body:
  FUNCTION mux(sel:INTEGER;vec:std_logic_vector) RETURN std_logic IS
  BEGIN
   RETURN vec(sel);
  END FUNCTION mux;
END PACKAGE BODY pk_example;
```

The package named pk_example above declares the constant cExtern. It will be available in the package body and everywhere where this package is declared. This is different from the constant cIntern declared within the package body below. It will be known only within this package body and below its declaration.

The constant delay is a so-called deferred constant, as no value is assigned to it in the package. The assignment is deferred (Latin: carried away) into the package body. As the package body may be compiled as last design unit, different delays (e.g. for fast, typical, slow parameters) can be passed to the design by compiling nothing else than the package body.

The external interface of the multiplexer function mux is declared in the package. In the package body function mux gets its executable body. If the declaration in the package is omitted, this function will be known only in the package body and after its declaration.

Signal big_array will be declared in any design unit where package pk_example is declared. It will be different signals which are independent from each other in any of those design units.

6.6 Arrays for State-Machine Design



(b) Data structures to realize the vectors of coefficients (a_i) and data words (s_i, ns_i) .

Fig. 6.6: (a) A model that requires a vector of bit-vectors and (b) a VHDL realization.

Fig. 6.6(a) illustrates a digital filter that requires vectors of bit-vectors. Fig. 6.6 (b) and listing 6.6 realize a single coefficient a_i alias CoefVector(i) as bit-vector of type t_coef. All coefficients are summarized in the constant CoefVector of type t_CoefVector. A single state or next-state is realized as bit-vector of data type t_state. A vector of such signals, e.g. named state or NextState, can be declared as signal of type t_StateVector.

Listing 6.6 (a): The package *pk filter*.

(Note: Some synthesizers cannot assign the values of "deferred constant" **CoefVector** in the package body. In this case remove the package body and assign the values in the package. Doing so you loose the capability of changing filter coefficients by recompiling the package bode only, you have to recompile the entire design.)

```
(1)
    : pk_filter
: Martin Schubert
(2)
   -- Module
(3)
   - -
       Designer
       Date last modified: 07.03.2011
(4)
   - -
               : Data Structures for Digital FIR Filter
(5)
       Purpose
                _____
(6)
    -----
(7)
   LIBRARY ieee; USE ieee.std logic 1164.ALL;
   PACKAGE pk_filter IS
(8)
(9)
     CONSTANT cDataInWidth:POSITIVE:=18; -- Input-Data BitWidth
                                      -- No of Input-Data fract. Bits
(10)
     CONSTANT cDataInFract:POSITIVE:=16;
     CONSTANT cDataOutWidth:POSITIVE:=18; -- Output-Data BitWidth
(11)
(12)
     CONSTANT cDataOutFract: POSITIVE:=16; -- No of Output-Data fract Bits
                                    -- Coefficient's BitWidth
     CONSTANT cCoefWidth:POSITIVE:=18;
(13)
                                       -- No of Coef's fractional Bits
(14)
     CONSTANT cCoefFract:POSITIVE:=18;
     SUBTYPE t DataIn IS std logic vector(cDataInWidth-1 DOWNTO 0);
(15)
     SUBTYPE t DataOut IS std logic_vector(cDataOutWidth-1 DOWNTO 0);
(16)
(17)
     SUBTYPE t coef IS std logic vector(cCoefWidth-1 DOWNTO 0);
(18)
              t CoefVector IS ARRAY (NATURAL RANGE <>) OF t coef;
     TYPE
(19)
     CONSTANT cTaps:POSITIVE:=33;
(20)
     CONSTANT CoefVector: t CoefVector(1 TO cTaps);
(21) END PACKAGE pk filter;
```

Listing 6.6 (b): The package body *pk_filter*.

```
(22) PACKAGE BODY pk_filter IS
(23) CONSTANT CoefVector: t_CoefVector(1 TO cTaps)
(24) := (OTHERS=>(cCoefWidth-6=>'1',OTHERS=>'0'));
(25) END PACKAGE BODY pk_filter;
```

6.7 Exercise: (Solutions at the end of this sub-chapter) $(\Sigma = 16P)$

Complete package pk_filter in listing 6.7.

Make sure that data type std_logic and vectors with elements of this type can be used according to the respective IEEE standard. (2P)

Complete line (10) such, that data type t_DataIn declares a vector representing a number with cDataInWidth bits of type std_logic. (1P)

Complete line (11) such, that data type t_coef declares a vector representing a number with cCoefWidth bits of type std_logic. (1P)

Complete line (12) such, that data type t_CoefVector declares a vector with elements type t_coef. The index range of this vector can be defined later with natural numbers. (2P)

No value is assigned to the constant in line (14). Where does this constant gets its value and what is the correct denomination of such a constant? (2P)

What do you write over an Entity to make all declarations in Package pk_filter available? (2P)

Listing 6.7: Package *pk filter*.

```
(1)
   (2)
   PACKAGE pk_filter IS
(3)
(4)
    CONSTANT cDataInWidth: POSITIVE:=8; -- Input-Data BitWidth
(5)
    CONSTANT cDataInFract:POSITIVE:=6;
                               -- No of Input-Data fract. Bits
    CONSTANT cDataOutWidth:POSITIVE:=18; -- Output-Data BitWidth
(6)
    CONSTANT cDataOutFract:POSITIVE:=16; -- No of Output-Data fract Bits
(7)
    CONSTANT cCoefWidth:POSITIVE:=18;
                              -- Coefficient's BitWidth
(8)
(9)
    CONSTANT cCoefFract:POSITIVE:=18; -- No of Coef's fractional Bits
(10)
    ....TYPE t_DataIn IS .....
(11)
    ...TYPE t coef
                  IS .....
(12)
    ...TYPE t_CoefVector IS .....
(13)
    CONSTANT cTaps:POSITIVE:=33;
(14)
    CONSTANT CoefVector: t_CoefVector(1 TO cTaps);
(15) END PACKAGE pk filter;
```

(**2P**)

Write a library statement at (16), that allows for the multiplication of numbers in the std_logic_vector-format using '*' (while library is known). (1P)
Line (17): Declaration of type t_product as funktion of constants such, that a signal of that type matches the product of t_DataIn and t_coef type signals. (1P)
Line (18): declaration of signal product such, that we can write the VHDL command: "product<=DataIn*CoefVector(i);" (2P)
Lines (20) and (21): Complete the declarations of iPl und iPh such, that (22) works respecting the vector lengths and the number of fractional bits using the respective named

constants in package pk_filter. (For the computation of iPl, iPh see document "FSM

Design for DSP Using Fixed-Point Numbers" [10]).

Solutions: LIBRARY ieee; USE ieee.std logic_1164.ALL; PACKAGE pk_filter IS CONSTANT cDataInWidth:POSITIVE:=8; (2) (3) (4) (5) (6) (7) (8) (9) -- Input-Data BitWidth -- No of Input-Data fract. Bits CONSTANT cDataInFract:POSITIVE:=6; CONSTANT cDataOutWidth:POSITIVE:=18; -- Output-Data BitWidth CONSTANT cDataOutFract:POSITIVE:=16; -- No of Output-Data fract Bits CONSTANT cCoefWidth:POSITIVE:=18; -- Coefficient's BitWidth -- No of Coef's fractional Bits CONSTANT cCoefFract:POSITIVE:=18; (10) (11) **SUB**TYPE t_DataIn IS std_logic_vector(cDataInWidth-1 DOWNTO t_coef IS std_logic_vector(cCoefWidth-1 DOWNTO 0); DOWNTO 0); SUBTYPE (12) t CoefVector IS ARRAY (NATURAL RANGE <>) OF t coef; TYPE CONSTANT cTaps:POSITIVE:=33; (13) (14) (15) CONSTANT CoefVector: t_CoefVector(1 TO cTaps); END PACKAGE pk_filter; (16) USE ieee.std_logic_signed."*"; (17) (18) TYPE t product:std logic vector(cDataInWidth+cCoefWidth-1 DOWNTO 0) SIGNAL product:t_product; SIGNAL DataOut:t DataOut; (19) CONSTANT iP1: INTEGER := cDataInFract + cCoefFract - cDataOutFract; CONSTANT iPh: INTEGER := iP1 + cDataOutWidth - 1; (20)(22)DataOut <= product(iPh DOWNTO iPl);</pre>

6.8 Mixing *INTEGER* and *std_logic_vector* Data Types

6.8.1 The Data Types *INTEGER*, *NATURAL*, *POSITIVE*

After synthesis any integer is a bit-vector. In this subchapter we illustrate bit-to-integer and integer-to-bit conversions.

VHDL requires INTEGER types to span a data range of at least -2 147 483 647 to +2 147 483 647. It is required that INTEGER'LEFT=-INTEGER'RIGHT. Furthermore there are the two predefined subtypes:

SUBTYPE NATURAL IS INTEGER RANGE 0 TO INTEGER'HIGH; SUBTYPE POSITIVE IS INTEGER RANGE 1 TO INTEGER'HIGH;

The *SUBTYPE* declaration passed the properties of *INTEGER* on to *NATURAL* and *POSITIVE*. Example: When operators such as "+", "-" or "*" are declared for *INTEGER* types, they are automatically declared for its subtypes also.

INTEGER is a 4 or 8 byte signed bit-vector, depending on the compiler. For example the following statements might synthesize to a 32- or 64-bit signal i and a 3-bit signal j:

SIGNAL i:INTEGER; SIGNAL j:INTEGER RANGE 0 TO 5;

If it is not sure if *INTEGER* data range is wide enough a bit-vector declaration should be used:

```
USE ieee.std_logic_signed."+";
...
SIGNAL a, b, y:std_logic_vector(127 DOWNTO 0);
...
Y <= a + b;</pre>
```

6.8.2 Synthesizable std_logic_vector – to – INTEGER Conversion

Both packes *std_logic_signed* and *std_logic_unsigned* within library *ieee* declare the following function, interpreting the first bit as sign bit or not, respectively:

```
function CONV_INTEGER(ARG: STD_LOGIC_VECTOR) return INTEGER;
```

Concatenation of a leading '0' bit will always deliver an unsigned interpretation. Example:

```
IntegValue <= CONV_INTEGER('0' & StdLogicVector);</pre>
```

Package std logic arith within library ieee declares function

function CONV_STD_LOGIC_VECTOR(ARG: INTEGER; SIZE: INTEGER) return STD_LOGIC_VECTOR;

A vector's number of elements can be obtained with attribute 'LENGTH, e.g.: vector'LENGTH.

In the following example we convert signal $int_in \rightarrow bits_io \rightarrow int_out \rightarrow bits_out$, where the last conversion to *bits out* serves for comparison with *bits io*.

```
LIBRARY ieee;
USE ieee.std logic 1164.ALL;
USE ieee.std_logic_signed.CONV_INTEGER; -- funktioniert
USE ieee.std_logic_arith.CONV_STD_LOGIC_VECTOR;
ENTITY conv_int_bitvec IS
END conv_int_bitvec;
ARCHITECTURE rtl_conv_int_bitvec OF conv_int_bitvec IS
  SIGNAL bits_io,bits_out:std_logic_vector(7 DOWNTO 0);
  SIGNAL int_in,int_out:INTEGER RANGE -128 TO 127;
BEGIN
  -- conversion integer <-> bitvector
 int in <= 5, -5 AFTER 1 ns, 100 AFTER 2 ns, -100 AFTER 3 ns;
 bits_io <= CONV_STD_LOGIC_VECTOR(int_in,bits_io'LENGTH);</pre>
  int_out <= CONV_INTEGER(bits_io);</pre>
 bits_out <= CONV_STD_LOGIC_VECTOR(int_out,bits_out'LENGTH);</pre>
END rtl_conv_int_bitvec;
```

6.8.3 Synthesizable Multiplication Using Operator '*':

Both packes *std_logic_signed* and *std_logic_unsigned* within library *ieee* declare the operator '*', interpreting the first bit as sign bit or not, respectively:

function "*"(L: STD_LOGIC_VECTOR; R: STD_LOGIC_VECTOR) return STD_LOGIC_VECTOR;

The length of the returned product vector must be the sum of the lengths of the operand vectors.

6.8.4 Synthesizable Addition Using Operator '+'

Both packes *std_logic_signed* and *std_logic_unsigned* within library *ieee* declare several overloadings of operator '+', interpreting the first bit as sign bit or not, respectively.

If both operands are of type *std_logic_vector* the length of the returned sum vector must be the length of the longer operand vector:

function "+"(L: STD_LOGIC_VECTOR; R: STD_LOGIC_VECTOR) return STD_LOGIC_VECTOR;

If there is only one *std_logic_vector* input data type, then the length of the returned sum vector must be the length of the input *std_logic_vector*:

```
function "+"(L: STD_LOGIC_VECTOR; R: INTEGER) return STD_LOGIC_VECTOR;
function "+"(L: INTEGER; R: STD_LOGIC_VECTOR) return STD_LOGIC_VECTOR;
function "+"(L: STD_LOGIC_VECTOR; R: STD_LOGIC) return STD_LOGIC_VECTOR;
function "+"(L: STD_LOGIC; R: STD_LOGIC_VECTOR) return STD_LOGIC_VECTOR;
function "+"(L: STD_LOGIC_VECTOR) return STD_LOGIC_VECTOR;
```

6.8.5 Synthesizable Subtraction Using Operator '-'

Operator '-' is declared much the same as operator '+':

Both packes *std_logic_signed* and *std_logic_unsigned* within library *ieee* declare several overloadings of operator '-', interpreting the first bit as sign bit or not, respectively.

If both operands are of type *std_logic_vector* the length of the returned sum vector must be the length of the longer operand vector.

function "-"(L: STD_LOGIC_VECTOR; R: STD_LOGIC_VECTOR) return STD_LOGIC_VECTOR;

If there is only one *std_logic_vector* input data type, then the length of the returned diffence vector must be the length of the input *std_logic_vector*:

```
function "-"(L: STD_LOGIC_VECTOR; R: INTEGER) return STD_LOGIC_VECTOR;
function "-"(L: INTEGER; R: STD_LOGIC_VECTOR) return STD_LOGIC_VECTOR;
function "-"(L: STD_LOGIC_VECTOR; R: STD_LOGIC) return STD_LOGIC_VECTOR;
function "-"(L: STD_LOGIC; R: STD_LOGIC_VECTOR) return STD_LOGIC_VECTOR;
function "-"(L: STD_LOGIC_VECTOR) return STD_LOGIC_VECTOR;
```

6.8.6 Synthesizable Comparisons

Both packes *std_logic_signed* and *std_logic_unsigned* within library *ieee* declare the several overloadings of comparison operators, interpreting the first bit as sign bit or not, respectively. Declared omparisions are '<', '<=', '>', '>=', '=' standing for less than, less than or equal, greater than, greater than or equal, equal and unequal, respectively. In the packages mentioned above any of that operators has the following three overloadings as as '<':

```
function "<"(L: STD_LOGIC_VECTOR; R: STD_LOGIC_VECTOR) return BOOLEAN;
function "<"(L: STD_LOGIC_VECTOR; R: INTEGER) return BOOLEAN;
function "<"(L: INTEGER; R: STD_LOGIC_VECTOR) return BOOLEAN;</pre>
```

7 Conclusions

This tutorial demonstrates the use of VHDL with the goal to implement a finite state machine (FSM) and data structures suitable to build a finite-impulse-response (FIR) filter.

8 References

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