

Empowering Embedded Systems

µC/OS-II μC/Probe

and The Freescale MC9S08 (Using the P&E DEMOQE Evaluation Board)

> **Application Note** AN-1208

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µC/OS-II **and** µC/Probe **on the Freescale MC9S08**

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1.00 Introduction

This document shows example code for using **µC/OS-II** on a Freescale MC9S08 processor. To demonstrate the MC9S08, we used a P&E DEMOQE evaluation board as shown in Figure 1-1.

We used the Freescale Codewarrior for Microcontrollers IDE version 6.0 to demonstrate this application. However, other tool-chains could be used.

Figure 1-1, P&E DEMOQE EVB

The application code is downloaded into Flash using a P&E Micro BDM Multilink. While the application is running, the 4 onboard LEDs sequentially illuminate and blink. μ C/Probe may be used to view various graphs depicting data read from the on-board accelerometer.

To use μ C/Probe with the DEMOQE128, download and install the trial version of the program from the Micriµm website as discussed below in the section 2.05, μ C/Probe. After programming your target with one of the included example projects, connect a RS-232 cable between your PC and the evaluation board, configure the RS-232 options (also covered below), and start running the program. The open data screens should begin to update. Please note that the communication baud rate for the MC9S08 should be configured to 38,400 baud and not higher.

1.01 Setting up the Hardware

Jumper Label Description J3 Set to USB Power
J4 Set to 3v Set to 3v J5 User may select J6 Across RXD and PTB0 J7 Across TXD and PTB1 J8 COM EN ON J9 All ON J11 PTC0, PTC1 both OFF J12 All ON $J13$ 0 $J14$ 0 $J15$ 1 J16 PTA7, PTC7, and PTA1 only J18 Reset_EN, LED_EN both ON J19 ON J20 SDL, SDA both ON J21 PTA0 only **Table 1-1**

The following table illustrates the recommended jumper configuration for the DEMOQE EVB.

1.02 Port Specific Details

Two **µC/OS-II** ports and example projects have been written for the MC9S08 to accommodate both the banked and non banked memory models. For both ports, the paging option must not be enabled within the Codewarrior project settings. This application note uses the terms non-banked and non-paged interchangeably. The terms banked and paged are also used in place of one another.

Care must be taken to ensure that the correct μ C/OS-II, and μ C/CPU files are compiled when running in either non-banked or banked memory model. Files from both ports must not be mixed together within any project. Separate directories have been created for files belonging to both the non-paged and paged μ C/OS-II and μ C/CPU ports. However, other modules such as **µC/Probe**, group the port files within the same directory and are distinguished by file name. Example **µC/Probe** file names for non-banked and banked assembly port files are as follows:

```
probe_rs232_nba.s
probe_rs232_ba.s
```
This example uses the on chip FLL. Once configured, the processor clock is set to 48MHz and the bus clock to 24MHz. The FLL settings may be changed from within BSP.c. by adjusting FLL_Init() accordingly. Note: In some cases, oscillator trimming may be desired. In order to adjust the port for a trimmed oscillator, the target oscillator frequency must be known and set within BSP.h. This may be accomplished by adjusting the macro BSP_INT_REF_FREQ_prior to compilation. The default value is set to 32768 Hertz.

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Additionally, this example assumes the use of TPM1, channel 0 for the μ C/OS-II periodic time tick interrupt. If TPM1 channel 0 is not available, this you may configure an alternate TPM1 channel by adjusting the macro named OS_TICK_TPM1_CH within BSP.h for channel numbers 0, 1 or 2. The file vectors.c will also require modification. See the section labeled "vectors.c" for more information.

All ISRs must be written as specified in the section labeled "Creating Interrupt Service Routines"

1.03 Directories and Files

The code and documentation of the port are placed in a directory structure according to "AN-2002, µC/OS-II Directory Structure". Files for both the banked and non banked memory models are provided by Micrium.

µC/OS-II:

\Micrium\Software\uCOS-II\Source

This directory contains the processor independent code for μ C/OS-II. The version used was 2.85.

\Micrium\Software\uCOS-II\Ports\MCS08\Paged\Codewarrior

This directory contains the standard processor specific files for a μ C/OS-II port assuming the Freescale Codewarrior IDE. These files could easily be modified to work with other tool chains; however, you would place the modified files in a different directory. Only one processor port, non-paged or paged may be compiled in to a μ C/OS-II project at any given time. Specifically, this directory contains the following files for the banked memory model:

 os_cpu.h os_cpu_a.s os_cpu_c.c

µC/Probe:

C:\Micrium\Software\uC-Probe\Target\Communication\Generic\OS\uCOS-II

This directory contains the OS dependent interface for the communication layer of **µC/Probe.** If you plan to run **µC/Probe** with a different RTOS, or without any RTOS, the following files would have to be adjusted accordingly:

probe_com_os.c

C:\Micrium\Software\uC-Probe\Target\Communication\Generic\RS-232\OS\uCOS-II

This directory contains OS dependent interface code for the RS-232 specific portion of **µC/Probe**, specifically the code necessary to generate an optional Rx packet parse task. If you plan to run μ C/Probe with a different RTOS, modifications to the files listed below will have to be made. If you are not running an RTOS, the following files may be excluded from the build.

probe_rs232_os.c

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C:\Micrium\Software\uC-Probe\Target\Communication\Generic\RS-232\Ports\Freescale\MC9S08

This directory contains the μ C/Probe hardware port files for the MC9S08 processor. Care must be taken to ensure that the correct assembly file, either probe $rs232$ ba.s (banked assembly) or probe rs232 nba.s (non banked assembly) is compiled into you project. Non-banked and banked assembly files must never be mixed within the same µC/Probe project.

probe_rs232c.c probe_rs232_ba.s probe_rs232_nba.s

C:\Micrium\Software\uC-Probe\Target\Communication\Generic\RS-232\Source

This directory contains target independent source code for the μ C/Probe RS-232 communication layer. Specifically, this directory contains the following files:

 probe_rs232.c probe_rs232.h

C:\Micrium\Software\uC-Probe\Target\Communication\Generic\Source

This directory contains target independent source code for the μ C/Probe communication layer. Specifically, this directory contains the following files:

 probe_com.c probe_com.h

C:\Micrium\Software\uC-Probe\Target\Plugins\uCOS-II

This directory contains the target independent source code for μ C/Probe. Specifically, this directory contains the following files:

 os_probe.c os_probe.h

µC/CPU:

\Micrium\Software\uC-CPU

This directory contains processor independent files for **µC/CPU**. **µC/CPU** contains code for entering and existing critical sections, as well as macro definitions for the 'C' programming datatypes used in most Micrium products. Specifically, this directory includes:

cpu_def.h

\Micrium\Software\uC-CPU\Ports\MCS08\Paged\Codewarrior

This directory contains processor port files for μ C/CPU. μ C/CPU contains code for entering and existing critical sections, as well as macro definitions for the 'C' programming datatypes used in most Micrium products. Care must be taken to ensure that the correct memory model port files for μ C/CPU are used when building your project. Only one port, either non-paged or paged, may be compiled into a given project at any time. The following directory contains the following files.

 cpu.h cpu_a.s

μC/LIB**:**

\Micrium\Software\uC-LIB

This directory contains **lib_def.h**, which provides #defines for useful constants (like DEF_TRUE and DEF_DISABLED) and macros.

\Micrium\Software\uC-LIB\Doc

This directory contains the documentation for **μC/LIB**.

Application Code:

\Micrium\Software\EvalBoards\Freescale\MC9S08QE128\DEMOQE\Paged\Codewarrior\OS-Probe This directory contains the Codewarrior project file, OS-Probe.mcp, for the banked version of AN-1208.

\Micrium\Software\EvalBoards\Freescale\MC9S08QE128\DEMOQE\Paged\Codewarrior \OS-Probe\Source

This directory contains the sample source code for using μ C/OS-II on the MC9S08 micro-controller. This directory contains the following files:

 app.c app_cfg.h probe_com_cfg.h derivative.h includes.h os_cfg.h start08.c vectors.c

app.c contains the application code, app_cfg.h contains application specific configuration information such as task priorities and stack sizes, probe com cfg.h contains μ C/Probe communication specific configuration, includes.h contains a master include file used by the application, os_cfg.h is the μ C/OS-II configuration file, derivative.h and start08.c are generated files for MCU compatibility and startup.

vectors.c contains the processor interrupt vector table. This array of Interrupt Service Routine addresses must be updated whenever a new interrupt is being configured on the system. Interrupt vectors that are not in use should be plugged with the appropriate Dummy ISR handler provided.

\Micrium\Software\EvalBoards\Freescale\MC9S08QE128\DEMOQE\Paged\Codewarrior\BSP

This directory contains the Board Support Package for the DEMOQE128 evaluation board and the MC9S08 MCU. While some of the code in this directory may work on other MC9S08 derivatives, routines that are hardware dependent such as LED_On() will require modification depending on the hardware design of your EVB. This directory contains:

 BSP.c BSP.h

BSP.c contains hardware specific source code for LED services, FLL initialization, µC/OS-II ticker initialization, and so on.

BSP.h contains prototypes for the BSP initialization and LED service functionality provided within BSP.c. The OS timer TPM1 channel selection and internal and external oscillator frequencies are also configured from within this file. You may adjust these settings by changing the macros OS_TICK_TPM1_CH, BSP_INT_REF_FREQ, and BSP_EXT_REF_FREQ respectively. BSP_EXT_REF_FREQ has been intentionally defined to 0 (Hz) since the DEMOQE128 EVB does not have an external oscillator on-board.

\Micrium\Software\EvalBoards\Freescale\MC9S08QE128\DEMOQE\Paged\Codewarrior

\OS-Probe\prm

 This directory contains the processor linker file. The user MUST remove all previously existing vector definitions from within this file in favor of those specified in vectors.c. Note: The linker file must be changed when porting to a different MC9S08 derivative.

1.04 Codewarrior IDE

We used the Freescale Codewarrior for Microcontrollers IDE version 6.0 to compile and run the MC9S08 example. You can of course use μ C/OS-II with other tools. Figures 1-3 shows the project source tree of the banked example project with all of the files visible in the project file source tree.

Figure 1-3, Codewarrior Source Tree

2.00 Example Code

As mentioned in the previous section, the test code for this board is found in the following directories and will be briefly described:

\Micrium\Software\EvalBoards\Freescale\MC9S08QE128\DEMOQE\Paged\Codewarrior\OS-Probe

It should be noted that processor header files and libraries are not included within the AN-1208 code archive since they are supplied by Freescale via the Codewarrior installation.

2.01 Example Code, app.c

app.c demonstrate some of the capabilities of **µC/OS-II**.

Listing 2-1, main()

```
void main (void) (1) 
{ 
   CPU_INT08U err; 
   BSP_IntDisAll(); (2) 
  OSInit(); (3)
   OSTaskCreateExt(AppStartTask, (4) 
             (void * )0, (OS_STK *)& AppStartTaskStk[APP_TASK_START_STK_SIZE - 1], 
              APP_TASK_START_PRIO, 
 APP_TASK_START_PRIO, 
 (OS_STK *)&AppStartTaskStk[0], 
              APP_TASK_START_STK_SIZE, 
             (void * )0, OS_TASK_OPT_STK_CHK | OS_TASK_OPT_STK_CLR); 
#if OS_TASK_NAME_SIZE > 11 
   OSTaskNameSet(APP_TASK_START_PRIO, "Startup Task", &err); (5) 
#endif 
   OSStart(); (6) 
}
```
- $L2-1(1)$ As with most C applications, the code starts in main().
- L2-1(2) We start off by calling a BSP function (see bsp.c) that will disable all interrupts. We do this to ensure that initialization doesn't get interrupted in case we do a 'warm restart'.
- L2-1(3) As with all **µC/OS-II** applications, you need to call $O(SInit()$ before creating any task or other kernel objects.

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- L2-1(4) We then create at least one task (in this case we used $OSTaskCreateExt()$ to specify additional information about your task to μ C/OS-II). It turns out that μ C/OS-II creates up to three tasks in σ sinit(). As a minimum, μ C/OS-II creates an idle task (OS_TaskIdle() which is internal to μ C/OS-II) and OS_TaskStat() (if you set OS_TASK_STAT_EN to 1 in os_cfg.h). OS_TaskStat() is also an internal task in **µC/OS-II**. The Timer task may be optional enabled or disabled from within os cfg.h.
- L2-1(5) As of V2.6x, you can now name μ C/OS-II tasks (and other kernel objects) and be able to display task names at run-time or, with a debugger. In this case, we name our first task 'Start Task'.
- L2-1(6) We finally start μ C/OS-II by calling OSStart(). μ C/OS-II will then start executing AppStartTask() since that's the highest priority task created. OSStart() does not return.

Listing 2-2, AppStartTask()

```
static void AppStartTask (void *p_arg) 
{ 
   (void)p_arg; 
  BSP\_Init(); (1)
#if OS TASK STAT EN > 0
  OSStatInit(); (2)
#endif 
#if (uC_PROBE_OS_PLUGIN > 0) (3)<br>
OSProbe_Init(); (4)
  OSProbe_Init();<br>OSProbe_SetCallback(AppProbeCallback); (5) (5)
   OSProbe_SetCallback(AppProbeCallback); (5) 
   OSProbe_SetDelay(50); (6) 
#endif 
#if (uC_PROBE_COM_MODULE > 0) 
  \text{ProbeCom\_Init}(); (7)<br>ProbeRS232 Init(38400); (8)
  ProbeRS232_Init(38400); (8)<br>ProbeRS232_RxIntEn(); (9)
  ProbeRS232_RxIntEn();
#endif 
  AppTaskCreate(); (10)
  LED Off(0); (11)
  while (DEF TRUE) \{ (12)
     for (j = 0; j < 4; j++) {
        for (i = 1; i \le 4; i++) {
           LED On(i); OSTimeDlyHMSM(0, 0, 0, 20); (13) 
           LED Off(i);
        OSTimeDlyHMSM(0, 0, 0, 20);
 } 
        for (i = 3; i >= 2; i--) {
            LED_On(i); 
            OSTimeDlyHMSM(0, 0, 0, 20); 
            LED_Off(i); 
        OSTimeDlyHMSM(0, 0, 0, 20);
 } 
      } 
     for (i = 0; i < 4; i++) {
        LED_0(1);LED_0(2);LED_0(3);LED_0(4); OSTimeDlyHMSM(0, 0, 0, 50);
```
LED $Off(1)$; LED_Off (2) ; $LED_Off(3);$ LED_Off(4); OSTimeDlyHMSM(0, 0, 0, 50); } } } $L2-2(1)$ BSP_Init() is called to initialize the Board Support Package – the I/Os, the tick interrupt, and so on. BSP_Init() will be discussed in the next section. L2-2(2) OSStatInit() computes how fast the CPU runs when OS TASK STAT EN is set to 1 in os_cfg.h. $L2-2(3)$ If uc_PROBE_OS_PLUGIN is defined in app_cfg.h, then include code to initialize the target independent portion of μ C/Probe. L2-2(4) Initialize the target independent portion of μ C/Probe. L2-2(5) Configure an optional call-back function for **µC/Probe**. The callback function is called each time μ C/Probe is queried to update variable values. L2-2(6) Set the delay in milliseconds that the **µC/Probe** task should delay before collecting information about task stacks and so forth. This task will not be necessary in future versions of uC/Probe. L2-2(7) Initialize the hardware independent portion of the μ C/Probe communication layer. L2-2(8) Initialize the hardware dependent communication channel. The baud rate for μ C/Probe RS-232 communication should not be set higher than 38,400 on the MC9S08. L2-2(9) Enable receive interrupts. This function call enables **uC/Probe** to begin processing requests from the host. L2-2(10) AppTaskCreate() is a user defined function for creating additional μ C/OS-II tasks. This function is not required and additional tasks could have been created directly within AppStartTask(). This function has been left empty in this example. L2-2(11) Shut off all onboard LEDs. L2-2(12) As with all task managed by μ C/OS-II, the task body must be in the form of an infinite loop. Tasks managed by $\mu\text{C/OS-II}$ must never be allowed to exit. Instead, tasks should be deleted using OSTaskDel() when they are no longer desired. This task performs the LED illumination and scrolling effect used within this application. L2-2(13) As μ C/OS-II tasks must either enter an infinite loop 'waiting' for some event to occur or terminate itself. In this case, we wait for time to expire as the 'event'. This is accomplished by calling OSTimeDlyHMSM().

2.02 Example Code, app_cfg.h

This file is used to configure:

- Additional **µC**/ module enables, specifically support for **µC/LIB** and **µC/Probe**
- **The** μ **C/OS-II** task priorities of each of the tasks in your application
- The stack size for each tasks

2.03 Example Code, includes.h

includes.h is a 'master' header file that contains #include directives to include other header files. This is done to make the code cleaner to read and easier to maintain.

2.04 Example Code, os_cfg.h

This file is used to configure μ C/OS-II and defines the maximum number of tasks that your application can have, which services will be enabled (semaphores, mailboxes, queues, etc.), the size of the idle and statistic task and more. In all, there are about 60 or so $\# \text{define }$'s that you can set in this file. Each entry is commented and additional information about the purpose of each #define can be found in the **µC/OS-II** book. os_cfg.h assumes you have **µC/OS-II** V2.83 or higher but also works with previous versions of **µC/OS-II**.

2.05 μC/Probe

µC/Probe is a Microsoft Windows program that displays the content of system variables on various user definable graphical elements such as simulated mechanical counters, graphs, on-screen LEDs and so on.

In order for μ C/Probe to display information about your application, an ELF file, must be generated by the user's compiler. The ELF file contains the names and addresses of all the global symbols referenced within the users embedded application. Only symbols that have been allocated memory, e.g. not allocated on the stack, are able to be monitored by μ C/Probe. Global and static variables are examples of variables that may be monitored.

The user places components (such as gauges, labels, and charts) into a Data Screen in a **µC/Probe** workspace. Each one of these controls is then assigned to one or more of the variables from the Symbol Browser. The Symbol Browser lists all symbols referenced from within the ELF file. Symbols associated with components placed on an open Data Screen will be updated after the user presses the start button (assuming the user's PC is connected to the target and the target is running).

µC/Probe currently interfaces with a target processor via JTAG, RS-232, UDP and USB. A small section of code resident on the target receives commands from the Windows application and responds to those commands. The commands ask for a certain number of bytes located at a certain address, for example, "Send 16 bytes beginning at 0x0040102C". The Windows application, upon receiving the response, updates the appropriate component(s) on the data screen(s) with the new values.

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Figure 2-1. µC/Probe **Windows Program**

To use **µC/Probe** with the example project (or your application), do the following:

1. **Download and Install** µC/Probe**.** A trial version of µC/Probe can be downloaded from the Micrium website at

<http://www.micrium.com/products/probe/probe.html>

2. Open µC/Probe. After downloading and installing this program, open the example µC/Probe workspace for µC/OS-II, named *OS-Probe.wsp*, which should be located in the AN-1208 Codewarrior project directory.

You may also open one of the sample workspaces that comes with μ C/Probe. The sample workspaces, located in the μ C/Probe target directory, contains generic workspaces for **µC/OS-II**, as well as other Micrium software modules.

3. **Connect Target to PC**. Currently, **µC/Probe** can use RS-232 to retrieve information from the target. You should connect a RS-232 cable between your target and computer.

Load Your ELF File. The example projects included with this application note are already configured to output an ELF file. (If you are using your own project, please refer to Appendix A of the μ C/Probe user manual for directions for generating an ELF file with your compiler.) Codewarrior generates an ELF file with a .abs extension. This file is located in a directory named BIN within the sample project directory.

To load this ELF file, right-click on the symbol browser and choose "Add Symbols". Navigate to the file directory, select the file, and choose "OK".

- 4. **Configure the RS-232 Options.** In **µC/Probe**, choose the "Options" menu item on the "Tools" menu. A dialog box as shown in Figure 6-2 (left) should appear. Choose the "RS-232" radio button. Next, select the "RS-232" item in the options tree, and choose the appropriate COM port and baud rate. The baud rate for the projects accompanying this application note is 38,400.
- 5.
- **6. Start Running**. You should now be ready to run **µC/Probe**. Just press the run button \blacktriangleright to see the variables in the open data screens update.

3.00 Board Support Package (BSP)

BSP stands for Board Support Package and provides functions to encapsulate common I/O access functions in order to make it easier for you to port your application code. In fact, you should be able to create other applications using the DEMOQE128 evaluation board and reuse these functions thus saving you a lot of time.

The BSP performs the following functions:

- Determine the MC9S08s CPU clock and bus frequencies
- Configure the LED I/Os for the DEMOQE128 EVB and MC9S08 CPU
- Configuration and handling of the **µC/OS-II** tick timer

The BSP for the DEMOQE128 is found in the follow directory.

\Micrium\Software\EvalBoards\Freescale\MC9S08QE128\DEMOQE\Paged\Codewarrior\BSP The BSP files are:

 BSP.c BSP.h

3.02 Board Support Package, bsp*.*

We will not be discussing every aspect of the BSP but only cover topics that require special attention.

Your application code must call BSP_Init() to initialize the BSP. BSP_Init() in turn calls other essential functions for initializing the MCU to the desired run-time state.

Listing 3-1, BSP_Init()

```
void BSP_Init (void) 
{ 
   SOPT1 \&= \sim SOPT1_COPE_MASK; (1)<br>FLL Init(); (2)
   FLL_Init(); (2)<br>
Tmr_TrickInit(); (3)
   \text{Trm}_{\text{inter}}\text{TrickInit}(); (3)<br>
LED_Init(); (4)
   LED_Init();
}
```
- L3-1(1) Disable the watchdog timer. It is enabled automatically after reset.
- L3-1(2) Initialize the FLL, set ICS_OUT to 50.33MHz using the internal oscillator.
- L3-1(3) Initialize the selected TPM1 channel for use with the μ C/OS-II time tick interrupt. The code for this function is described below.
- L3-1(4) Initialize the general purpose I/O pins used for controlling the onboard LEDs.

Listing 3-2, Tmr_TickInit()

```
static void Tmr_TickInit (void) 
{ 
     CPU_INT32U bus_freq; 
     bus_freq = BSP_CPU_ClkFreq(); (1) 
    bus_freq / = 2iOSTickCnts = (CPU INT16U)(bus free / (8 * OS TICKS PER SEC)); (2)\begin{array}{lll}\n\text{SCGCl} & & | = & \text{SCGCl\_TPM1\_MASK} \n\end{array} \tag{3} \tag{3} \begin{array}{lll}\n\text{TPM1SC} & & & \\
\text{SCH} & &TPMISC = 0; (4)
TPM1MOD = 0; (5)
#if OS_TICK_TPM1_CH == 0 (6)<br>TPM1COSC = TPM1COSC MS0A MASK | (7)
 TPM1C0SC = TPM1C0SC_MS0A_MASK | (7) 
 TPM1C0SC_CH0IE_MASK; 
    TPM1C0V = TPM1CNT + OSTickCnts; (8)
#endif 
#if OS_TICK_TPM1_CH == 1 
 TPM1C1SC = TPM1C1SC_MS1A_MASK | 
 TPM1C1SC_CH1IE_MASK; 
     TPM1C1V = TPM1CNT + OSTickCnts; 
#endif 
#if OS_TICK_TPM1_CH == 2 
    TPM1C2SC = TPM1C2SC_MS2A_MASK | 
                     TPM1C2SC_CH2IE_MASK; 
   TPM1C2V = TPM1CNT + OSTickCnts;
#endif 
    TPM1SC |= (1 << TPM1SC_CLKSx_BITNUM) | (9) 
                    (3 \leq \text{TPMISC} PS BITNUM);
}
```
- L3-2(1) Get the CPU operating frequency (ICS OUT) in Hz. This number is divided by 2 in order to generate the bus frequency. The bus frequency is the reference frequency for the TPM.
- L3-2(2) Compute the number of timer increments necessary to generate the desired µC/OS-II time tick period. The '8' represents the TPM prescaler that is configured in step 10.
- L3-2(3) Enable the TPM1 module clock.
- L3-2(4) Stop and reset the TPM1 counter.
- L3-2(5) Set the modulus count to 0x00 allowing the timer to freerun from 0x00 to 0xFFFF.
- L3-2(6) If $OS_TICK_TPML_CH$ in BSP.h is defined as 0, 1, or 2, then configure the corresponding TPM1 channel for use with the μ C/OS-II periodic time tick interrupt.
- L3-2(7) Configure the timer channel for output compare and enable TPM1 channel 0 interrupts.
- L3-2(8) Write the match register with the current value of the TPM1 counter plus the number of ticks until the next desired match.
- L3-2(9) Configure the TPM1 reference clock source, set the prescaler and enable the timer.

Listing 3-3, Tmr_TickISR_Handler()

```
void Tmr_TickISR_Handler (void) 
{ 
\begin{array}{lll} \# \text{if OS\_TICK\_TPM1\_CH & =& 0 \ \end{array} \tag{1} \ \text{TPM1C0SC} \quad \& = & \text{TFM1C0SC CH0F MASK} \tag{2} \end{array}TPM1C0SC \t <= \t TPM1C0SC_CH0F_MASK; (2)
 TPM1C0V += OSTickCnts; (3) 
#endif 
#if OS_TICK_TPM1_CH == 1 
 TPM1C1SC &= ~TPM1C1SC_CH1F_MASK; 
 TPM1C1V += OSTickCnts; 
#endif 
#if OS_TICK_TPM1_CH == 2 
 TPM1C2SC &= ~TPM1C2SC_CH2F_MASK; 
 TPM1C2V += OSTickCnts; 
#endif 
    OSTimeTick(); (4) 
}
```
This function is called from an assembly interrupt service routine which informs μ C/OS-II of the interrupt and calls the 'C' code interrupt handler. See os cpu a.s and the section labeled "Creating Interrupt Service Routines" for more information.

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The TPM generates a match interrupt when the up-counter value reaches the value stored within the timer channel match register. After an interrupt occurs, the match register is incremented to the next value for which a time tick interrupt is desired. The timer is allowed to free-run and overflow without error when necessary.

Figure 3-1, OS Tick Timer Operation

When the selected TPM issues an interrupt, the processor vectors to the associated TPM module and channel interrupt service routine. The ISR calls Tmr_TickISR_Handler() as described above in Listing 3-3.

You should note that ALL of your ISRs should be written in assembly and call an ISR handler function written in C. The ISR handler should be prototyped as 'void MyISR Handler (void)'. The assembly interrupt service routine is necessary in order to maintain a fully preemptive RTOS. If the processor is allowed to jump directly to an ISR written in C, then any task made ready to run as the result of that ISR may have to wait until the interrupted task is resumed and the scheduler runs. This is non deterministic and is determined by the application code length. Instead, the ISR handler returns to the assembly interrupt service routine which immediately calls the scheduler preventing the restoration of a lower priority task context when a higher priority task should run upon completion of the ISR. You should note that ALL of your ISRs should be written in assembly and
function written in C. The ISR handler should be prototyped as 'vooical'. The assembly interrupt service routine is necessary in order
preemptive RTOS

3.03 Configuring the FLL

The FLL is an on chip peripheral capable of boosting the processor clock and bus frequencies higher than the frequency provided by either an external or internal oscillator. Before attempting to reconfigure the FLL from BSP.c you should consult the MC9S08 datasheet and understand the MCU's absolute maximum ratings. The absolute maximum ratings must be followed in order to prevent the possibility of damaging the device.

The MC9S08 has a maximum processor clock of 59.77MHz. The bus clock is always ½ of the processor clock (ICS_OUT) and must never exceed 25.16 MHz. Therefore, when running ICS_OUT at the maximum frequency, it becomes necessary to further divide the bus clock to prevent overclocking of the peripheral bus. Further bus clock division may be performed within FLL Init() by setting the BDIV bits greater than 0.

There is no external oscillator on the DEMOQE128, therefore, all clock references are derived from the MC9S08 internal oscillator and FLL.

The FLL output frequency is computed as follows:

```
(32768Hz * Pre-selected Multiplier)
```
Where the pre-selected multiplier has been hard coded within $BSP.c$, FLL Init()as 1536 by the setting of the DCO range to high-range and the DMX32 bit to 0.

Note: BSP_CPU_ClkFreq() returns a 32 bit unsigned integer representation of ICS_OUT, the processor clock frequency. Dividing this value by 2 will yield the bus clock frequency during runtime. It is recommended that users call this function in order to compute peripheral clock dividers as opposed to hard coding them. This ensures system integrity should the clock settings require modification at a later time. The use of this function is demonstrated in $\text{Tr}_{\text{Tr}}\text{TickInit}()$.

3.04 vectors.c

vectors.c contains the interrupt vector table for the application. The interrupt vector table is necessary so that the processor knows the address of the interrupt service routine to jump to when a specific interrupt occurs. Failure to properly plug the interrupt vector table with the address of a valid handler may cause the application to crash. If a wrong, but valid, interrupt handler address is specified for vector number 'n' and the interrupt occurs, the interrupt source will not be cleared the processor will execute the same interrupt service routine indefinitely.

Care should be taken when working with the interrupt vector table.

For convenience, dummy interrupt service routines have been provided for all 32 vectors excluding reset. When an interrupt vector is not in use, the dummy ISR for that vector should be set. In the case of a spurious interrupt, the processor will vector to the dummy ISR and loop indefinitely. Should this occur, you may be able to debug the application and catch the processor in the dummy interrupt service routine thus identifying the source of the spurious interrupt. The correct action may then be taken to correct the application to prevent this type of error in the future.

When plugging the interrupt vector table with a new ISR address, a 'C' prototype in the form of:

extern void near MyISR(void);

Must be provided at the top of the file. The name of the ISR may then be plugged into the correct location of the interrupt vector table.

The following vectors are used by μ C/OS-II and should not be modified:

- Vector 1: SWI. Used to perform the **µC/OS-II** context switch.
- Vector 4: TPM1 channel 0. This may be adjusted to one of the other TPM1 channel vectors if desired. The macro named OS_TICK_TPM1_CH in BSP.h will also have to be adjusted accordingly.

3.05 Creating Interrupt Service Routines

All interrupt service routines must contain a short assembly routine. The address of the assembly routine is used to plug the interrupt vector table, while the content is designed to notify μ C/OS-II of the interrupt and call the user supplied interrupt handler written in either assembly or 'C' code.

The prototype specified at the top of Vectors.c (See section 3.04 above) is the 'C' code prototype for the following assembly interrupt service routine. It is this prototype that allows you to plug the interrupt vector table with the name (address) of the ISR from 'C'.

As a reminder, the prototype is written as follows:

extern void near MyISR(void);

Of course, the name of the ISR would change each time a new ISR is declared since two ISR's of the same name cannot exist in the system simultaneously.

The format of an interrupt service routine for the BANKED memory model is as follows:

Listing 3-4, MyISR

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- L3-4(1) Force the contents of the assembly file, perhaps named: myisr_a.s, into NON_BANKED memory. This is critical since the processor only has a 16 bit address bus. Vectors that are accidentally placed into banked memory will have a 24 bit address (8 bit page number + 16 bit address) and will overflow the slot in the interrupt vector table. The interrupt handler written in 'C' may be placed in banked memory.
- L3-4(2) A label for the PPAGE register memory address is defined. This label is used in place of hard coding the PPAGE register address multiple times within the source code.
- L3-4(3) XDEF is a Codewarrior assembly directive for prototyping assembly routines thereby making them visible from C files such as vectors.c. This directive is not necessarily portable to other assemblers.
- L3-4(4) (4), (5), (6), and (7) are external references to variables and or functions defined in 'C'. These variables are referenced from the context of the assembly ISR and must therefore be declared external such that they are visible to the assembler and ISR file. This directive is not necessarily portable to other assemblers.
- L3-4(8) MyISR is the name of the interrupt service routine and must match the name specified in L3-4(3). This ISR name is also external from within vectors.c.
- L3-4(8) Store the H register on the stack of the task that was interrupted. This register is not automatically stored and restored by the processor interrupt mechanism and must be pushed and pulled from the stack separately.
- L3-4(9) Load the value of the PPAGE register into the A register.
- L3-4(10) Store the current value of PPAGE on to the stack of the task that was interrupted. This register is not automatically stored and restored by the processor interrupt mechanism and must be pushed and pulled from the stack separately. Note: enabling the PPAGE checkbox within the compiler options in Codewarrior, will cause the compiler to add code for storing and restoring PPAGE on and off the stack when an interrupt occurs. Since not all compilers support this feature, the storing and restoring of PPAGE is manually specified within the port. Enabling this feature will cause PPAGE to be pushed twice on to the stack and cause the system to fail.

- L3-4(11) Store the H register on to the stack of the task that was interrupted. This register is not automatically stored and restored by the processor interrupt mechanism and must be pushed and pulled from the stack separately.
- L3-4(12) (12), (13), and (14), increment $OSLntNesting$. This notifies $\mu\text{C/OS-II}$ that at least one interrupt is in progress and that the scheduler should not schedule any new tasks to run until all nested interrupts have completed (e.g. OSIntNesting equals 0). It also informs the OS of the interrupt thereby allowing certain μ C/OS-II API calls to return error codes if called from within the context of an ISR.
- L3-4(15) Load a copy of OSIntNesting from memory into a register so a comparison may be made.
- L3-4(16) Check OsIntNesting to see if its value is 1. If so, then this is the only interrupt in progress and no nested interrupts are pending completion.
- L3-4(17) $(17 24)$, If no interrupts have been nested then the scheduler is free to schedule a new task when the ISR completes. Therefore the address of the current task TCB (Task Control Block) is obtained and the current stack pointer is saved to the current tasks TCB. (OSTCBCur->OSTCBStkPtr = Stack Pointer)

 If interrupt have been nested, skip storing the current tasks stack pointer back into its task control block and jump to $MyISR1$. Note: the name of the ISR and the labels used within it must be changed for each new ISR implemented in the system. For convenience, the number '1' is added to the end of the ISR name in order to create a unique and convenient label to jump to.

- L3-4(25) Call the user defined ISR handler. Generally the ISR handler is defined and prototyped in 'C'. The ISR handler may be located in banked memory.
- L3-4(26) Call OSIntExit) . This informs **µC/OS-II** about the end of the interrupt. This is effectively the same as decrementing OSIntNesting and calling the scheduler. A context switch to a task different than the one that was running before the interrupt may or may not occur.
- L3-4(27) Restore the H register.
- L3-4(28) Pull the value of PPAGE of the stack of the task that is to be resumed.
- L3-4(29) Restore the PPAGE register.
- L3-4(30) Return to the highest priority task that is ready to run.

The difference between the NON-banked and the banked memory model from an assembly ISR perspective is the use of JSR vs. the CALL instruction. CALL uses a 24 bit address to jump to and return from the called function, JSR only uses 16 bit addresses and is only suitable for calling functions in non-banked memory. Since one can never be sure of a functions linked address ahead of time, it becomes necessary to substitute JSR for CALL in the banked version of the **µC/OS-II** port and associated user interrupt service routines. Likewise, RTS, return from subroutine must be substituted with RTC, return from call. Lastly, the NON-banked port does not need to save or restore the PPAGE register before or after context switches. This is however necessary in order to resume execution from within a function that was interrupted from a banked page in memory.

For more information, one may compare the non-banked and banked versions of os_cpu_a.s. Specifically, one should examine $Tmx_TickISR$ which is the $\mu C/OS-II$ time tick ISR. All interrupts should be modeled exactly as shown with the exception of the labels used to identify the sub-routine. Additional examples are available upon request.

3.05 Context Switching on the MC9S08

When a new task is created, a stack frame for that task is initialized from within os cpu c.c, OSTaskStkInit(). The values placed on the initial stack frame represent the desired values of the processor registers when the task runs for the first time.

The following code shows stack initialization for a banked memory project. Non-banked memory projects do not need to push a value for PPAGE on to the tasks stack.

Listing 3-5, OSTaskStkInit()

```
OS_STK *OSTaskStkInit (void (*task)(void *pd), void *p_arg, OS_STK *ptos, INT16U opt) 
{ 
    INT16U *wstk; 
    INT8U *bstk; 
   wstk = (INT16U * )ptos; (1)
   *--wstk = (INT16U)p_{arg};<br>*--wstk = (INT16U)(task); (3)
   *--wstk = (INT16U)(task);<br>
bstk = (INT8U *)wstk; (4)
   bstk = (INT8U * )wstk; (4)<br>*-bstk = (TNT8U)0x22; (5)
   *--bstk = (INT8U) 0x22;*--bstk = (INT8U)0xAA; (6)<br>*--bstk = (0x00); (7)
   *--bstk = (0x00);<br>*--bstk = (INT8U)((INT32U)(task) >> 16);<br>(8)
   *--bstk = (INT8U) ((INT32U) (task) >> 16);
   *--bstk = (INT8U)0x11; (9)<br>
return ((OS STK *b)stk); (10)
   return ((OS_STK\hspace{0.1mm}^*)bstk);
}
```
- L3-5(1) Cast a pointer to the top of the stack. The top of the new tasks stack is known from the call to either OSTaskCreate() or OSTaskCreateExt().
- L3-5(2) Push the argument for the task function on to the stack. The argument is an optional argument passed to either OSTaskCreate() or OSTaskCreateExt() and is often passed as NULL.
- L3-5(3) Push the tasks (function) address on to the stack. The address is a 16 bit value and does not include the page number that contains the task.
- L3-5(4) Cast the current top of stack pointer to an 8 bit pointer since the remaining values to be pushed are only 8 bits each.
- L3-5(5) Push the initial value for the X register on to the stack. We choose $0x22$ since it's the second register to be restored. This value is easy to identify during debugging.
- L3-5(6) Push the initial value for the A register on to the stack. We choose 0xAA since the value corresponds to the name of the register. This value is easy to identify during debugging.
- L3-5(7) Push the initial value for the CCR register on to the stack. Interrupts for the tasks context are enabled by means of clearing the I-bit for the initial CCR value. Since the

task has yet to run, we initialize the CCR to 0 to indicate that no condition codes have been set, and that global interrupts should be enabled when the task is running.

- L3-5(8) Push the PPAGE value of the tasks address on to the stack.
- L3-5(9) Push the initial value for the H register on to the stack. We chose $0x11$ since H is the first register to be restored. This value is easy to identify during debugging.
- L3-5(10) Return a pointer to the new top of stack.

Note: since a task body is in the form of an infinite loop, the return address is never utilized and remains on the stack along with p_{arg} indefinitely.

When the scheduler switches task context, the RTI (Return From Interrupt) instruction is used to restore all of the system registers from the stack pointer provided.

Since the RTI instruction does not restore H or PPAGE, they must be manually restored before using RTI. Again, the non-banked port does not require the manual saving or restoration of PPAGE.

After task stack initialization, the new tasks stack looks as follows:

Stack grows toward low memory

Figure 3-2

When μ C/OS-II scheduler determines that a context switch is necessary, it calls a macro named OS_TASK_SWITCH()located within os_cpu.h. This macro is defined as __asm swi which in turn causes the scheduler to invoke a software interrupt. vectors.c loads the SWI interrupt vector with the address of $OSCtxSw$ () which is located within $os_cpu_a.s.$

Upon entering OSCtxSw(), SWI pushes the majority of the interrupted tasks context on to the interrupted tasks stack. Notice the absence of the PPAGE and H register as shown in figure 3-3 below. These registers must be manually saved and restored by the **µC/OS-II** port. For nonbanked projects, the PPAGE is exempt from this requirement.

CCR A X PC High Byte PC Low Byte … … Stack Pointer Register -> **CCR Top of Stack Low Memory**

Stack grows toward low memory

Figure 3-3

High Memory

Listing 3-6, OSCtxSw()

L3-6(1) Obtain the PPAGE register value associated with the program counter of the preempted task. Push PPAGE on to the preempted tasks stack.

- L3-6(2) Push the H register on to the preempted tasks stack. The stack frame for the preempted task is now complete.
- L3-6(3) The next group of instructions loads the preempted tasks stack pointer into the 16 bit HX register. The stack pointer is then pushed on to the preempted tasks stack for temporary storage. A pointer to the preempted tasks TCB is then loaded into HX . This pointer also points to the storage location for the tasks stack pointer. The tasks stack pointer is stored in this location each time the task is swapped out in favor of a higher priority task that is ready to run.
- L3-6(4) The next group of instructions pulls the preempted tasks stack pointer off of the preempted tasks stack 1 byte at a time. Each byte is stored within the holding place located within the preempted tasks TCB. Later, when the this preempted task becomes the highest priority task that is ready to run, the stack pointer will be recovered from the TCB and used to restore this tasks context.
- L3-6(2) Call the μ C/OS-II task switch hook. This function is defined within os_cpu_c.c and may be used to notify the application of a context switch. In newer example applications, the task switch hook calls an application defined task switch hook to avoid modification to files used by the μ C/OS-II port. Application hooks may be enabled by defining OS_APP_HOOKS_EN to 1 within os_cfg.h. Upon doing so, several functions will need to be defined within the user application to handle the various hook function made available by μ C/OS-II to the user application.
- L3-6(6) Update the current running task priority to reflect the priority of the task that is about to be swapped in.
- L3-6(7) Update the currently running task TCB pointer to reflect the address of the TCB that is about to be swapped in; store this value in the HX register. As mentioned in L3- 6(4), the new task to be swapped in stores the value of its last known stack pointer in the first location of the TCB. Therefore, the address of the current running task (though it is not yet running) also points to the stack pointer associated with the task that will run next. The pointer is dereferenced and the stack pointer address is also stored in HX. Lastly, the stack pointer address is copied into the stack pointer register.
- L3-6(8) Restore the new tasks H register. Recall, that the H register is not automatically restored. Before the task may be run, the stack frame must be reduced from that of figure of 3-2 to that of 3-3.
- L3-6(9) Obtain and restore the new tasks PPAGE register. Recall, that the PPAGE register is not automatically restored by the RTI (Return from Interrupt) instruction. In fact, for the non-banked **µC/OS-II** port, PPAGE must not be restored. Once PPAGE has been removed from the stack, the new tasks stack looks like that of figure 3-3 and is ready to be restored automatically by the RTI instruction.
- L3-6(10) Return from interrupt. This instruction automatically pops PC , X , A , and CCR from the current stack pointed to by the stack pointer register. The new task is now running from where it left off prior to its last preemption.

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References

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