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Contents

| Safety Considerations | | | 4 |
|---------------------------|--|-----------------|------------|
| Introduction BIS C Ident | ification System | | 5-7 |
| Application BIS-C60_1 F | Processor, Basic knowledge for application. | | 8/9 |
| BUS interface: InterBus | | | 10/11 |
| Compatibility with BIS C | C-6_1 processor | | 12 |
| Function Description: | Communication with the processor | | 13 |
| · | In- and output buffer on INTERBUS | | 14/15 |
| | Output buffer, configuration and explanation | on | 16-20 |
| | Input buffer, configuration and explanation | ١ | 21-24 |
| | Configuring the BIS C-60_1 processor | | 25-27 |
| | Processing code tags | | |
| | Reading and writing | | 28 |
| | Codetag Present / Auto-Read | | 29 |
| | Reading and writing in dynamic mo | | |
| | Reading and writing with simultane | ous data transn | nission 30 |
| | CRC initialization | | 31 |
| | Mixed Data Access | | |
| | Reading and writing with program | | 34 |
| | Examples for protocol sequence | | |
| Read/Write Times | | | 53/54 |
| LED Display | | | |
| | | BIS C-6001 | |
| Mounting Head / Proces | sor | 57 | 70 |
| | Interface information | | |
| Installing the connectior | cables / Mounting the PG connection | 59/60 | |
| Interface information / W | /iring diagrams | 61-65 | 71-76 |
| Changing the EEPROM. | | 66 | / / |
| | | | |
| | | | |
| Appendix, ASCII Table. | | | 81 |
| | | PALLIE | |

4

Safety Considerations

Approved operation

Series BIS C-60_1 processors along with the other BIS C system components comprise an identification system and may only be used for this purpose in an industrial environment in conformity with Class A of the EMC Law.

Installation and Operation

Installation and operation should be carried out by technically trained personnel only. Unauthorized access and improper use will lead to loss of warranty and liability claims.

When installing the processor, consult the section on wiring diagrams carefully. Special caution must be used when wiring the processor to external controllers, particularly with respect to selection and polarity of the signals and power supply.

Only approved power supplies may be used with the processor. See the section on Technical Data for details.

Use and Checking

The relevant safety procedures must be followed when using the Identification System. In particular, steps must be taken to ensure that no danger to persons or equipment can arise should a fault occur in the Identification System.

This includes maintaining the published ambient operating conditions and regular checking of the functionality of the Identification System with all its associated components.

Fault Conditions

As soon as there is evidence that the Identification System is not functioning properly, it should be taken out of service and protected against unauthorized use.

Scope

This description is valid for series BIS C-6001-023...03-KL2 processors and both the ST8 and ST9 versions of series BIS C-6021-023-050-03-ST_.

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4 E

Introduction **BIS C Identification System**

This manual is designed to assist the user in setting up the control program and installing and starting up the components of the BIS C-60_1 Identification System, and to assure rapid, trouble-free operation.

Principles

The BIS C-60_1 Identification System belongs in the category of

-contact systems for reading and writing

This dual function permits applications for not only transporting information in fixed-programmed code tags, but also for gathering and passing along up-to-date information as well.



If 2 read/write heads are connected to a BIS C-60_1 processor, both heads can be operated independently of each other. This means for example that you can read a code tag from one head while writing to another code tag at the other head.

Applications

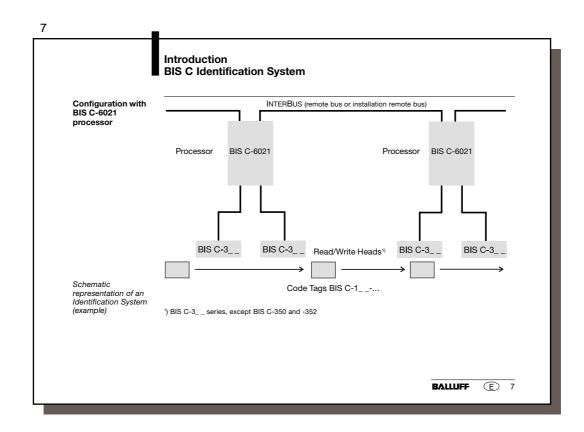
Some of the notable areas of application include

- for controlling material flow in production processes (e.g. in model-specific processes), for workpiece conveying in transfer lines, in data gathering for quality assurance , for gathering safety-related data,
- in tool coding and monitoring;
- in equipment organization;
- in storage systems for monitoring inventory movement;
- in transporting and conveying systems;
- in waste management for quantity-based fee assessment.

BALLUFF E 5



6 Introduction **BIS C Identification System** System Components The main components of the BIS C Identification System are - Processor. - Read/Write Heads and - Code Tags. INTERBUS (remote bus) Configuration with processor Processor BIS C-6001 Processor BIS C-6001 with Adapter with Adapter BIS C-670 BIS C-650 Processor BIS C-6001 Read/Write Read/Write mit Kopf BIS C-35 BIS C-3__ BIS C-65_ BIS C-3__ Schematic representation of an Code Tags BIS C-1_ _-... Identification System (example) $^{\mbox{\tiny 1}}\mbox{)}$ BIS C-3_ $_$ series, except BIS C-350 and -352 2) BIS C-350 or -352 only 6 E BALLUFF



BIS C-60_1 Processor Basic knowledge for application

Selecting System Components

8

The BIS C-6001 processor has a plastic housing. Connections are made through a terminal strip, with the cables secured by PG fittings. A single read/write head from BIS C-65_series can be directly mounted to the processor, which creates a compact unit. If the BIS C-650 adapter is attached instead of the BIS C-65_read/write head, two read/write heads may be cable connected. If the BIS C-670 adapter is attached, one read/write head may be cable connected.

The BIS C-6021 processor has a metal housing. Connection is made through round connectors. Two read/write heads can be cable connected to the processor.

Series BIS C-60_1 processors have in addition a digital input. The input has various functions depending on the configuration (see Parametering).

Whether the compact version of the processor with integrated read/write head makes sense or whether the external solution is preferred depends primarily on the spatial arrangement of the components. There are no functional limitations. All read/write heads are suitable for both static and dynamic reading. Distance and relative velocity are based on which code tag is selected. Additional information on the read/write heads in series BIS C-65_ and series BIS C-3__ including all the possible code tag/read-write head combinations can be found in the manuals for the respective read/write heads.

The system components are electrically supplied by the processor. The code tag represents an free-standing unit and needs no line-carried power. It receives its energy from the read/ write head. The latter constantly sends out a carrier signal which supplies the code head as soon as the required distance between the two is reached. The read/write operation takes place during this phase. Reading and writing may be dynamic or static.

BIS C-60_1 Processor Basic knowledge for application

Control Function

The processor writes data from the host system to the code tag or reads data from the tag through the read/write head and prepares it for the host system. Host systems may include:

- a host computer (e.g. industrial PC) or
- a programmable logic controller (PLC)

Data checking

When sending data between the read/write head and the code tag a procedure is required for recognizing whether the data were correctly read or written.

The processor is supplied with standard Balluff procedure of double reading and comparing. In addition to this procedure a second alternative is available: CRC-16 data checking.

Here a test code is written to the code tag, allowing data to be checked for validity at any time or location.

| Advantages of CRC-16 | Advantages of double reading |
|---|--|
| Data checking even during the non-active phase (CT outside read/write head zone). | No bytes on the code tag need to be reserved for storing a check code. |
| Shorter read times since each page is read only once. | Shorter write times since no CRC needs to be written. |



It is not permitted to operate the system using both check procedures!

BALLUFF E 9



10

BUS interface: INTERBUS

INTERBUS

Communication between the BIS C-60_1 processor and the host system is via INTERBUS.

The INTERBUS system consists of three components:

- the wiring module (rack card for industrial PC or PLC),
 Bus terminal as network node and/or
 the I/O modules (here the BIS C-60_1 processor).

Depending on the wiring module up to a maximum of 63 BIS C-60_1 processors can be

The BIS C-6001 processor is used as a remote bus station. The BIS C-6021 can be used as a remote bus or installation bus station.



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Important hints for use with PLC:

In some control systems the PROFIBUS-DP data area is not synchronously transmitted with the updating of the input/output content. If more than 2 bytes of data are sent, a mechanism must be used which guarantees that the data in the PLC and the data in the BIS C are always identical!

2nd alternative: Set 2nd bit header

and alternative: Set 2nd bit header
Data exchange between PLC and BIS is controlled by the so-called bit header. This is always the
first byte of the respective read/write head in the data buffer. This bit header exists both in the
input range (data from BIS to the PLC) and in the output range (data from the PLC to the BIS). If
this bit header is also sent as the last byte, a comparison of these two bytes can be used to guarantee the consistency of the transmitted data.

In this method the PLC cycle is unaffected nor is the bus access time changed. All that is required is that a byte in the data buffer be used for the 2nd bit header instead of for user data.

This 2nd alternative is the Balluff recommended setting (factory default).

BUS interface: INTERBUS

Address setting is done on the module (not on the I/O modules), i.e. not on the BIS $C-60_1$ processor). There are two types of addressing possible:

- logical addressing, and
 physical addressing.

Logical Addressing

Logical addressing permits free addressing of each module.

Advantage: high security and flexibility; Disadvantage: more difficult at setup.

For the BIS C-60_1 Identification System use:

I/O Module Type IDENT-No. IN-Address (Byte) OUT-Address (Byte) Processor BIS C-60_1 03 16 16

Physical Addressing

Advantage: easy to configure at setup;
Disadvantage: changes in module location when power was off are recognized upon initialization, but are not made known to the user.

BALLUFF E 11



12

Compatibility with BIS C-6_1 processor

Setting compatibility Compatibility with the BIS C-6_1 processors is established using terminal X5 and a jumper.



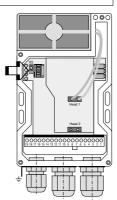
If the BIS C-60_1 processor is set to be compatible with the BIS C-601 or BIS C-621, all settings for data exchange must be made as described in the sections on parametering, function description, protocol sequence and LED display in the user's manual for the BIS C-6_1 processor! This user's manual can be mailed on request, or you may download it from the Internet at www.balluff.de.

123

| Jumper setting | Processor |
|----------------|---------------------------|
| Terminal X5 | compatible with BIS C-6_1 |
| 1-2 | no |
| 2-3 | yes |

In the illustration compatibility with the BIS C-6_1 is not set.

To open the cover of the BIS C-6001 processor, see \boxdot 58, and for BIS C-6021 see \boxdot 71.



Terminal X5 (with cover removed)

Function Description Communication with the processor

Basic Procedure

Communication between the host system and the processor takes place using a fixed protocol sequence. Data integrity from the control to the processor and vice-versa is indicated by a control bit. This bit is used to implement a handshake between the control and the processor.

Following is a simplified representation of the sequence of a job sent from the control to the

- The control sends a command designator to the processor together with the associated command parameters and sets a bit (AV bit). This bit indicates to the processor that the transmitted data are valid and that the job is now beginning.
- 2. The processor takes the job and sets a bit (AA bit), which indicates this to the control.
- If an additional exchange of data between the control and the processor is required to carry out the job, each uses a bit (TI bit and TO bit) to indicate that the control / processor is now ready for additional data exchange or has accepted the received data.
- 4. Once the processor has carried out the job correctly, it sets a bit (AE bit).
- 5. Once the control has accepted all the important data, it indicates this to the processor by resetting the bit that was set at the beginning (AV bit).
- 6. The processor now in turn sets all the control bits that were set during the sequence (AA bit, AE bit) and is ready for the next job.

BALLUFF



14

Function description In- and output buffer on INTERBUS

Input and Output Buffers

To transmit the control bit, commands and data between the BIS C-60 1 processor and the host system, the latter must provide two fields. These two fields are:

- the output buffer

For the control bit (bit header) and controller commands sent to the BIS identification system,

For the data to be written, and For configuring the BIS C-60_1 processor

- the input buffer

For the control bit (bit header) of the BIS C-60_1 processor, For the data to be read,

For the ID's and error codes coming from the BIS identification system, and For reading out the configuration data.

The total buffer size of the BIS C-60_1 is 16 bytes for the input buffer and 16 bytes for the output buffer. This total buffer size is divided into 2 sectors

Buffer sector 1 for Read/Write Head 1: 8 bytes input buffer, 8 bytes output buffer Buffer sector 2 for Read/Write Head 2: 8 bytes input buffer, 8 bytes output buffer

Please note the basic procedure on 11 13 and 28...34 and the examples on pages 11 35...52.

14 E

Function description In- and output buffer on INTERBUS

Input and Output Buffers (continued)

Example: Using a PLC, the buffer sector for the BIS C-60_1 will start at input byte EB32 and output byte AB 32.

EB 0 / AB 0 Memory map: PLC: BIS: PLC buffer Read/write head 1 (R/W1): Input buffer Output buffer Head R/W 1 Subaddress 0 from EB 32 to EB 39 from AB 32 to AB 39 buffer for R/W 1 Subaddress 7 Read/write head 2 (R/W2): Input buffer Output buffer Head R/W 2 from EB 40 to EB 47 from AB 40 to AB 47 Subaddress 0 buffer for R/W 2 Subaddress 7

Please note the basic procedure on 13 and 28...34 and the examples on pages 11 35...52.

| Note that these buffers can be in two different | Sequence 1 | Sequence 2 |
|---|---------------|---------------|
| sequences depending on the type of control. | Subaddress 00 | Subaddress 01 |
| | 01 | 00 |
| The following description is based on sequence 1! | 02 | 03 |
| | 03 | 02 |
| | 04 | 05 |
| | 05 | 04 |
| | 06 | 07 |
| | 07 | ne |

BALLUFF

E 15

16

Function Description Output buffer, configuration and explanation

Configuration of the output buffer for one (1) read/write head

| Bit No. | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | |
|--------------------|--------------------------|-------------|----------|----|------|----|-----------|----------|----------|
| Subaddress | | | | | | | | | |
| 00нех = Bit Header | CT | П | | | | GR | | AV | Bit Name |
| 01Hex | Com | nand Desi | gnator | or | Data | or | Config. | 1st byte | |
| 02Hex | Start Address (Low Byte) | | | or | Data | or | Config. 1 | 2nd Byte | |
| 03Hex | Start A | ddress (Hiç | gh Byte) | or | Data | or | Config. | 3rd Byte | |
| 04нех | No. of | Bytes (Lo | v Byte) | or | Data | or | Config. | 4th Byte | |
| 05Hex | No. of | Bytes (Hig | h Byte | or | Data | or | Config. | 5th Byte | |
| 06Hex | | | | | Data | or | Config. | 6th Byte | |
| 07Hex | 2nd Bit | Header (a: | s above) | or | Data | | | | |

Description of Output Buffer

| Sub- address | Bit Name | Meaning | Function Descri | iption | |
|---------------------------------|-------------|-------------------------|---|------------|--|
| 00 _{hex} Bit Header | CT | Code tag type 0 1 | Select code tag 32 Byte bloc 64 Byte bloc | k size | for code tag type: BIS C-102, -03, -04, -05 BIS C-110, -11, -30 |
| | TI | Toggle-Bit In | for reading: for writing: | new Cor | ntroller is ready to accept n/additional data. ntroller has prepared n/additional data. |
| | GR AV | Ground state Command | | , | n to go to ground state. that a job is waiting. |

Please note the basic procedure on 13 and 28...34 and the examples on pages 35...52.

BALLUFF

(continued next 1)

Function Description Output buffer, configuration and explanation

Description of Output Buffer (continued)

| Sub- address | Meaning | Function Description |
|-------------------|-------------------|--|
| 01 _{Hex} | Command de | esignator |
| | 00 _{Hex} | No command present |
| | 01 _{Hex} | Read code tag |
| | 02 _{Hex} | Write to code tag |
| | 04 _{Hex} | Configure processor |
| | 05 _{Hex} | Read configuration data |
| | 06HEX | Store program in the EEPROM for the Mixed Data Access function |
| | 07 _{HEX} | Store the start address for the Auto-Read function in the EEPROM |
| | 12 _{Hex} | Initialize the CRC-16 data check |
| | 21 _{Hex} | Read code tag using Mixed Data Access function (corresponding to the program stored in the EEPROM) |
| | 22 _{Hex} | Write to code tag using the Mixed Data Access function (corresponding to the program stored in the EEPROM) |
| or | Configuration | n 1st byte |
| | 00нех | Default value (factory setting). Changes depending on the configuration. |
| or: | Data | for writing to the code tag |

Please note the basic procedure on 11 13 and 28...34 and the examples on pages 11 35...52.

BALLUFF E 17



18

Function Description Output buffer, configuration and explanation

Description of Output Buffer (continued)

| Sub- address | Meaning | Function Description | | | |
|-----------------|--------------------------------------|---|--|--|--|
| 02 Hex | Start address (Low Byte) | Address at which reading from or writing to the code tag begins. (The Low Byte includes the address range from 0 to 255). | | | |
| or: | Start address (Low Byte) | Address for the Auto-Read function, starting at which the cotag is to be read. The value is stored in the EEPROM. (The Lo Byte covers the address range from 0 to 255). | | | |
| or: | Program No | Number of the program to be stored in the EEPROM in conjunction with command ID 06 _{Hex} for Mixed Data Access function. | | | |
| or: | Program No. | Number of the program stored in the EEPROM for read or write operations in conjunction with command ID 21 _{Hex} or 22 _{Hex} for t Mixed Data Access function. | | | |
| or | Configuration 2 80 _{Hex} | and byte Default value (factory setting) Changes depending on the configuration. | | | |
| or: | Data | for writing to the code tag. | | | |

(continued next 🖺)

Please note the basic procedure on 11 13 and 28...34 and the examples on pages 11 35...52.

18 E

Function Description Output buffer, configuration and explanation

Description of Output Buffer (continued)

| Sub- address | Meaning | Function Description | | | |
|-----------------|--------------------------------------|---|--|--|--|
| 03 Hex | Start address (High Byte) | Address for reading from or writing to the code tag. (The High Byte is additionally used for the address range from 256 to 8,191) | | | |
| or: | Start address (High Byte) | Address for the Auto-Read function, starting at which the code tag is to be read. The value is stored in the EEPROM. (The Hig Byte is also required for the address range from 256 to 8,191). | | | |
| or: | Configuration 3 00 _{Hex} | rd byte Default value (factory setting) This value must not be changed! | | | |
| or: | Data | for writing to the code tag | | | |
| 04 Hex | No. of bytes (Low Byte) | Number of bytes to read or write beginning with the start address (the Low Byte includes from 1 to 255 bytes). | | | |
| or: | Configuration 4 82 _{Hex} | th byte Default value (factory setting) Changes depending on the configuration. | | | |
| or: | Data | for writing to the code tag | | | |

Please note the basic procedure on 11 13 and 28...34 and the examples on pages 11 35...52.

BALLUFF E 19



20

Function Description Output buffer, configuration and explanation

Description of Output Buffer (continued)

| Sub- | Meaning | Function Description | | | | |
|---------------|-----------------------------|---|--|--|--|--|
| address | | | | | | |
| 05 Hex | No. of bytes (High Byte) | Number of bytes to read or write beginning with the start addres (the High Byte is additionally used for the range between 256 ar 8,191 bytes). | | | | |
| or: | Configuration 5th byte | | | | | |
| | 00нех | Default value (factory setting) Changes depending on the configuration. | | | | |
| or: | Data | for writing to the code tag. | | | | |
| 06 Hex | Configuration 6 | th byte | | | | |
| | 00нех | Default value (factory setting) This value must not be changed! | | | | |
| or: | Data | for writing to the code tag. | | | | |
| 07 Hex | 2nd Bit header | The data are valid if the 1st and 2nd bit header are identical. | | | | |
| or: | Data | for writing to the code tag. | | | | |

Please note the basic procedure on 11 13 and 28...34 and the examples on pages 11 35...52.

20 E

Function Description Input buffer, configuration and explanation

Configuration of the input buffer for one (1) Read/Write head

| Bit No. | 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | |
|--------------------|--------------------|---------|-----------|----------|---------|----------|----------|----------|----------|
| Subaddress | | | | | | | | | |
| 00нех = Bit Header | BB | HF | TO | IN | AF | Æ | AA | CP | Bit name |
| 01Hex | Error code or Data | | | | or | Config. | 1st byte | | |
| 02Hex | Data | | | | | or | Config. | 2nd Byte | |
| 03Hex | Data | | | | or | Config. | 3rd Byte | | |
| 04Hex | Data | | | or | Config. | 4th Byte | | | |
| 05Hex | Data | | | or | Config. | 5th Byte | | | |
| 06Hex | Data | | | or | Config. | 6th Byte | | | |
| 07 _{Hex} | | 2nd Bit | Header (a | s above) | | or | D | ata |] |

Description of Input Buffer

Sub-Bit Meaning Function Description address Name **00**Hex ВВ Ready The BIS Identification System is in the Ready state. Bit Header HF Cable break from read/write head or Head Error no read/write head connected. for read: BIS has new/additional data ready. for write: BIS is ready to accept new/additional data. TO Toggle-Bit Out (continued on next 1)

Please note the basic procedure on 13 and 28...34 and the examples on pages 11 35...52.

BALLUFF E 21



22

Function Description Input buffer, configuration and explanation

Description of Input Buffer (continued)

| Sub- address | Bit Name | Meaning | Function Description |
|-----------------|-------------|-----------------|---|
| 00 Hex | (contin | ued) | |
| Bit Header | ĺΝ | Input | If the parameter "Input IN" is 1, this bit indicates the state of the Input. $\label{eq:input}$ |
| | AF | Command Error | The command was incorrectly processed or aborted. |
| | AE | Command end | The command was finished without error. |
| | AA | Command start | The command was recognized and started. |
| | CP | Codetag Present | Code tag present within the active zone of the read/write head. |

| - | | |
|-----------------|-------------------|---|
| Sub- address | Meaning | Function Description |
| 01 Hex | Error code | Error number is entered if command was incorrectly processed or aborted. Only valid with AF bit! |
| | 00 _{Hex} | No error. |
| | 01 _{Hex} | Reading or writing not possible because no code tag is present in the active zone of a read/write head. |
| | 02 _{Hex} | Read error. |
| | 03нех | Code tag was removed from the active zone of the read/write head while it was being read. |
| | 04 _{Hex} | Write error. |
| | 05нех | Code tag was removed from the active zone of the read/write head while it was being written. |
| | 07 _{Hex} | AV bit is set but the command designator is missing or invalid. Number of bytes is 00 _{Hex} . |
| | (continued on | |

Please note the basic procedure on 11 13 and 28...34 and the examples on pages 11 35...52.

BALLUFF

Function Description Input buffer, configuration and explanation

Description of Input Buffer (continued)

| Sub- address | Meaning | Function Description | | | |
|-------------------|------------------------|--|--|--|--|
| 01 _{Hex} | Error code (continued) | | | | |
| | 09 _{Hex} | Cable break to select read/write head, or head not connected. | | | |
| | 0C _{Hex} | The EEPROM cannot be read/programmed. | | | |
| | 0D _{Hex} | Faulty communication with the code tag. | | | |
| | ОЕнех | The CRC of the read data does not coincide with the CRC of the code tag. | | | |
| | 0F _{Hex} | Contents of the 1st and 2nd bit header (1st and last bytes) of the output buffers are not identical (2nd bit header must be served). | | | |
| | 11 _{Hex} | Invoking a function that is not possible, since the processor is in "compatible with BIS C-6_1" mode. | | | |
| or: | Configuratio | n 1st byte | | | |
| | 00нех | Default value (factory setting). Changes depending on the configuration. | | | |
| or: | Data | Data which was read from the code tag. | | | |

Please note the basic procedure on 11 13 and 28...34 and the examples on pages 11 35...52.

BALLUFF E 23



24

Function Description Input buffer, configuration and explanation

Description of Input Buffer (continued)

| Sub- address | Meaning | Function Description |
|-----------------|--------------------------------------|---|
| 02 Hex | Configuration 2 80 _{Hex} | ind byte Default value (factory setting). Do not change! |
| or: | Data | Data which was read from the code tag. |
| 03 Hex | Configuration 3 00 _{нех} | ord byte Default value (factory setting). Changes depending on the configuration. |
| or: | Data | Data which was read from the code tag. |
| 04 Hex | Configuration 4 82 _{Hex} | th byte Default value (factory setting). Changes depending on the configuration. |
| or: | Data | Data which was read from the code tag. |
| 05 Hex | Configuration 5 | ith byte Default value (factory setting). Changes depending on the configuration. |
| or: | Data | Data which was read from the code tag. |
| 06 Hex | Configuration 5 | ith byte Default value (factory setting). Do not change! |
| or: | Data | Data which was read from the code tag. |
| 07 Hex | 2nd Bit header | The data are valid if the 1st and 2nd bit headers are in agreement. |
| or: | Data | Data which was read from the code tag. |

Please note the basic procedure on 11 13 and 28...34 and the examples on pages 11 35...52.

BALLUFF

Function Description Configuring the BIS C-60_1 processor

Configuration,

The following functions can be activated / deactivated through the configuration:

- CRC-16 data check:

If this function is activated, the correctness of the read or written data is ensured by a CRC-16 data check (see 2 9).

- Simultaneous data transmission for both read/write heads:

With simultaneous data transmission shorter read/write times can be achieved depending on the amount of data to be read/written and the type of controller.

Dynamic operation on Read/Write Head 1 or 2:

If dynamic operation is parametered, a read/write job can be sent even though there is no code tag in the active zone of the head. As soon as a code tag passes by the head, the command is immediately carried out.

- "Auto-Read" for Read/Write Head 1 or 2:

If this function is activated, the processor reads out the first (max. 31) bytes from the code tag starting at a defined start address as soon as the tag enters the active zone of the read/ write head. The start address must first have been stored in the processor's EEPROM with the command ID 07_{Hex}

 2nd bit header at end of in- and output buffer:
 The 2nd bit header (factory setting) prevents data from being accepted by the bus as long as it is not fully updated.

 Display state of the digital input in the bit header of the input buffer:
 If this function is activated, the IN-bit displays the state of the digital input of the processor: $IN = 0 \rightarrow digital input low; IN = 1 \rightarrow digital input high$

Reset BIS C-60_1 processor through the digital input:
 If this function is activated, the processor is reset when the digital input is set to high.

BALLUFF E 25



26

Function Description Configuring the BIS C-60_1 processor

Configuration

The BIS C-60_1 processor is configured by the controller using the output buffer. The configuration data are arranged within 6 configuration bytes that are sent to the BIS C-60_1 processor using the command identifier 04_{Hex} (see Example 13 on ☐ 51). Command identifier 05_{Hex} is used to read out the current device configuration (see Example 14 on 1 52).



To input the configuration, all 6 bytes must be entered in Hex. Only the named bits are permitted to be changed. If any of the other bits are changed, there is no assurance that the BIS C-60_1 $\,$ will function properly.

The default values of the 6 bytes are (factory setting):

2nd byte 3rd byte 5th byte 6th byte 1st byte 4th byte 00000000 Binary 000**0**0**0**00 100<u>00</u>000 00000000 **10**0000**1**0 000<u>00</u>000 bit 7 bit 2 bit 3 bit 4 These are used for bit 5 bit 5 hit 8 bit 5

configuration:

Having the following 1st byte, bit 5 Activate CRC-16 data checking

> 1st byte, bit 3 Activate simultaneous data transmission for both read/write heads

2nd byte, bit 5 Dynamic mode on read/write head 1 (for effects on read/write times, see 11 53/54)

Activate Auto-Read function starting at specified address after CT-Present 2nd byte, bit 4 for Head 1 (the amount data read is 6 bytes for a double bit header or 7 bytes for a single bit header)

Bit state: 0 = no 1 = yes

BALLUFF

Function Description Configuring the BIS C-60_1 processor

Configuration (continued)

Bit state: 0 = no 1 = yes

Arrange a 2nd bit header at the end of the input and output buffers. 4th byte, bit 8, 4th byte, bit 7 Display state of the digital input in the bit header of the input buffers

4th byte, bit 2 Reset the BIS C-60_1 processor through the digital input

Dynamic mode on read/write head 2 (for effects on read/write times, see 1153/54) 5th byte, bit 5

5th byte, bit 4 Activate Auto-Read function starting at specified address after CT-Present

for Head 2 (the amount data read is 6 bytes for a double bit header or 7 bytes for a single bit header)

BALLUFF E 27



28

Function Description Processing code tags

Reading and writing

To carry out a read or write job, the code tag must be located in the active zone of the read/

A read/write job has the following sequence (see examples on $\square 37...43$):

- 1. The host sends to the output buffer:

 - the command designator to subaddress 01_{Hex},
 the start address for reading or writing to subaddress 02_{Hex}/03_{Hex},
 the number of bytes for reading or writing to subaddress 04_{Hex}/05_{Hex},
 - the CT bit according to the code tag type (block size),
 and sets the AV bit in the bit header to high.
- 2. The processor:
 - takes the request (AA in the bit header of the input buffer to high), begins to transport the data;

begins to transport the data; read = from code tag to input buffer, write = from output buffer to code tag. Larger data quantities are sent in blocks: block size of 6 bytes with 2nd bit header,

block size of 7 bytes without 2nd bit header.
The toggle bits in the two bit headers are used as a kind of handshaking between the host and the BIS C-60_1 processor.

3. The processor has processed the command correctly (AE bit in the bit header of the input buffer). If an error occurred during execution of the command, an error number will be written to subaddress 01Hex of the input buffer and the AF bit in the bit header of the input buffer will be set.

28 E

Function Description Processing code tags

Codetag Present

As soon as the code tag enters the active one of the read/write head, the processor indicates this by setting the CP bit (Codetag present)



To accelerate the reading of small amounts of data, the ID system makes the first bytes of the code tag available in the input buffer of the respective read/write head as soon as the tag is detected (6 bytes for double bit header, 7 bites for single bit header).

The data are only valid after the rising edge of the CP bit in the bit header of the input buffer. They remain valid until the falling edge of the CP bit, or until the controller issues a new job.

Auto-Read

If the Auto-Read function is activated, data are automatically read (6 bytes for a double bit header / 7 bytes for a single bit header) beginning with a start address as soon as a code tag is recognized. The read process begins at the start address that was specified by command identifier 07Hex. Each head can have its own start address assigned. The start addresses are stored in the EEPROM of the BIS C-60_1 processor.



To obtain correct data output, use command identifier 07_{Hex} for each partial buffer Head 1 and/or

If the Auto-Read function is not activated, the processor runs in standard mode and sends 6 bytes for a double bit header or 7 bytes for a single bit header starting with code tag address 0.

BALLUFF



30

Function Description Processing code tags

Reading and writing in dynamic mode

In normal operation a read/write job is rejected by the processor BIS C-60_1 by setting the AF bit and an error number if there is no code tag in the active zone of the read/write head. If dynamic mode is configured, the processor accepts the read/write job and stores it. When a code tag is recognized, the stored job is carried out.

Reading and writing with simultaneous data transmission

Reading without simultaneous data transmission: In the case of a read job the processor first reads our all requested data from the code tag after receiving the start address and the desired number of bytes, and then sets the AE bit. Then the data read from the code tag are written to the input buffer. In the case of larger data amounts this is done in blocks, controlled by the handshake with the toggle bits as described on \(\Delta\) 28.

Reading with simultaneous data transmission: In the case of a read job the processor begins to send the data to the input buffer as soon as the first 6 bytes (or 7 bytes for a single bit header) have been read from the code tag beginning with the start address, and indicates this by inverting the TO bit. As soon as the controller inverts the TI bit, the processor sends the data that have been read in the meantime to the input buffer. This is repeated until th processor has read all the desired data from the code tag. Now the processor sets the AE bit and outputs the remaining data to the input buffer.

Writing without simultaneous data transmission: In the case of a write job the process waits until it has received all the data that need to be written from the controller. Only then are the data written to the code tag as described on \bigcirc 28.

Writing with simultaneous data transmission: In the case of a write job the processor begins to write the data to the code tag as soon as it has received the first data to be written from the controller's output buffer. Once all the data have been written to the code tag, the AE bit is set.



Function Description Processing code tags

CRC initialization

To be able to use the CRC check, the code tag must first be initialized with the command identifier 12_{Hex} (see ☐ 35). The CRC initialization is used like a normal write job. The latter is rejected (with an error message) if the processor recognizes that the code tag does not contain the correct CRC. Code tags as shipped from the factory (all data are 0) can immediately be programmed with a CRC check.

If CRC-16 data checking is activated, a special error message is output to the interface whenever a CRC error is detected.

If the error message is not caused by a failed write request, it may be assumed that one or more memory cells on the code tag is defective. That code tag must then be replaced.

If the CRC error is however due to a failed write request, you must reinitialize the code tag in order to continue using it.

The checksum is written to the code tag as a 2-byte wide datum. Two bytes per page are 'lost', i.e., the page size becomes 30 bytes or 62 bytes depending on code tag type (setup of page size see \(\frac{1}{2} \) 16). This means that the actual usable number of bytes is reduced:

| Code tag type | | Usable bytes |
|---------------|---|--------------|
| 128 bytes | = | 120 bytes |
| 256 bytes | = | 240 bytes |
| 511 bytes *) | = | 450 bytes |
| 1023 bytes *) | = | 930 bytes |
| 2047 bytes *) | = | 1922 bytes |
| 2048 bytes | = | 1984 bytes |
| 8192 bytes | = | 7936 bytes |

*) The last code tag page for these EEPROM-based code tags is not available.

BALLUFF E 31



32

Function Description Processing code tags

Mixed Data Access

Small read/write programs can be stored in the BIS C-60_1 processor's EEPROM.

The Mixed Data Access function is useful when the required information is stored on the code tag at various addresses. This function makes it possible to read out this "mixed", i.e. non-contiguously stored data from the code tag in a single procedure and using just one com-

Up to 10 programs with up to 25 instructions can be stored. Each program instruction contains a "start address" and a "number of bytes" specification. The amount of data for reading may not exceed 2 kB.

The command identifier 06Hex is used to send the read/write program to the BIS C-60_1 processor. One program per command can be stored. All 25 program records plus an additional 2 bytes with FFHex FFHex as a terminator must always be sent. This means a total of **104 bytes** of information per program must be sent (including the command identifier and program number).



The individual program records must all be contiguous. They must be sent one after the other and be terminated with FFHexFFHex as a terminator. It is recommended that the remaining, unused memory sector be filled with FFHexFFHex.

If an address range is selected twice, the data will also be output twice.



Function Description Processing code tags

Mixed Data Access

The following shows the structure of a program:

| Program structure | Subaddress | Value | Range |
|--|-------------------|-------------------|-----------------|
| Command designator | 01 _{HEX} | 06 _{Hex} | |
| Program record Program number | 02HEX | 01 _{Hex} | 01Hex to 0AHex |
| 1st data record: | OLI III. | O THE | O THE CO OF THE |
| Start address Low Byte | 03 _{Hex} | | |
| Start address High Byte | 04HEX | | |
| Number of bytes Low Byte | 05HEX | | |
| Number of bytes High Byte 2nd data record: | 06нех | | |
| 25th data record: | | | |
| Start address Low Byte | 03HEX | | |
| Start address High Byte | 04 _{HEX} | | |
| Number of bytes Low Byte | 05 _{HEX} | | |
| Number of bytes High Byte | 06HEX | | |
| Terminator | FFHEX FFHEX | | |

To store a second program, repeat this process.

The procedure for writing these settings to the EEPROM is described in the 9th example on $\ \, \stackrel{\textstyle \hookrightarrow}{\square} \ \,$ 45...47.

Replacing the EEPROM is described on \$\bigcap\$ 66 for BIS C-6001 and on \$\bigcap\$ 77 for BIS C-6021.

BALLUFF E 33



34

Function Description Processing code tags

Read from code tag, with program

The command identifier 21_{Hex} can be used to read out the program records stored in the program from the code tag. The user must document exactly which data are to be read from where and with what number of bytes for the respective program (see example 10 on \square 48)

Write to code tag, with program

The command identifier 22_{Hex} can be used to write the program records stored in the program to the code tag. The user must document exactly which data are to be written from where and with what number of bytes for the respective program (see example 11 on \square 49)

Function Description Examples for protocol sequence

Example No. 1

Initializing the code tag for the CRC-16 data checking

For configuring with double bit header!

The processing of this command is similar to a write command. Start address and number of bytes have to correspond to the maximum number of data to be used. In this example the complete memory range of a code tag with 128 bytes shall be used (BIS C-1__03/L with 32 byte block size). Because 2 bytes are used for the CRC only 120 bytes can be used as data bytes, hence: start address = 0, number of bytes = 120.

BIS C-60_1 Identification System: 2.) Process subaddresses of the input buffer in order

1.) Process subaddresses of the output buffer in order shown:

| 01 _{Hex} | Command designator 12Hex | |
|-------------------|---------------------------------|--|
| 02 _{Hex} | Start address 00 _{Hex} | |
| 03 _{Hex} | Start address 00 _{Hex} | |
| 04 _{Hex} | No. of bytes 78 _{Hex} | |
| 05 _{Hex} | No. of bytes 00 _{Hex} | |
| 00Hex/07Hex | Set AV-Bit, CT-Bit to 0 | |

Set AA-Bit, invert TO-Bit 00_{Hex}/07_{Hex}

3.) Process subaddresses of the output buffer:

| - | · |
|--------------------------------------|-----------------------------|
| 0106 _{Hex} | Enter first 6 bytes of data |
| 00 _{Hex} /07 _{Hex} | Invert TI-Bit |

4.) Process subaddresses of the output buffer:

| 0106 _{Hex} | Copy first 6 data bytes | | |
|--|-------------------------|--|--|
| Process subaddress of the input buffer: | | | |
| 00 _{Hex} /07 _{Hex} Invert TO-Bit | | | |

...To be continued until the complete memory range is written. See next [].

5.) Process subaddresses of the output buffer: 01...06_{Hex} Enter the second 6 data bytes Invert TI-Bit

6.) Process subaddresses of the output buffer: Copy second 6 data bytes Process subaddress of the input buffer 00нех/07нех Invert TO-Bit

BALLUFF E 35



36

Function Description Examples for protocol sequence

41.)Process subaddresses of the output buffer:

| 42.)Process | subaddresses | of | the | output | buffer: |
|-------------|--------------|----|-----|--------|---------|
|-------------|--------------|----|-----|--------|---------|

| Example No. 1 |
|---------------|
| (continued) |
| |

| | Enter the remaining data byte |
|--------------------------------------|-------------------------------|
| 00 _{Hex} /07 _{Hex} | Invert TI-Bit |

01...06_{Hex} Copy the remaining data byte Process subaddress of the input buffer 00Hex/07Hex Set AE-Bit

For configuring with double bit header!

43.) Process subaddresses of the output buffer:

00_{Hex}/07_{Hex} Reset AV-Bit

44.)Process subaddresses of the input buffer:

00_{Hex}/07_{Hex} Reset AA-Bit and AE-Bit

36 E



Function Description Examples for protocol sequence

Example No. 2

Read 17 bytes starting at code tag address 10 (code tag type with 32 byte block size): BIS C-60_1 Identification System: Host:

For configuring with double bit header!

1.) Process subaddresses of the output buffer in order shown

Command designator 01 He 02_{Hex} Start address Low Byte 0A He Start address High Byte 00 He 03_{He} No. of bytes Low Byte 11 He 04_{Hex} No. of bytes High Byte 00 00_{Hex}/07_{Hex} CT-Bit to 0 (32 Byte block size), 2.) Process subaddresses of the input buffer in order

| 00нех/07нех | Set AA-Bit |
|---------------------|-----------------------------|
| 0106 _{Hex} | Enter first 6 bytes of data |
| 00Hex/07Hex | Set AE-Bit |

3.) Process subaddresses of the input buffer:

01...06_{Hex} Copy first 6 data bytes Process subaddress of the output buffer: 00_{Hex}/07_{Hex} Invert TI-Bit

4.) Process subaddresses of the input buffer:

| 0106Hex | Enter the second 6 data bytes |
|--------------------------------------|-------------------------------|
| 00 _{Hex} /07 _{Hex} | Invert TO-Bit |

5.) Process subaddresses of the input buffer:

01...06_{Hex} Copy second 6 data bytes Process subaddress of the output buffer: 00_{Hex}/07_{Hex} Invert TI-Bit

6.) Process subaddresses of the input buffer:

| - | · |
|--------------------------------------|----------------------------------|
| | Enter the remaining 5 data bytes |
| 00 _{Hex} /07 _{Hex} | Invert TO-Bit |

7.) Process subaddresses of the input buffer:

01...05Hex Copy the remaining 5 data bytes Process subaddress of the output buffer: 00_{Hex}/07_{Hex} Reset AV-Bit

8.) Process subaddresses of the input buffer:

00_{Hex}/07_{Hex} Reset AA-Bit and AE-Bit



BALLUFF E 37

38

Function Description Examples for protocol sequence

Example No. 3 with simultaneous data transmission

Read 17 bytes starting at code tag address 10, with simultaneous data transmission like 2nd example but (code tag type with 32 byte block size):

While the read job is being carried out and as soon as the input buffer is filled, the first data are sent. The AE bit is not set until the "Read" operation is completed by the processor.

For configuring with double bit header!

The reply "Job End" = AE bit is reliably set no later than before the last data are sent. The exact time depends on the requested data amount, the input buffer size and the timing of the controller. This is indicated in the following by the note Set AE-Bit (in italics).

Host:

1.) Process subaddresses of the output buffer in order shown:

| 01 _{Hex} | Command designator 01 _{Hex} |
|--------------------------------------|---|
| 02 _{Hex} | Start address Low Byte 0A Hex |
| 03 _{Hex} | Start address High Byte 00 Hex |
| 04 _{Hex} | No. of bytes Low Byte 11 Hex |
| 05 _{Hex} | No. of bytes High Byte 00 Hex |
| 00 _{Hex} /07 _{Hex} | CT-Bit to 0 (32 Byte block size), set AV-Bit |

BIS C-60 1 Identification System:

2.) Process subaddresses of the input buffer in order shown:

| 00 _{Hex} /07 _{Hex} | Set AA-Bit |
|--------------------------------------|-----------------------------|
| 0106 _{Hex} | Enter first 6 bytes of data |
| $00_{\text{Hex}}/07_{\text{Hex}}$ | Invert TO-Bit |
| $00_{\text{Hex}}/07_{\text{Hex}}$ | Set AE-Bit |

3.) Process subaddresses of the input buffer:

| 0106 _{Hex} | Copy first 6 data bytes |
|---------------------|----------------------------------|
| Process s | subaddress of the output buffer: |
| 00/07 | Invest TI Dit |

4.) Process subaddresses of the input buffer:

| 0106 _{Hex} | Enter the second 6 data bytes |
|--------------------------------------|-------------------------------|
| 00 _{Hex} /07 _{Hex} | Invert TO-Bit |
| 00нех/07нех | Set AE-Bit |

Continued on next \Box .

38 E



Function Description Examples for protocol sequence

Example No. 3

like 2nd example but with simultaneous data transmission

For configuring with double bit header!

5.) Process subaddresses of the input buffer:

| 0106Hex | Copy second 6 data bytes |
|--|--------------------------|
| Process subaddress of the output buffer: | |
| 00Hex/07Hex | Invert TI-Bit |
| | |

7.) Process subaddresses of the input buffer:

| 0105 _{Hex} | Copy the remaining 5 data bytes |
|--|---------------------------------|
| Process subaddress of the output buffer: | |
| 00Hex/07Hex | Reset AV-Bit |

6.) Process subaddresses of the input buffer:

| 0105 _{Hex} | Enter the remaining 5 data bytes |
|--------------------------------------|----------------------------------|
| 00 _{Hex} /07 _{Hex} | Invert TO-Bit |
| 00 _{Hex} /07 _{Hex} | Set AE-Bit |

8.) Process subaddresses of the input buffer:

00_{Hex}/07_{Hex} Reset AA-Bit and AE-Bit

BALLUFF E 39



40

Function Description Examples for protocol sequence

Example No. 4

Read 30 bytes starting at code tag address 10 with read error (code tag type with 64 byte block size):

For configuring with double bit header!

| 01 _{Hex} | Command designator 01 _{Hex} |
|--------------------------------------|---|
| 02Hex | Start address Low Byte 0AHex |
| 03 _{Hex} | Start address High Byte 00 Hex |
| 04 _{Hex} | No. of bytes Low Byte 1E _{Hex} |
| 05 _{Hex} | No. of bytes High Byte 00 Hex |
| 00 _{Hex} /07 _{Hex} | Set CT-Bit to 1 (64 Byte block size), set AV-Bit |

BIS C-60_1 Identification System:

1.) Process subaddresses of the output buffer in the order shown:

2.) Process subaddresses of the input buffer in the order shown:

If an error occurs right away:

| $00_{\text{Hex}}/07_{\text{Hex}}$ | Set AA-Bit |
|--------------------------------------|--------------------|
| 01 _{Hex} | Enter error number |
| 00 _{Hex} /07 _{Hex} | Set AF-Bit |

3.) Process subaddress of the input buffer:

| 01 _{Hex} | Copy error number |
|---|-------------------|
| Process subaddress of the output buffer: | |
| 00 _{Hex} /07 _{Hex} Reset AV-Bit | |
| | |

4.) Process subaddresses of the input buffer:

00_{Hev}/07_{Hex} Reset AA-Bit and AF-Bit

Function Description Examples for protocol sequence

Example No. 5, like 4th example but with simultaneous

data transmission

For configuring with double bit header!

Read 30 bytes starting at code tag address 10, with read error and simultaneous data transmission (code tag type with 64 byte block size):

If an error occurs, the AF bit is set instead of the AE-Bit, with a corresponding error number. When the AF-BIT is set the job is interrupted and declared to be ended.

Host:

1.) Process subaddresses of the output buffer in the order shown:

2.) Process subaddresses of the input buffer in the order shown:

| 01 _{Hex} | Command designator 01Hex |
|--------------------------------------|---|
| 02 _{Hex} | Start address Low Byte 0A _{Hex} |
| 03Hex | Start address High Byte 00 Hex |
| 04 _{Hex} | No. of bytes Low Byte 1E _{Hex} |
| 05Hex | No. of bytes High Byte 00 Hex |
| 00 _{Hex} /07 _{Hex} | Set CT-Bit to 1 (64 Byte block size), set AV-Bit |

3.) Process subaddress of the input buffer:

| 01 _{Hex} | Copy error number |
|--|-------------------|
| Process subaddress of the output buffer: | |
| $00_{\text{Hex}}/07_{\text{Hex}}$ | Reset AV-Bit |

BIS C-60_1 Identification System:

If an error occurs right away:

| I | 00Hex/07Hex | Set AA-Bit |
|---|--------------------------------------|--------------------|
| | 01Hex | Enter error number |
| ſ | 00 _{Hex} /07 _{Hex} | Set AF-Bit |

4.) Process subaddresses of the input buffer:

| Reset AA-Bit and AF-Bit |
|-------------------------|



An error can also occur after the data have already been sent (see example on the next 🗅).

BALLUFF E 41



42

Function Description Examples for protocol sequence

Example No. 6

Read 30 bytes starting at code tag address 10, with read error and simultaneous data transmission (code tag type with 64 byte block size):

For configuring with double bit header and 8-byte buffer size!

If an error occurs after data have started to be sent, the AF-Bit is set instead of the AE-Bit along with the corresponding error number. The error message AF is dominant. It cannot be specified which data are incorrect. When the AF-BIT is set the job is interrupted and declared to be ended.

BIS C-60_1 Identification System:

1.) Process subaddresses of the output buffer in the order shown:

| 01 _{Hex} | Command designator 01 _{Hex} |
|-------------------|---|
| 02 _{Hex} | Start address Low Byte 0A _{Hex} |
| 03 _{Hex} | Start address High Byte 00 Hex |
| 04 _{Hex} | No. of bytes Low Byte 1E _{Hex} |
| 05 _{Hex} | No. of bytes High Byte 00 Hex |
| 00нех/07нех | Set CT-Bit to 1 (64 Byte block size), set AV-Bit |

3.) Process subaddress of the input buffer:

| 0106Hex | Copy first 6 data bytes |
|--|-------------------------|
| Process subaddress of the output buffer: | |
| 00 _{Hex} /07 _{Hex} | Invert TI-Bit |

5.) Process subaddress of the input buffer:

| 01 _{Hex} | Copy error number |
|--|-------------------|
| Process subaddress of the output buffer: | |
| 00Hex/07Hex | Reset AV-Bit |

2.) Process subaddresses of the input buffer in the order shown:

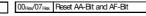
| 00 _{Hex} /07 _{Hex} | Set AA-Bit |
|--------------------------------------|------------------------------|
| | Enter the first 6 data bytes |
| 00 _{Hex} /07 _{Hex} | Invert TO-Bit |
| | |

4.) Process subaddresses of the input buffer: If an error has occurred:

| 01 _{Hex} | Enter error number |
|-------------------|--------------------|

| 01 _{Hex} | Enter error number |
|--------------------------------------|--------------------|
| 00 _{Hex} /07 _{Hex} | Set AF-Bit |

6.) Process subaddresses of the input buffer:



Function Description Examples for protocol sequence

Example No. 7

For configuring with double bit header!

Write 16 bytes starting at code tag address 20 (code tag type with 32 byte block size): BIS C-60_1 Identification System: Host:

1.) Process subaddresses of the output buffer in the

| order shown: | | |
|--------------------------------------|---|----|
| 01 _{Hex} | Command designator 02 _{Hex} | 00 |
| 02 _{Hex} /03 _{Hex} | Start address 14 _{Hex} / 00 _{Hex} | |
| 04 _{Hex} /05 _{Hex} | No. of bytes 10 _{Hex} / 00 _{Hex} | |
| 00 _{Hex} /07 _{Hex} | CT-Bit to 0 (32 Byte block size), set AV-Bit | |

2.) Process subaddresses of the input buffer in the

OHex/07Hex Set AA-Bit, invert TO-Bit

3.) Process subaddresses of the output buffer:

Enter the first 6 data bytes 00Hex/07Hex Invert TI-Bit

4.) Process subaddresses of the output buffer:

01...06Hex Copy the first 6 data bytes Process subaddress of the input buffer: 00_{Hex}/07_{Hex} Invert TO-Bit

5.) Process subaddresses of the output buffer:

01...06_{Hex} Enter the second 6 data bytes Invert TI-Bit

6.) Process subaddresses of the output buffer:

01...06Hex Copy the second 6 data bytes Process subaddress of the input buffer: 00Hex/07Hex Invert TO-Bit

7.) Process subaddresses of the output buffer:

01...04_{Hex} Enter the remaining 4 data bytes 00_{Hex}/07_{Hex} Invert TI-Bit

8.) Process subaddresses of the output buffer:

01...04_{Hex} Copy the remaining 4 data bytes Process subaddress of the input buffer: 00нех/07нех Set AE-Bit

9.) Process subaddresses of the output buffer:

00_{Hex}/07_{Hex} Reset AV-Bit

10.)Process subaddresses of the input buffer: 00_{Hex}/07_{Hex} Reset AA-Bit and AE-Bit

BALLUFF E 43



44

Function Description Examples for protocol sequence

Example No. 8 with Auto-Read function

For configuring with double bit header!

Programming start address 50 (code tag type with 32 byte block size):

1.) Process subaddresses of the output buffer in the order shown:

| 01 _{Hex} | Command designator 07Hex |
|-------------------|---|
| 02 _{Hex} | Start address Low Byte 32 _{Hex} |
| 03 _{Hex} | Start address High Byte 00 Hex |
| | CT-Bit to 0 (32 Byte block size), set AV-Bit |

BIS C-60 1 Identification System:

2.) Process subaddresses of the input buffer: 00Hex/07Hex Set AA-Bit and AE-Bit

3.) Process subaddresses of the output buffer:

00Hex/07Hex Reset AV-Bit

4.) Process subaddresses of the input buffer:

00Hex/07Hex Reset AA-Bit and AE-Bit



Function Description Examples for protocol sequence

Example No. 9 Mixed Data Access

Storing a program for reading out 3 data records:

For configuring with double bit header!

5 1st data record Start address Number of bytes Start address Start address Number of bytes Number of bytes 75 312 3 17 2nd data record 3rd data record

Total number of bytes exchanged in the operation: 27 bytes

All 104 bytes are written for the programming.

Host:

1.) Process subaddresses of the output buffer in the 2.) Process subaddresses of the input buffer:

| 01 _{Hex} | Command designator 06 _{Hex} |
|-------------------|--|
| 02 _{Hex} | Program number 01 _{Hex} |
| | CT-Bit to 0 or 1 (depending on block size), set AV-Bit |

3.) Process subaddresses of the output buffer:

| 01 _{Hex} | 1st start address | (Low Byte) 05 _{Hex} |
|-------------------|---------------------|-------------------------------|
| 02 _{Hex} | | (High Byte) 00 _{Hex} |
| 03нех | 1st number of bytes | (Low Byte) 07Hex |
| 04 _{Hex} | | (High Byte) 00 _{Hex} |
| 05 _{Hex} | 2nd start address | (Low Byte) 4B _{Hex} |
| 06 _{Hex} | | (High Byte) 00 _{Hex} |
| 00um/07um | Invert TI-Rit | |

00_{Hex}/07_{Hex} Set AA-Bit, invert TO-Bit

4.) Process subaddresses of the input buffer:

| 00 _{Hex} /07 _{Hex} | Invert TO-Bit |
|--------------------------------------|---------------|
| | |
| | |
| | |
| | |
| | |

Host:

Continued on next].

BALLUFF E 45



46

Function Description Examples for protocol sequence

Example No. 9 Mixed Data Access (continued)

For configuring with double bit header!

5.) Process subaddresses of the output buffer:

| 01 _{Hex} 02 _{Hex} | 2nd number of bytes | (Low Byte) 03 _{Hex} (High Byte) 00 _{Hex} |
|--|------------------------|---|
| 03 _{Hex} | 3rd start address | (Low Byte) 38 _{Hex} |
| 04 _{Hex} | | (High Byte) 01 Hex |
| 05 _{Hex} | 3rd number of | (Low Byte) 11 _{Hex} |
| 06 _{Hex} | bytes | (High Byte) 00 _{Hex} |
| 00 _{Hex} /07 _{Hex} | Invert TI-Bit | |

- 7.) Process subaddresses of the output buffer:
- 01_{Hex}/02_{Hex} Terminator FF_{Hex}/FF_{Hex} 03_{Hex}/04_{Hex} (not used) FF_{Hex}/FF_{Hex} 05_{Hex}/06_{Hex} (not used) 00_{Hex}/07_{Hex} Invert TI-Bit

6.) Process subaddresses of the input buffer:

00_{Hex}/07_{Hex} Invert TO-Bit

8.) Process subaddresses of the input buffer:

00_{Hex}/07_{Hex} Invert TO-Bit

Fill all unused start addresses and number of bytes with FFHex!

Continued on next].



Function Description Examples for protocol sequence

Example No. 9 Mixed Data Access (continued)

35.)Process subaddresses of the output buffer:

36.)Process subaddresses of the input buffer:

| - | | - |
|--------------------------------------|--------------------|--------------------------------------|
| 01Hex/02Hex | (nicht verwendet) | FFHex/FFHex |
| 03 _{Hex} /04 _{Hex} | (nicht verwendet) | FF _{Hex} /FF _{Hex} |
| 05Hex/06Hex | (nicht verwendet) | FFHex/FFHex |
| 00 _{Hex} /07 _{Hex} | TI-Bit invertieren | |

00нех/07нех AE-Bit setzen

For configuring with double bit header!

37.)Process subaddresses of the output buffer:

38.)Process subaddresses of the input buffer:

00_{Hex}/07_{Hex} AV-Bit rücksetzen

00Hex/07Hex AA-Bit und AE-Bit rücksetzen



We recommend that you carefully document which parameters are used for start addresses and number of bytes for writing/reading the desired data records.

The data are sequenced in the exact order specified in the program.

BALLUFF E 47



48

Function Description Examples for protocol sequence

Example No. 10 Mixed Data Access For configuring with double bit header!

Read code tag using Program No. 1 (code tag type with 32 byte block size):

Host:

1.) Process subaddresses of the output buffer in the 2.) Process subaddresses of the input buffer in the

| | order shown. | |
|-------|---|--|
| 01Hex | Command designator 21Hex | |
| 02Hex | Program number 01 _{Hex} | |
| | CT-Bit to 0 (32 byte block size), set AV-Bit | |

BIS C-60_1 Identification System:

| Order Shown. | |
|--------------------------------------|-----------------------------|
| 00 _{Hex} /07 _{Hex} | Set AA-Bit |
| 0106Hex | Enter first 6 bytes of data |
| 00 _{Hex} /07 _{Hex} | Set AE-Bit |

| 0., 11000033 3 | abadaresses of the input baller. |
|---------------------|----------------------------------|
| 0106 _{Hex} | Copy first 6 data bytes |
| Process s | subaddress of the output buffer: |
| 000/07 | Invert TI_Rit |

4.) Process subaddresses of the output buffer:

| 0106 _{Hex} | Enter the second 6 data bytes |
|--------------------------------------|-------------------------------|
| 00 _{Hex} /07 _{Hex} | Invert TO-Bit |

... A total of 27 bytes of data are exchanged. For the remainder of the procedure, see Example 2 on $\hfill 37.$



Dynamic mode is turned off while the Mixed Data Access program is being run.

Function Description Examples for protocol sequence

Example No. 11 Mixed Data Access

Write code tag using Program No. 1 (code tag type with 32 byte block size):

Host:

For configuring with double bit header!

BIS C-60_1 Identification System: 1.) Process subaddresses of the output buffer in the 2.) Process subaddresses of the input buffer in the

order shown:

| 01 _{Hex} | Command designator 22 _{Hex} |
|-------------------|---|
| 02 _{Hex} | Program number 01 _{Hex} |
| | CT-Bit to 0 (32 byte block size), set AV-Bit |

| order enemn | |
|--------------------------------------|-----------------------------|
| 00нек/07нек | Set AA-Bit |
| 0106 _{Hex} | Enter first 6 bytes of data |
| 00 _{Hex} /07 _{Hex} | Set AE-Bit |

3.) Process subaddresses of the input buffer:

| 0106 _{Hex} | Copy first 6 data bytes |
|--|-------------------------|
| Process subaddress of the output buffer: | |
| 00 _{Hex} /07 _{Hex} Invert TI-Bit | |

4.) Process subaddresses of the input buffer:

| 0106 _{Hex} | Enter the second 6 data bytes |
|--------------------------------------|-------------------------------|
| 00 _{Hex} /07 _{Hex} | Invert TO-Bit |

... A total of 27 bytes of data are exchanged.
For the remainder of the procedure, see Example 7 on 143.



Dynamic mode is turned off while the Mixed Data Access program is being run.

BALLUFF E 49



50

Function Description Examples for protocol sequence

Example No. 12

Put the relevant read/write head into ground state:

For configuring with double bit header!

Both read/write heads can be independently set to the ground state.

BIS C-60_1 Identification System:

1.) Process subaddresses of the output buffer:

2.) Go to ground state; Process subaddresses of the input buffer:

| 00um/07um | Set GR-Rit |
|-----------|------------|

00_{Hex}/07_{Hex} Reset BB-Bit

4.) Process subaddresses of the input buffer:

3.) Process subaddresses of the output buffer:

00_{Hex}/07_{Hex} Reset GR-Bit

00Hex/07Hex Set BB-Bit

Function Description Examples for protocol sequence

Example No. 13

Program configuration data:

Configuration data can be programmed in both buffers - for Head 1 and Head 2 - as desired.

For configuring with double bit header!

BIS C-60_1 Identification System:

1.) Process subaddresses of the output buffer in the order shown:

2.) Process subaddresses of the input buffer in the order shown.

| | Command designator 04 _{Hex} |
|--------------------------------------|--------------------------------------|
| 00 _{Hex} /07 _{Hex} | Set AV-Bit |

3.) Process subaddresses of the output buffer:

| 1 | 0106 _{Hex} | Enter the 6 configuration bytes |
|---|--------------------------------------|---------------------------------|
| | 00 _{Hex} /07 _{Hex} | Invert TI-Bit |

4.) Process subaddresses of the output buffer:

00_{Hex}/07_{Hex} Set AA-Bit, invert TO-Bit

| 0106 _{Hex} Enter the 6 configuration bytes | | |
|---|------------|--|
| Process subaddresses of the input buffer | | |
| 00 _{Hex} /07 _{Hex} | Set AE-Bit | |

| | | subaddresses of the output buller. | |
|----|-----------|------------------------------------|---|
| 00 | Hey/07Hey | Reset AV-Bit | _ |

6.) Process subaddresses of the input buffer:

00Hev/07Hex Reset AA-Bit and AE-Bit

BALLUFF E 51



52

Function Description Examples for protocol sequence

Example No. 14

Read-out programmed configuration data:

For configuring with double bit header!

Host:

BIS C-60_1 Identification System:

1.) Process subaddresses of the output buffer in the order shown:

2.) Process subaddresses of the input buffer in the order shown:

| 01 _{Hex} | Command designator 05 _{Hex} |
|-------------------|--------------------------------------|
| 00Hex/07Hex | Set AV-Bit |

| 00 _{Hex} | Set AA-Bit |
|-------------------|---------------------------------|
| 0106Hex | Enter the 6 configuration bytes |
| 00 _{Hex} | Set AE-Bit |

3.) Process subaddresses of the input buffer::

| - | • |
|---------------------|---------------------------------|
| 0106 _{Hex} | Copy the 6 configuration bytes |
| | ubaddress of the output buffer: |
| 00Hex/07Hex | Reset AV-Bit |

4.) Process subaddresses of the output buffer:

00Hex/07Hex Reset AA-Bit and AE-Bit

Read/Write Times

Read times from code tag to processor in processor in static mode (parametering: 2nd byte, bit 5 = 0, without CRC-16 data

check) Write times from processor to code tag in static mode

(parametering: 2nd byte, bit 5 = 0, without CRC-16 data

check)

For double read and compare:

| Code tag with 32 byte blocks | | |
|-------------------------------------|----------------|--|
| No. of bytes | Read time [ms] | |
| from 0 to 31 | 110 | |
| for each additional 32 bytes add | 120 | |
| from 0 to 255 | = 950 | |

Code tag with 64 byte blocks No. of bytes Read time [ms] from 0 to 63 220 for each additional 64 bytes add 230 from 0 to 2047 = 7350

Including readback and compare:

| Code tag with 32 byte blocks | | |
|------------------------------|------------------|--|
| No. of bytes | Write time [ms] | |
| from 0 to 31 | 110 + n * 10 | |
| for 32 bytes or more | y * 120 + n * 10 | |

| ocks |
|------------------|
| Write time [ms] |
| 220 + n * 10 |
| y * 230 + n * 10 |
| |

n = number of contiguous bytes to write y = number of blocks to be processed

Example: 17 bytes from address 187 have to be written. Code tag with 32 bytes per block. The blocks 5 and 6 will be processed since the start address 187 is in block 5 and the end address 203 in block 6. t=2*120+17*10=410 ms



The indicated times apply after the code tag has been recognized. If the code tag is not yet recognized, an additional 45 ms for building the required energy field until the code tag is recognized must be added.

BALLUFF E 53



54

Read/Write Times

Read times from code tag to processor in dynamic mode (parametering: 2nd byte, bit 5 = 1, without CRC-16 data check)

Read times within the 1st block for dual read and compare:

The indicated times apply after the code tag has been recognized. If the code tag is not yet recognized, an additional 45 ms for building the required energy field until the code tag is recognized must be added.

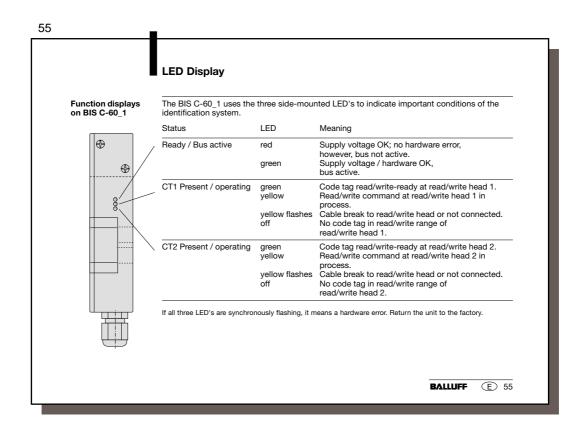
| Code tag with 32 byte blocks | |
|------------------------------|----------------|
| No. of bytes | Read time [ms] |
| from 0 to 3 | 14 |
| for each additional | |
| byte add | 3.5 |
| from 0 to 31 | 112 |

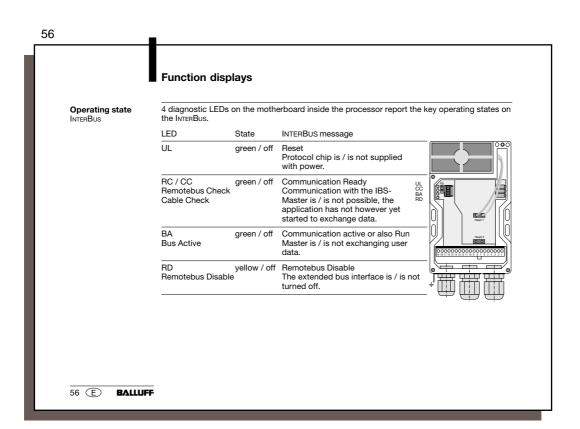
| Code tag with 64 byte blocks | |
|------------------------------|----------------|
| No. of bytes | Read time [ms] |
| from 0 to 3 | 14 |
| for each additional | |
| byte add | 3.5 |
| from 0 to 63 | 224 |

m = highest address to be read

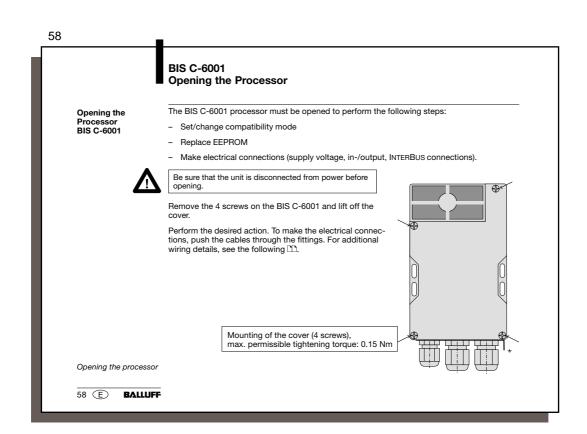
Formula: t = (m + 1) * 3.5 ms

Example: Read 11 bytes starting at address 9, i.e. the highest address to be read is 19. This corresponds to 70 ms.





57 BIS C-6001 **Mounting Head / Processor** Depending on model, the processor is equipped with a read/write head or the adapter for offset read/write heads. Both the read/write head and the adapter can be rotated by the user by + or -90 deg. to the desired Mounting the read/write head or adapter position (see drawing). Be sure that power is off first. Loosen Caution: both screws (indicated with arrows). Carefully pull the head wires inside! or adapter out towards the side (direction of arrow, right drawing). Caution: wires inside! 8 Reattach at the desired orientation and screw tight again. 16.8 The processor is attached Mounting the BIS C-6001 using 4 lateral mounting holes. 46.6 BALLUFF **E** 57



BIS C-6001 Installing the connection cables

Make connections on the BIS C-6001 processor

The BIS C-6001 processor must be opened in order to make the connections for the supply voltage, the digital input and the INTERBUS connections (see \square 58).

First be sure that the unit is disconnected from power.

Remove the 4 screws on the BIS C-6001 and lift off the cover.

Guide the two INTERBUS cables through the PG 11 fittings (see 🗋 60). For additional information on wiring, see the following 11.

Push the cable for supply voltage and for the digital input through the PG 9 fitting.

Close up the processor.

If the processor is equipped with an adapter:

- BIS C-650: Connect the read/write heads to terminals Head 1 and Head 2.
 BIS C-670: Connect the read/write head to terminal Head 1.

BALLUFF E 59

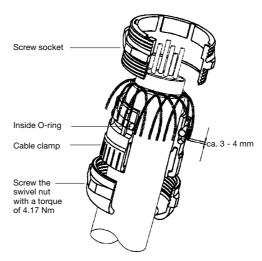


60

BIS C-6001 **Mounting the PG Connection**

Mounting the PG Connection on the processor BIS C-6001

After connecting the (field) bus leads to the termional block, make sure that the shield has proper connection to the PG housing.

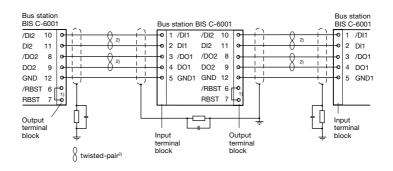


60 E

BIS C-6001 Interface Information / Wiring Diagrams

Remote bus cable and interfaces for INTERBUS

To insert the BIS C-6001 processors into the serial INTERBUS, the terminal strip provides terminals 1...5 for the input interface and terminals 8...12 for the output interface of the INTERBUS. The following drawing shows the wiring when the BIS C-6001 processors need to be connected together.



- Leave the jumper in the BIS C-6001 if an additional station is to follow. Remove it if no additional station follows.
- 2) The differential signals DO and /DO as well as DI and /DI must be twisted-pair. Recommended cable: LiYCY 3x2x0.25 mm² (AWG 24); maximum cable capacitance: 120 pF/m

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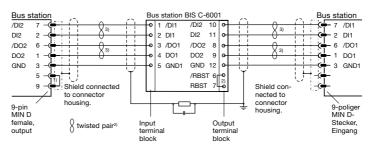


62

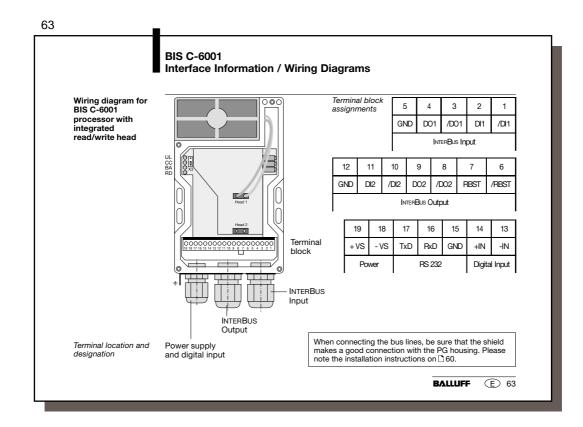
BIS C-6001 Interface Information / Wiring Diagrams

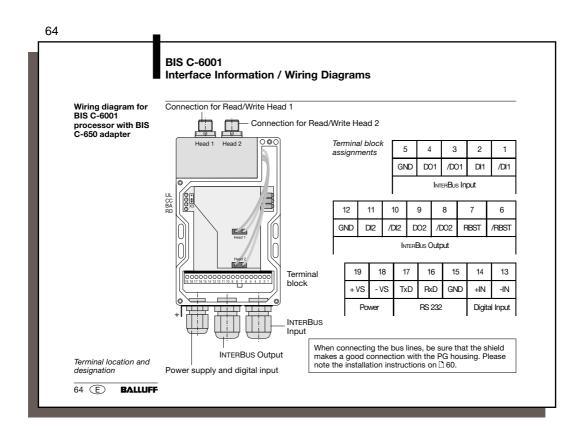
Remote bus cable and interfaces for INTERBUS (continued)

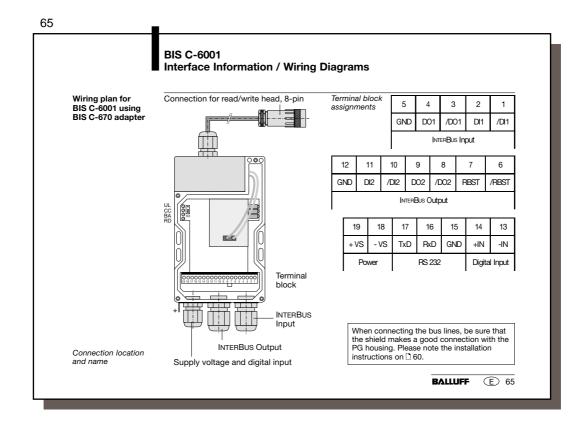
To insert the BIS C-6001 processors into the serial INTERBUS, the terminal strip provides terminals 1...5 for the input interface and terminals 1...5 for the output interface of the INTERBUS. The following drawing shows the wiring when the BIS C-6001 interface needs to be connected using a 9-pin terminal (e.g. to a BIS C-6021).

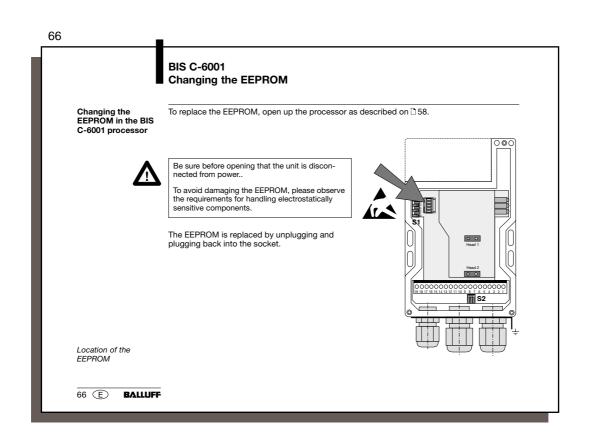


- Connect the jumper in the connector if another station is to follow. Remove it if no additional station follows.
- 2) The jumper remains in the BIS C-601 if an another station is to follow. Remove it if no additional station follows.
- 3) The differential signals DO and /DO as well as DI and /DI must be twisted-pair. Recommended cable: LiYCY 3x2x0.25 mm² (AWG 24); maximum cable capacitance: 120 pF/m









BIS C-6001

Technical Data Housing
Dimensions with read/write head BIS C-652
Dimensions with adapter BIS C-650 Plastic PS Dimensions, Weight 169 x 90 x 35 mm 184 x 90 x 35 mm Weight 400 g 0 °C to +60 °C Operating Ambient Temperature Conditions 19-pin 1 x PG 9 fitting (metal) Connections Terminal Block Cable Entry for supply voltage 2 x PG 11 fittings (metal) 4 to 8 mm for PG 9 5 to 10 mm for PG 11 for INTERBUS, in-/output Cable Diameter Wire gauge 0.14 to 1 mm² 0.25 to 0.34 mm² with end crimps **Enclosure Rating Enclosure Rating** IP 65 (when connected) Supply Voltage V_s, input DC 24 V ± 20 % Electrical Ripple Current Draw ≤ 10 % ≤ 400 mA Connections INTERBUS, In-/ and Outputs serial interface for remote bus station, Ident-No. 03, 16 bytes IN, 16 bytes OUT Digital Input (+IN, -IN) Control voltage active Terminal block, Optocoupler isolated 4 V to 40 V Control voltage inactive Input current at 24 V 1.5 V to -40 V 11 mA +IN-Delay time, typ. 5 ms

68

BIS C-6001 **Technical Data**

Electrical Connections (continued)

Function Displays

Service interface

RS 232

Read/Write Head*

option for mounted adapter BIS C-650*

integrated, BIS C-65_ and following; 2 x connectors 4-pin (male) for all read/write heads BIS C-3 with 4-pin connector (female), not BIS C-350 and BIS C-352 1 x connector 4-pin (male) for read/write heads BIS C-350 / BIS C-352

BALLUFF E 67

option for mounted adapter BIS C-670*

BIS operating states (LED in housing)

Ready / Bus active CT1 Present / Operating

INTERBUS state (LED on side of housing)

LED red / green LED green / yellow LED green / yellow CT2 Present / Operating LED green Reset

LED green Cable Check LED green LED yellow Bus active Remotebus Disable



The CE-Mark is your assurance that our products are in conformance with the EC-Guideline

89/336/EEC (EMC-Guideline)

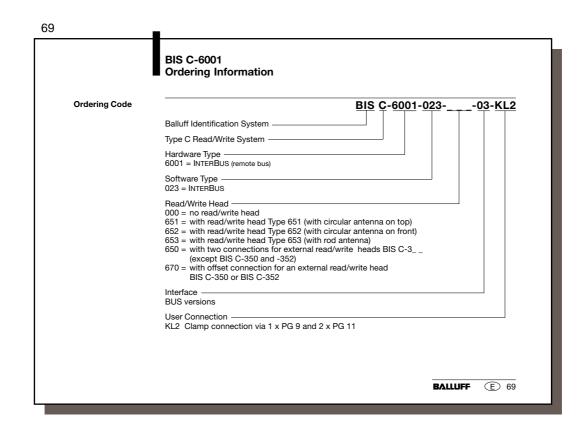
and the EMC Law. Testing in our EMC Laboratory, which is accredited by the DATech for Testing of Electromagnetic Compatibility, has confirmed that Balluff products meet the EMC requirements of the Generic Standard

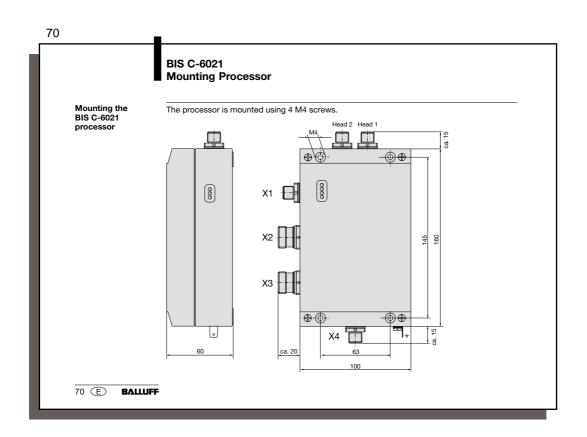
EN 50081-2 (Emission) and EN 50082-2 (Noise Immunity).

68 E

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* rotatable by 90 degrees





71 BIS C-6021 Opening the processor / Interface information To convert the processor for the power supply via the installation remote bus instead of sup-Opening the BIS C-6021 plying it via X1, the internal connections have to be changed. processor Ensure that the device is turned off. Remove the 4 screws on the BIS C-6021 and lift off the cover. Conversion see the following 🗀. Connection diagram for BIS C-6021 Head 2 Head 1 Connection for read/write head 2 Connection for read/write head 1 **(**) **⊕** ⊕ opening that the unit is disconnected from power. Digital input, X1 power supply for ST9 remote bus version INTERBUS output INTERBUS input Mounting of the cover (4 screws), • • • **⊕**⊕ max. permissible tightening torque: 0.15 Nm Service interface X4 Connection locations Function ground FE and names BALLUFF E 71

BIS C-6021 Interface Information / Wiring Diagrams

To make the connections for the InterBus, the supply voltage and the digital input, connect the pre-assembled cable to the processor. For additional wiring information, see the following [12].

Connect the read/write heads to the terminals for Head 1 and Head 2.

Connecting on the remote bus or installation remote bus

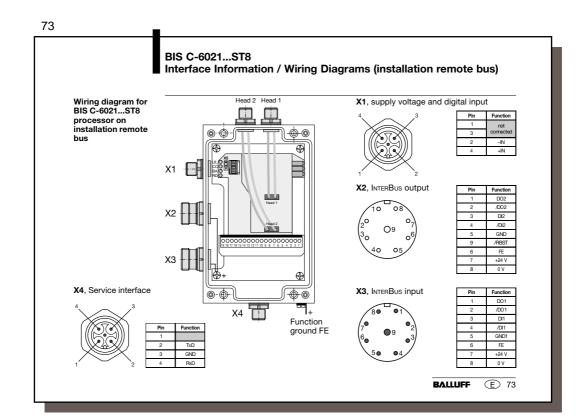
72

The BIS C-6021...ST8 processor is intended for use on the installation remote bus, which provides the supply voltage over the bus. The BIS C-6021...ST9 processor is intended for use on the remote bus, whereby the supply voltage for the processor is brought in through X1.

Connect the incoming INTERBUS cable to the INTERBUS input X3. Connect the outgoing INTERBUS cable to the INTERBUS output X2.

If this remote bus station is the last one on the bus, the INTERBUS output X2 must be closed off with a threaded cap to maintain the enclosure rating.

Please note the load capacity of the INTERBUS cable and verify during operation that the supply voltage is maintained at the processor (see Technical Data for specifications).



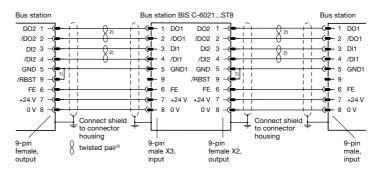
BIS C-6021...ST8

Wiring diagram for BIS C-6021...ST8 processor on installation remote bus (continued)

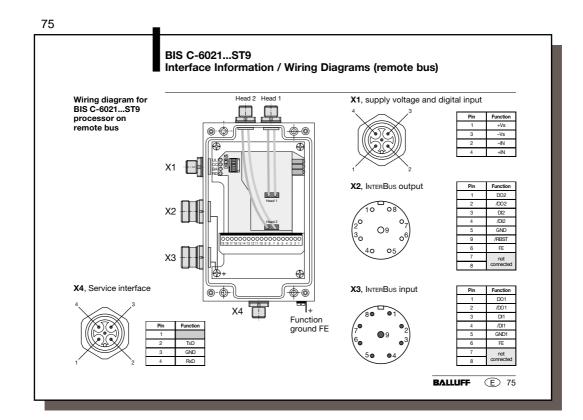
74

To integrate BIS C-6021 processor into the serial INTERBUS, two terminals are provided on the housing, X2 as INTERBUS output and X3 as INTERBUS input. For installation remote bus, the stations are supplied through the bus.

Interface Information / Wiring Diagrams (installation remote bus)



- Connect the jumper in the connector if another station is to follow. Remove it if no additional station follows.
- ²) The differential signals DO and /DO as well as DI and /DI must be twisted-pair. Recommended cable: LiYCY 3x2x0.25 mm² (AWG 24); maximum cable capacitance: 120 pF/m

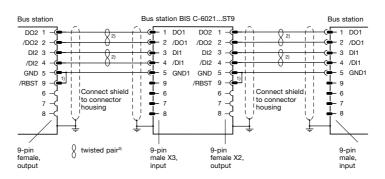


BIS C-6021...ST9 Interface Information / Wiring Diagrams (remote bus)

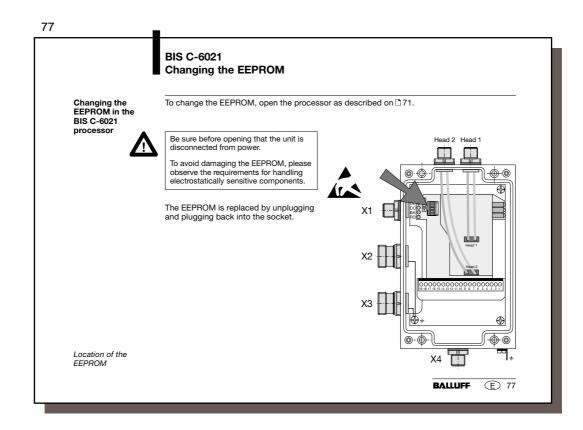
Wiring diagram for BIS C-6021...ST9 processor on remote bus (continued)

76

To integrate BIS C-6021 processor into the serial INTERBUS, two terminals are provided on the housing, X2 as INTERBUS output and X3 as INTERBUS input. For remote bus, the stations are not supplied through the bus.



- Connect the jumper in the connector if another station is to follow. Remove it if no additional station follows.
- ²) The differential signals DO and /DO as well as DI and /DI must be twisted-pair. Recommended cable: LiYCY 3x2x0.25 mm² (AWG 24); maximum cable capacitance: 120 pF/m



78 BIS C-6021 **Technical Data** Housing Dimensions, weight Metal Dimensions Weight 190 x 120 x 60 mm 820 g Operating conditions Ambient temperature $0\,^{\circ}\text{C}$ to +60 $^{\circ}\text{C}$ Connection type Integral connector X1 5-pin (male) Integral connector Head 1, Head 2 Round connector for X2 4-pin (male) 9-pin (female) 9-pin (male) 4-pin (male) Round connector for X3 Integral connector X4 Enclosure Protection class IP 65 (when connected) Electrical Supply voltage V_s DC 24 V ± 20 % connections Ripple Current draw ≤ 10 % ≤ 400 mA Connections for supply voltage \mathbf{V}_{S} at INTERBUS input X3, output X2 with installation remote bus with remote bus serial interface for remote bus station, Ident-No. 03, 16 bytes IN, 16 bytes OUT (with BIS C-621 mode: 8 Byte IN, 8 Byte OUT) Output X2, input X3, INTERBUS via integrated adapter with 2 x connectors for all read/write heads BIS C-3__ with 4-pin connector (female), excluding BIS C-350 and BIS C-352 Head 1, Head 2, Read/Write Head 78 E BALLUFF

79 BIS C-6021 **Technical Data** galvanically isolated (optocoupler) 4 V to 40 V 1.5 V to -40 V Digital input X1 (+IN, -IN) Electrical Control voltage active Control voltage inactive Input current at 24 V Delay time, typ Connections (continued) 11 mA 5 ms Service interface X4 RS 232 LED red / green LED green / yellow LED green / yellow BIS operating states (LED in housing) Ready / Bus active CT1 Present / Operating **Function Displays** CT2 Present / Operating LED green Reset Cable Check InterBus state (LED on side LED green LED green LED yellow Bus active Remotebus Disable of housing)

The CE-Mark is your assurance that our products are in conformance with the EC-Guideline 89/336/EEC (EMC-Guideline)

and the EMC Law. Testing in our EMC Laboratory, which is accredited by the DATech for Testing of Electromagnetic Compatibility, has confirmed that Balluff products meet the

EMC requirements of the Generic Standard EN 50081-2 (Emission) and EN 50082-2 (Noise Immunity).

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E 79

+IN T W

80 **BIS C-6021 Ordering Information** Ordering code BIS C-6021-023-050-03-ST_ Balluff Identification System Type C Read/Write System Hardware Type 6021 = metal housing, INTERBUS
(remote bus or installation remote bus) Software Type - 023 = INTERBUS Adapter 050 = with two connections for external read/write heads BIS C-3_ (except BIS C-350 and -352) Interface 03 = BUS versions User Connection

ST8 = Connector version (installation remote bus)

ST9 = Connector version (remote bus)

(2 round connector for power supply/digital input and service interface, 2 round connectors for INTERBUS) Accessory Type Ordering code (optional, not included) for X1 for X2 for X3 BKS-S79-00 Mating connector BKS-S83-00 BKS-S84-00 BKS-S 10-3 BES 12-SM-2 Mating connector for X4 Protective cap Protective cap for X1, Head 1, Head 2, X4 for X2 80 E BALLUFF

