

UM10120

Chapter 5: Vectored Interrupt Controller (VIC)

Rev. 01 — 24 June 2005 User manual

5.1 Features

- **•** ARM PrimeCell™ Vectored Interrupt Controller
- **•** 32 interrupt request inputs
- **•** 16 vectored IRQ interrupts
- **•** 16 priority levels dynamically assigned to interrupt requests
- **•** Software interrupt generation

5.2 Description

The Vectored Interrupt Controller (VIC) takes 32 interrupt request inputs and programmably assigns them into 3 categories, FIQ, vectored IRQ, and non-vectored IRQ. The programmable assignment scheme means that priorities of interrupts from the various peripherals can be dynamically assigned and adjusted.

Fast Interrupt reQuest (FIQ) requests have the highest priority. If more than one request is assigned to FIQ, the VIC ORs the requests to produce the FIQ signal to the ARM processor. The fastest possible FIQ latency is achieved when only one request is classified as FIQ, because then the FIQ service routine can simply start dealing with that device. But if more than one request is assigned to the FIQ class, the FIQ service routine can read a word from the VIC that identifies which FIQ source(s) is (are) requesting an interrupt.

Vectored IRQs have the middle priority, but only 16 of the 32 requests can be assigned to this category. Any of the 32 requests can be assigned to any of the 16 vectored IRQ slots, among which slot 0 has the highest priority and slot 15 has the lowest.

Non-vectored IRQs have the lowest priority.

The VIC ORs the requests from all the vectored and non-vectored IRQs to produce the IRQ signal to the ARM processor. The IRQ service routine can start by reading a register from the VIC and jumping there. If any of the vectored IRQs are requesting, the VIC provides the address of the highest-priority requesting IRQs service routine, otherwise it provides the address of a default routine that is shared by all the non-vectored IRQs. The default routine can read another VIC register to see what IRQs are active.

All registers in the VIC are word registers. Byte and halfword reads and write are not supported.

Additional information on the Vectored Interrupt Controller is available in the ARM PrimeCell™ Vectored Interrupt Controller (PL190) documentation.

5.3 Register description

The VIC implements the registers shown in [Table](#page-1-0) 33. More detailed descriptions follow.

© Koninklijke Philips Electronics N.V. 2005. All rights reserved.

Table 33: VIC register map

[1] Reset value relects the data stored in used bits only. It does not include reserved bits content.

5.4 VIC registers

The following section describes the VIC registers in the order in which they are used in the VIC logic, from those closest to the interrupt request inputs to those most abstracted for use by software. For most people, this is also the best order to read about the registers when learning the VIC.

5.4.1 Software Interrupt register (VICSoftInt - 0xFFFF F018)

The contents of this register are ORed with the 32 interrupt requests from the various peripherals, before any other logic is applied.

Table 35: Software Interrupt register (VICSoftInt - address 0xFFFF F018) bit description

5.4.2 Software Interrupt Clear register (VICSoftIntClear - 0xFFFF F01C)

This register allows software to clear one or more bits in the Software Interrupt register, without having to first read it.

Table 36: Software Interrupt Clear register (VICSoftIntClear - address 0xFFFF F01C) bit allocation *Reset value: 0x0000 0000*

Table 37: Software Interrupt Clear register (VICSoftIntClear - address 0xFFFF F01C) bit description

5.4.3 Raw Interrupt status register (VICRawIntr - 0xFFFF F008)

This is a read only register. This register reads out the state of the 32 interrupt requests and software interrupts, regardless of enabling or classification.

Table 38: Raw Interrupt status register (VICRawIntr - address 0xFFFF F008) bit allocation *Reset value: 0x0000 0000*

Bit	31	30	29	28	27	26	25	24
Symbol	٠	$\overline{}$	۰	٠	٠	۰	$\overline{}$	$\overline{}$
Access	RO	RO	RO	RO	RO	RO	RO	RO
Bit	23	22	21	20	19	18	17	16
Symbol	۰	$\overline{}$	AD ₁	BOD	I2C1	AD ₀	EINT ₃	EINT ₂
Access	RO	RO	RO	RO	RO	RO	RO	RO
Bit	15	14	13	12	11	10	9	8
Symbol	EINT ₁	EINT ₀	RTC	PLL	SPI1/SSP	SPI ₀	I2C0	PWM0
Access	RO	RO	RO	RO	RO	RO	RO	RO
Bit	$\overline{7}$	6	5	4	3	$\overline{2}$	1	$\bf{0}$
Symbol	UART1	UART0	TIMER1	TIMER0	ARMCore1	ARMCore0	$\overline{}$	WDT
Access	RO	RO	RO.	RO	RO	RO	RO	RO

Table 39: Raw Interrupt status register (VICRawIntr - address 0xFFFF F008) bit description

5.4.4 Interrupt Enable register (VICIntEnable - 0xFFFF F010)

This is a read/write accessible register. This register controls which of the 32 interrupt requests and software interrupts contribute to FIQ or IRQ.

Table 40: Interrupt Enable register (VICIntEnable - address 0xFFFF F010) bit allocation

Reset value: 0x0000 0000

5.4.5 Interrupt Enable Clear register (VICIntEnClear - 0xFFFF F014)

This is a write only register. This register allows software to clear one or more bits in the Interrupt Enable register (see Section 5.4.4 "Interrupt Enable register (VICIntEnable -0xFFFF [F010\)" on page 52](#page-4-0)), without having to first read it.

Table 42: Software Interrupt Clear register (VICIntEnClear - address 0xFFFF F014) bit allocation *Reset value: 0x0000 0000*

Bit	31	30	29	28	27	26	25	24
Symbol	$\overline{}$	۰	٠	$\overline{}$	$\overline{}$	۰	$\overline{}$	$\overline{}$
Access	WO	WO	WO	WO	WO	WO	WO	WO
Bit	23	22	21	20	19	18	17	16
Symbol	$\overline{}$	$\overline{}$	AD ₁	BOD	12C1	AD ₀	EINT ₃	EINT ₂
Access	WO	WO	WO	WO	WO	WO	WO	WO
Bit	15	14	13	12	11	10	9	8
Symbol	EINT ₁	EINT ₀	RTC	PLL	SPI1/SSP	SPI ₀	12C ₀	PWM0
Access	WO	WO	WO	WO	WO	WO	WO	WO
Bit	$\overline{7}$	6	5	4	3	$\overline{2}$	1	$\bf{0}$
Symbol	UART1	UART0	TIMER1	TIMER0	ARMCore1	ARMCore0	$\overline{}$	WDT

Table 43: Software Interrupt Clear register (VICIntEnClear - address 0xFFFF F014) bit description

5.4.6 Interrupt Select register (VICIntSelect - 0xFFFF F00C)

This is a read/write accessible register. This register classifies each of the 32 interrupt requests as contributing to FIQ or IRQ.

Reset value: 0x0000 0000

5.4.7 IRQ Status register (VICIRQStatus - 0xFFFF F000)

This is a read only register. This register reads out the state of those interrupt requests that are enabled and classified as IRQ. It does not differentiate between vectored and non-vectored IRQs.

Table 46: IRQ Status register (VICIRQStatus - address 0xFFFF F000) bit allocation *Reset value: 0x0000 0000*

Table 47: IRQ Status register (VICIRQStatus - address 0xFFFF F000) bit description

5.4.8 FIQ Status register (VICFIQStatus - 0xFFFF F004)

This is a read only register. This register reads out the state of those interrupt requests that are enabled and classified as FIQ. If more than one request is classified as FIQ, the FIQ service routine can read this register to see which request(s) is (are) active.

Table 49: FIQ Status register (VICFIQStatus - address 0xFFFF F004) bit description

5.4.9 Vector Control registers 0-15 (VICvectCntl0-15 - 0xFFFF F200-23C)

These are a read/write accessible registers. Each of these registers controls one of the 16 vectored IRQ slots. Slot 0 has the highest priority and slot 15 the lowest. Note that disabling a vectored IRQ slot in one of the VICVectCntl registers does not disable the interrupt itself, the interrupt is simply changed to the non-vectored form.

Table 50: Vector Control registers 0-15 (VICvectCntl0-15 - 0xFFFF F200-23C) bit description

5.4.10 Vector Address registers 0-15 (VICVectAddr0-15 - 0xFFFF F100-13C)

These are a read/write accessible registers. These registers hold the addresses of the Interrupt Service routines (ISRs) for the 16 vectored IRQ slots.

5.4.11 Default Vector Address register (VICDefVectAddr - 0xFFFF F034)

This is a read/write accessible register. This register holds the address of the Interrupt Service routine (ISR) for non-vectored IRQs.

5.4.12 Vector Address register (VICVectAddr - 0xFFFF F030)

This is a read/write accessible register. When an IRQ interrupt occurs, the IRQ service routine can read this register and jump to the value read.

Table 53: Vector Address register (VICVectAddr - address 0xFFFF F030) bit description

5.4.13 Protection Enable register (VICProtection - 0xFFFF F020)

This is a read/write accessible register. This one-bit register controls access to the VIC registers by software running in User mode.

Table 54: Protection Enable register (VICProtection - address 0xFFFF F020) bit description

5.5 Interrupt sources

[Table](#page-9-0) 55 lists the interrupt sources for each peripheral function. Each peripheral device has one interrupt line connected to the Vectored Interrupt Controller, but may have several internal interrupt flags. Individual interrupt flags may also represent more than one interrupt source.

Table 55: Connection of interrupt sources to the Vectored Interrupt Controller (VIC)

Block	Flag(s)	Mask	VIC Channel # and Hex
WDT	Watchdog Interrupt (WDINT)	0	0x0000 0001
	Reserved for Software Interrupts only	1	0x0000 0002
ARM Core	Embedded ICE, DbgCommRx	$\overline{2}$	0x0000 0004
ARM Core	Embedded ICE, DbgCommTX	3	0x0000 0008
TIMER0	Match 0 - 3 (MR0, MR1, MR2, MR3)	4	0x0000 0010
	Capture 0 - 3 (CR0, CR1, CR2, CR3)		
TIMER1	Match 0 - 3 (MR0, MR1, MR2, MR3)	5	0x0000 0020
	Capture 0 - 3 (CR0, CR1, CR2, CR3)		
UART0	Rx Line Status (RLS)	6	0x0000 0040
	Transmit Holding Register Empty (THRE)		
	Rx Data Available (RDA)		
	Character Time-out Indicator (CTI)		
UART ₁	Rx Line Status (RLS)	7	0x0000 0080
	Transmit Holding Register Empty (THRE)		
	Rx Data Available (RDA)		
	Character Time-out Indicator (CTI)		
	Modem Status Interrupt (MSI)[1]		
PWM ₀	Match 0 - 6 (MR0, MR1, MR2, MR3, MR4, MR5, MR6)	8	0x0000 0100
I ² CO	SI (state change)	9	0x0000 0200
SPI ₀	SPI Interrupt Flag (SPIF)	10	0x0000 0400
	Mode Fault (MODF)		
SPI1 (SSP)	TX FIFO at least half empty (TXRIS)	11	0x0000 0800
	Rx FIFO at least half full (RXRIS)		
	Receive Timeout condition (RTRIS)		
	Receive overrun (RORRIS)		
PLL	PLL Lock (PLOCK)	12	0x0000 1000
RTC	Counter Increment (RTCCIF)	13	0x0000 2000
	Alarm (RTCALF)		
System Control	External Interrupt 0 (EINT0)	14	0x0000 4000
	External Interrupt 1 (EINT1)	15	0x0000 8000
	External Interrupt 2 (EINT2)	16	0x0001 0000
	External Interrupt 3 (EINT3)	17	0x0002 0000
ADC0	A/D Converter 0 end of conversion	18	0x0004 0000

Philips Semiconductors UM10120

[1] LPC2134/6/8 Only.

5.6 Spurious interrupts

Spurious interrupts are possible in the ARM7TDMI based microcontrollers such as the LPC2131/2/4/6/8 due to asynchronous interrupt handling. The asynchronous character of the interrupt processing has its roots in the interaction of the core and the VIC. If the VIC state is changed between the moments when the core detects an interrupt, and the core actually processes an interrupt, problems may be generated.

Real-life applications may experience the following scenarios:

- 1. VIC decides there is an IRQ interrupt and sends the IRQ signal to the core.
- 2. Core latches the IRQ state.
- 3. Processing continues for a few cycles due to pipelining.
- 4. Core loads IRQ address from VIC.

Furthermore, It is possible that the VIC state has changed during step 3. For example, VIC was modified so that the interrupt that triggered the sequence starting with step 1) is no longer pending -interrupt got disabled in the executed code. In this case, the VIC will not be able to clearly identify the interrupt that generated the interrupt request, and as a result the VIC will return the default interrupt VicDefVectAddr (0xFFFF F034).

This potentially disastrous chain of events can be prevented in two ways:

- 1. Application code should be set up in a way to prevent the spurious interrupts from occurring. Simple guarding of changes to the VIC may not be enough since, for example, glitches on level sensitive interrupts can also cause spurious interrupts.
- 2. VIC default handler should be set up and tested properly.

5.6.1 Details and case studies on spurious interrupts

This chapter contains details that can be obtained from the official ARM website (http://www.arm.com), FAQ section under the "Technical Support" link: http://www.arm.com/support/faqip/3677.html.

What happens if an interrupt occurs as it is being disabled?

Applies to: ARM7TDMI

If an interrupt is received by the core during execution of an instruction that disables interrupts, the ARM7 family will still take the interrupt. This occurs for both IRQ and FIQ interrupts.

For example, consider the following instruction sequence:

```
MRS r0, cpsr
ORR r0, r0, #I_Bit:OR:F_Bit ;disable IRQ and FIQ interrupts
MSR cpsr_c, r0
```
If an IRQ interrupt is received during execution of the MSR instruction, then the behavior will be as follows:

• The IRQ interrupt is latched.

© Koninklijke Philips Electronics N.V. 2005. All rights reserved.

- **•** The MSR cpsr, r0 executes to completion setting both the I bit and the F bit in the CPSR.
- **•** The IRQ interrupt is taken because the core was committed to taking the interrupt exception before the I bit was set in the CPSR.
- **•** The CPSR (with the I bit and F bit set) is moved to the SPSR_IRQ.

This means that, on entry to the IRQ interrupt service routine, you can see the unusual effect that an IRQ interrupt has just been taken while the I bit in the SPSR is set. In the example above, the F bit will also be set in both the CPSR and SPSR. This means that FIQs are disabled upon entry to the IRQ service routine, and will remain so until explicitly re-enabled. FIQs will not be reenabled automatically by the IRQ return sequence.

Although the example shows both IRQ and FIQ interrupts being disabled, similar behavior occurs when only one of the two interrupt types is being disabled. The fact that the core processes the IRQ after completion of the MSR instruction which disables IRQs does not normally cause a problem, since an interrupt arriving just one cycle earlier would be expected to be taken. When the interrupt routine returns with an instruction like:

SUBS pc, lr, #4

the SPSR IRQ is restored to the CPSR. The CPSR will now have the I bit and F bit set, and therefore execution will continue with all interrupts disabled. However, this can cause problems in the following cases:

Problem 1: A particular routine maybe called as an IRQ handler, or as a regular subroutine. In the latter case, the system guarantees that IRQs would have been disabled prior to the routine being called. The routine exploits this restriction to determine how it was called (by examining the I bit of the SPSR), and returns using the appropriate instruction. If the routine is entered due to an IRQ being received during execution of the MSR instruction which disables IRQs, then the I bit in the SPSR will be set. The routine would therefore assume that it could not have been entered via an IRQ.

Problem 2: FIQs and IRQs are both disabled by the same write to the CPSR. In this case, if an IRQ is received during the CPSR write, FIQs will be disabled for the execution time of the IRQ handler. This may not be acceptable in a system where FIQs must not be disabled for more than a few cycles.

5.6.2 Workaround

There are 3 suggested workarounds. Which of these is most applicable will depend upon the requirements of the particular system.

5.6.3 Solution 1: test for an IRQ received during a write to disable IRQs

Add code similar to the following at the start of the interrupt routine.

 ; are next enabled. ; Rest of interrupt routine

This code will test for the situation where the IRQ was received during a write to disable IRQs. If this is the case, the code returns immediately - resulting in the IRQ not being acknowledged (cleared), and further IRQs being disabled.

Similar code may also be applied to the FIQ handler, in order to resolve the first issue.

This is the recommended workaround, as it overcomes both problems mentioned above. However, in the case of problem two, it does add several cycles to the maximum length of time FIQs will be disabled.

5.6.4 Solution 2: disable IRQs and FIQs using separate writes to the CPSR

MRS r0, cpsr ORR r0, r0, #I Bit ;disable IROs MSR cpsr_c, r0 ORR r0, r0, #F_Bit ;disable FIQs MSR cpsr_c, r0

This is the best workaround where the maximum time for which FIQs are disabled is critical (it does not increase this time at all). However, it does not solve problem one, and requires extra instructions at every point where IRQs and FIQs are disabled together.

5.6.5 Solution 3: re-enable FIQs at the beginning of the IRQ handler

As the required state of all bits in the c field of the CPSR are known, this can be most efficiently be achieved by writing an immediate value to CPSR_C, for example:

```
MSR cpsr_c, #I_Bit:OR:irq_MODE ;IRQ should be disabled
                                    ;FIQ enabled
                                    ;ARM state, IRQ mode
```
This requires only the IRQ handler to be modified, and FIQs may be re-enabled more quickly than by using workaround 1. However, this should only be used if the system can guarantee that FIQs are never disabled while IRQs are enabled. It does not address problem one.

5.7 VIC usage notes

If user code is running from an on-chip RAM and an application uses interrupts, interrupt vectors must be re-mapped to on-chip address 0x0. This is necessary because all the exception vectors are located at addresses 0x0 and above. This is easily achieved by configuring the MEMMAP register (see [Section 3.6.1 "Memory Mapping control register](#page--1-0) [\(MEMMAP - 0xE01F](#page--1-0) C040)" on page 25) to User RAM mode. Application code should be linked such that at 0x4000 0000 the Interrupt Vector Table (IVT) will reside.

Although multiple sources can be selected (VICIntSelect) to generate FIQ request, only one interrupt service routine should be dedicated to service all available/present FIQ request(s). Therefore, if more than one interrupt sources are classified as FIQ the FIQ interrupt service routine must read VICFIQStatus to decide based on this content what to

© Koninklijke Philips Electronics N.V. 2005. All rights reserved.

do and how to process the interrupt request. However, it is recommended that only one interrupt source should be classified as FIQ. Classifying more than one interrupt sources as FIQ will increase the interrupt latency.

Following the completion of the desired interrupt service routine, clearing of the interrupt flag on the peripheral level will propagate to corresponding bits in VIC registers (VICRawIntr, VICFIQStatus and VICIRQStatus). Also, before the next interrupt can be serviced, it is necessary that write is performed into the VICVectAddr register before the return from interrupt is executed. This write will clear the respective interrupt flag in the internal interrupt priority hardware.

In order to disable the interrupt at the VIC you need to clear corresponding bit in the VICIntEnClr register, which in turn clears the related bit in the VICIntEnable register. This also applies to the VICSoftInt and VICSoftIntClear in which VICSoftIntClear will clear the respective bits in VICSoftInt. For example, if VICSoftInt = 0x0000 0005 and bit 0 has to be cleared, VICSoftIntClear = 0x0000 0001 will accomplish this. Before the new clear operation on the same bit in VICSoftInt using writing into VICSoftIntClear is performed in the future, VICSoftIntClear = 0x0000 0000 must be assigned. Therefore writing 1 to any bit in Clear register will have one-time-effect in the destination register.

If the watchdog is enabled for interrupt on underflow or invalid feed sequence only then there is no way of clearing the interrupt. The only way you could perform return from interrupt is by disabling the interrupt at the VIC (using VICIntEnClr).

Example:

Assuming that UART0 and SPI0 are generating interrupt requests that are classified as vectored IRQs (UART0 being on the higher level than SPI0), while UART1 and I2C are generating non-vectored IRQs, the following could be one possibility for VIC setup:

After any of IRQ requests (SPI0, I2C, UART0 or UART1) is made, microcontroller will redirect code execution to the address specified at location 0x0000 0018. For vectored and non-vectored IRQ's the following instruction could be placed at 0x0000 0018:

LDR pc, [pc,#-0xFF0]

This instruction loads PC with the address that is present in VICVectAddr register.

In case UART0 request has been made, VICVectAddr will be identical to VICVectAddr0, while in case SPI0 request has been made value from VICVectAddr1 will be found here. If neither UART0 nor SPI0 have generated IRQ request but UART1 and/or I2C were the reason, content of VICVectAddr will be identical to VICDefVectAddr.