This paper says that the logical thing would be that the samples should be measured again if the environment changes.

They have studied six different environments in two cases:

- ‐ The kNN from each sample were computed only among the other samples from the same environment (self – self).
- ‐ The kNN from each sample were computed using samples from all environments $(self - all).$

They conclude that the self‐self case is better than the self‐all case in terms of accuracy, but the standard deviation of the self‐self case is noticeably larger than that of the self‐all case.

They also study the K which provides the best results and they talk about the hardware they used, the protocol, the problems they found such us some ACK's were lost, etc

K-Nearest-Neighbor Analysis of Received Signal Strength Distance Estimation Across Environments

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*Abstract***— Prior studies investigating the use of non-parametric models for ranging and localization via received signal strength have been restricted usually to one or two relatively similar environments, and have used primarily to 802.11 network interfaces. This paper discusses methods for and results of ranging experiments with signal strength measurements in a sensor network as the environment is varied. The K-nearest-neighbor distance estimation error is computed for both ranging and localization scenarios. Degradation of the mean 80th percentile error from 30 feet to 50 feet is seen when multiple environments are considered; however, the standard deviation improves from 16 feet to 10 feet. The localization results are similar.**

I. INTRODUCTION

Localization of sensors is a common challenge in the operation of indoor and short-range outdoor wireless sensor networks. A cost-effective solution is to use the received signal strength (RSS) across the wireless channel. Since mathematical propagation models of indoor RSS are still considered both quite complex and not consistently accurate [1], most RSS algorithms to date have used non-parametric approaches that pre-sample the RSS in the operating environment. Many algorithms have been proposed for this task, but the results of [2] show that they all perform comparably due, most probably, to inherent uncertainty in the environment.

Past work has shown a median $(50th$ percentile) estimation error of 9.64 feet [3], 10.4 feet [4], and 10 feet [2] when the training samples were taken from the same environment as the data samples. [3] found a median error of 14.1 feet for a parametric approach, and [5] found a RMS error of 6 feet for a different parametric approach in a dense network of sensors. [2] shows that a median percentile error of 10 feet and a $97th$ percentile error of 30 feet can be expected when using any non-parametric algorithm on a pre-sampled environment.

Sampling an environment to train a non-parametric algorithm can be a time-consuming task. Furthermore, if the environment changes after sampling, conventional logic would say to re-sample the environment. We investigate the error that might arise from a drastic change in environment without a change in training samples.

We show in this paper that a good baseline for RSS based algorithms in a MICA2-based sensor network is an 80th percentile error of approximately 30 to 40 feet. A degradation in accuracy of 20 feet is seen when a generic sample set from multiple environments is used, but the variance of the error decreases. We also discuss practical challenges of design and implementation that we encountered during our experiment.

II. OVERVIEW

MPR400 MICA2 sensor motes from Crossbow [6] are used for all experiments in this paper. They have a Chipcon CC1000 radio, which uses FSK modulation and provides a Received Signal Strength Indicator (RSSI) output that is sampled by a 10-bit ADC. [7] All of our data was taken at 915.998 MHz. TinyOS [8] was running on the motes and Sensor-MAC (S-MAC) was used for communication. We developed a custom algorithm for collecting and recording the data.

We recorded the output value of the RSSI for communications at 255 different power levels as the Physical Separation Distance (PSD) was increased from 2 feet to 150 feet in increments of 2 feet. Communication was not possible beyond 150 feet. We analyzed the resulting 255 dimensional dataset $\sin \varrho$ a leave-one-out K-nearest-neighbor (kNN) algorithm to determine the cumulative distribution function (CDF) of the distance estimation error. We also used the distance estimates from the kNN analysis to simulate a localization problem where a number of "beacon" motes, whose absolute locations are known, are used to estimate the location of an unfixed mote using only RSSI-based distance estimates to the beacons.

The leave-one-out kNN analysis followed [9]. It involves taking one sample as a "data sample" and computing a distance metric in the signal space to other samples, denoted as "training samples." The signal space distances (SSD) are sorted, and the K samples with the smallest SSD are chosen as the "K Nearest Neighbors." The PSD's for the kNN are then combined to estimate the PSD of the data sample. This is compared with the actual PSD to calculate the distance estimation error. The cumulative distribution function (CDF) of this error was plotted, along with the mean and standard deviation of the CDF's across environments. This analysis was run in two scenarios: (1) the data sample was compared only to training samples from its own environment (self-self); (2) the data sample was compared to training samples from all environments (self-all).

The localization algorithm took a random set of samples from within one environment and placed them at random in a 2-D plane such that they were located at the correct PSD from an "unfixed" mote at the origin. Using the kNN estimated PSD's for the random set of samples, we minimized an error function to estimate the unfixed mote's location.

Samples were taken in 6 environments chosen for their variation in terms of contruction materials, architectural design, and amount and type of "clutter." We did not do any extensive modeling to determine precise multipath and fading properties because one purpose of this experiment is to determine how well we can estimate PSD without complex environmental studies. A description of each environment is given below, along with some pictures in Figures 1 and 2.

A. Environment A

A 150 x 80 x 12 foot agricultural building housing 400 beef cattle: The structure has an open ceiling and open 8-foot studs sitting on top of 4-foot concrete sidewalls. The sensors were located on a concrete driveway running down the middle of the building that is flanked by cattle pens.

B. Environment B

A 220 x 44 x 10 foot dairy building housing 120 dairy cattle: The structure has a 10-foot, covered ceiling and open, 8-foot stud walls sitting on top of 2-foot concrete walls. There is a large feed bunk in the middle of the barn that contains conveyors and electric motors. The walls are lined with 8 foot deep freestalls made of metal tubing and rubber floor matresses. The sensors were located in the freestalls during the experiment. Cattle were kept on the opposite side of the feed bunk to safeguard the sensors.

C. Environment C

A 150 x 100 x 12 foot dairy building housing approximately 120 dairy cattle: The structure has steel I-beams running straight up from the ground to a 12-foot eave, at which point they angle inward and run the rest of the way up the roof. The sides of the barn are open and the ends and roof are covered with steel siding. A picture of the data capture setup for this is shown in Figure 1. Some variations in this environment were created by opening and closing garage doors, driving tractors through the cattle pens, and by the presence of varying numbers of cattle.

D. Environment D

A 44 x 20 x 8 foot milking parlor where dairy cattle are milked: It is strikingly different from the other environments studied because the sensors did not have line-ofsight communication. The walls are made of ceramic-coated concrete masonry blocks. The room is filled with large, angled, stainless-steel obstructions designed to hold cattle. A picture of this complex environment is in Figure 2.

E. Environment E

A 250 x 8 x 8 foot hallway on the third floor of the Materials Science and Electrical Engineering (MSEE) building at Purdue University: The most important feature of this environment is the constant flow of people through the hallway.

Fig. 1. Set-up for taking samples in Environment C

Fig. 2. Taking Samples in Environment D

F. Environment F

An open, grassy field in West Lafayette, Indiana.

III. EXPERIMENTAL SETUP

This experiment was performed twice, with the second iteration containing all the improvements learned from experience with the first. The challenges we encountered and our solutions are presented in this section.

A. Hardware

The MICA2 motes were connected to an Apple Powerbook G4 laptop using a MIB510 serial programming board from Crossbow. The motes were set on cloth-topped camping tables that were 1.25 feet tall. A plastic, 300-foot measuring tape was used for measuring PSD and was kept stationary throughout each experiment. A picture of the setup appears in Figure 1.

B. Network Protocols

The most important requirement placed on our network protocol design comes from the kNN analysis. Any "missing" data are assumed to be missing because the transmission power was too low for communication. Therefore, the network must not allow any data to be lost during transit.

We chose S-MAC over the standard TinyOS B-MAC protocol because B-MAC has a maximum packet payload size of 29 bytes, while S-MAC will allow up to 250 bytes. Our first algorithm using B-MAC took 5 hours to take one sample at each PSD; our final algorithm using S-MAC took 5 hours to take 10 samples at each PSD, a factor of 10 improvement.

The largest source of error when using S-MAC was **ACK's** that were lost relatively frequently. One major drawback of using larger packet sizes is that there can be no TCP-like transport layer to guarantee delivery of all packets because there is only 4 KB of RAM available to a program on the motes (to prevent corruption of the call stack during execution). This is not enough to store a number of large packets for the purpose of retransmission. If smaller packets are used, then retransmission buffers can be maintained in a lower level of the communication stack.

It is important to note that this problem cannot be eliminated simply by centralized polling by the computer. The signal patterns of the MICA2 motes have distinct nulls that vary with environment and can cause communication to be unreliable even where one would expect it to be fine, thus allowing ACK's to be lost for reasons other than collisions.

Our final network algorithm involves a ring-like configuration of five sensor motes: the *base, sender, receiver, sender relay,* and *receiver relay*.

The network configuration described below is used only for collecting the data that is generated by the *sender* and *receiver*. It was designed to ensure that no data is lost between these two motes and the computer, and also to ensure that there is little to no signal contention between motes during the transmissions from the sender to the receiver. Our results will hold for sensor grids of any density, provided the sensors have been synchronized and avoid transmitting simultaneously. If there are simultaneous transmissions, shadowing will occur, and the results will likely be degraded.

C. Base

The base mote is attached to the programming board that connects to the computer and is used to relay information between the computer and the network. The communication to the computer was handled through a custom serial interface that we wrote for transferring arbitrarily large amounts of information between the base station and a C program running on the computer.

D. Sender

The sender mote waits to receive a control message from the sender relay. The sender then transmits a "start message" to the receiver and waits for a "ready message." Then it broadcasts sample packets at each of the 255 available transmission power levels. A sample packet contains the transmitter battery voltage, the transmission power level, the PSD, and the sequence number of this sample. Once the 255 transmissions have taken place, a "done message" is sent to the receiver. If there are more samples left to take at this PSD, then it repeats the process again by sending another start message.

E. Receiver

The receiver mote waits for a start message from the sender. It responds with a ready message if its sample buffer is empty. Upon reception of a sample packet, it records the RSSI and its own battery voltage level, as well as the information received within the sample packet itself. It does this until a done packet is received from the sender. It organizes the stored data into "record messages" and sends them to the receiver relay for delivery to the computer.

F. Sender Relay

The sender relay mote receives messages from the base on the control port and retransmits them to the sender mote.

G. Receiver Relay

The receiver relay mote receives messages from receiver on the record port, and retransmits them to the base station mote.

H. Miscellaneous Network Problems and Solutions

The existence of relay motes was an unanticipated need for this experiment. We quickly learned that plugging the base station mote into the programming board changes its communication characteristics such that it cannot communicate beyond about 30 feet. We decided to create one relay mote that could be moved easily to allow for longer communication between the base station and the other motes. This still did not fix the problem, as we would frequently experience cases where the relay would have communication to either the sender or the receiver, but not both. We therefore settled on the final ring-like setup where each mote that must communicate with the computer has its own relay.

In the end, our final algorithm had two main problems. (1) $ACK's$ were still lost, causing some deadlock situations between the sender and the recevier. After resetting the two motes, the problem would be temporarily fixed. (2) The receiver and receiver relay would inexplicably enter states where both were trying to send but neither was really doing anything. If we waited for about one minute, then the transmission would eventually go through and everything would resume. Odd things would happen at this point on the computer, such as multiple receptions of the same packet from the receiver relay, and some packets that were sent were never received. After resetting the two motes and repeating the sample, everything usually worked correctly.

The lesson is that any algorithm we write in the future will have a built-in reset message type that will instantly reset all motes in the network, because manually resetting everything when the motes are 150 feet apart wastes a lot of time.

I. K-Nearest-Neighbor Analysis

We used Matlab to read all of the data files and separate them into 255 dimension samples, where each transmitter power level is one dimension. A summary of the amount of data taken for each environment appears in Table I. We then converted the RSSI readings from their raw ADC values to a more meaningful dB value that factors in the receiver battery power level using a formula provided in [7].

The squared Euclidean distance from each sample's RSSI values to all other samples was computed and sorted for each sample in order to find the first K nearest neighbors in signal space, for K from 1 to 100. We combined the PSD of the K nearest neighbors as follows:

where w_k is the weighting factor for neighbor k , s_k is the SSD of neighbor k, \hat{d} is the kNN distance estimate, d_k is the PSD of neighbor k , $E_{ranging}$ is the distance estimation error, and d is the actual PSD of the data sample. We calculated percentiles of the distance estimation error and plotted them as a function of K in order to find the best K value.

Using the best K-value from above, we looked at two different scenarios: where the kNN from each sample were computed only among the other samples from the same environment (self-self), and where the kNN for each sample were computed using samples from all environments (selfall). These curves were plotted for each environment, and then the mean and standard deviation of these curves were plotted against each other.

After plotting the CDF's for each environment, it became clear that the environments where the maximum PSD was small performed artificially well; their maximum possible error was much smaller than that of the other environments. To make a valid comparison, we normalized each of the CDF curves by the maximum PSD of each environment.

J. Localization Algorithm

Using the kNN distance estimates, we then separated the data in each file into sets of samples with the same PSD. Then, generating a random permutation of the available PSD's in each environment, we selected one of the samples at random for each of the distances in the permuation.

In order to simulate an environment where stationary beacon motes are transmitting to an unfixed mote, we placed the unfixed mote at the origin (0,0) of a 2-D plane. Then, each of the samples chosen from the random permutation were assigned a uniformly random number from zero to 2π . The beacon motes were then placed at their PSD from the origin but with with the random polar angle.

TABLE I AMOUNT OF DATA PER ENVIRONMENT

Environment	Max. Distance	RSSI Measurements	255-D Samples
	146	149,960	710
В	144	146,426	720
C	150	153,906	747
D	40	36,833	181
E	146	141,055	728
F	68	56,530	331
All	150	684,710	3417

The resulting setup was fed into the function:

where N is the number of beacon motes, x_i and y_i are the known x and y coordinates of beacon i, \hat{x} and \hat{y} are the unknown x and y coordinates of the unfixed mote, and d_i is the kNN estimated distance from beacon i to the unfixed mote.

The function f is then minimized over \hat{x} and \hat{y} to result in an estimate of the unfixed mote's location. The location error was computed as: $E_{localization} = \sqrt{\hat{x}^2 + \hat{y}^2}$ Note that the actual coordinates of the unfixed mote do not appear in the $E_{localization}$ equation because they are $(0,0)$.

IV. RESULTS

Our final data set contained 3417 samples of 255 RSSI values from 6 environments. We took 10 samples at each PSD except for the final distance where communication failed before 10 samples could be taken. The indoor, line-of-sight environments allowed communication nearing 150 feet due to waveguide effects, whereas the outdoor and non-line-ofsight environments lost communication at 68 and 40 feet. A summary of the amount of data taken is given in Table I.

Figure 3 shows the results of our attempts to find the best value for K. \overline{As} expected, the 50th, 80th, and 90th percentile curves all start to approach a constant value around $K = 9$ or 10. Our weighting method weights the neighbors that are closer in signal space more heavily than those far away. As K increases, we are adding more neighbors that are farther away, so their contribution to the end result is smaller and smaller. [3] used a simple direct averaging weighting scheme, and they saw heavy degradation in performance as K became large. We plotted both schemes, and they performed almost identically for small K values, but the final scheme we used was clearly superior for larger values of K. We chose $K = 9$ for the rest of our analysis because it is near the steady-state values, but is also small enough to speed up computation.

Using $K = 9$, we plotted the CDF of the distance estimation error for the self-self case (Figure 4) and the self-all case (Figure 5) for each environment. For the self-self case, each

Fig. 3. 50th, 80th, and 90th Percentile Values vs. K for Self-All Case

Fig. 4. CDF of Distance Estimation Errors for Self-Self Scenario

curve was normalized by the maximum PSD for that environment, so that the x-axis of the CDF plots are actually in units of percentage of maximum communication distance. In other words, if the desired environment allows for communication up to 100 feet, then the x-axis should be scaled by multiplying by 100. Note that this is only valid for ranges on the order of those studied here.

The self-self plot resembles the results of [2], as expected. The open field and milking parlor were the best environments for this case. This is probably because they were both static – no moving people, animals or objects during the experiment. It is also possible that our method of dividing the error by the maximum PSD in the set does not fully account for the problem of smaller maximum possible error.

The self-all plot, however, shows the open field and milking parlor environments to be by far the worst of the environments. The smaller maximum PSD in each set is probably leading to

Fig. 5. CDF of Distance Estimation Errors for Self-All Scenario

artificially poor performance by including nearest neighbors with PSD's that are outside the realm of possibility for that environment.

In order to better investigate the degradation in performance due to using a generic set of samples, we computed the mean and standard deviation of the curves in both cases and plotted them in Figure 6.

Fig. 6. Mean CDF +/- 1 Std of Distance Estimation Error Across Environments

It is obvious from Figure 6 that the self-self case is better than the self-all case in terms of accuracy, but the standard deviation of the self-self case is noticeably larger than that of the self-all case.

If we assume that a standard environment allows for communication up to 150 feet, then we expect to be able to predict the PSD between two sensor motes with 80%

confidence to within 30 feet if the samples are taken from the same environment, and 50 feet if the samples are taken from generic environments. The self-self case has an 80th percentile standard deviation of 16.1 feet, and the self-all case has an 80th percentile standard deviation of 9.8 feet, an improvement of almost 40%.

At first, the improvement in precision seems counterintuitive. To understand this result, one must think further about what it means to have a smaller standard deviation in the context of a kNN estimate. If we look at a sample (sample A) with a PSD of 6 feet, for instance, we will get one estimate of the distance based on the other samples around it. If we have another sample (sample B) at 6 feet that is relatively close in signal space to sample A, then B's estimate will be "far" from A's only if a small change in signal space location results in a very different set of neighbors. "Different" in the sense that the new neighbors have much different PSD's than those of the neighbors of the first sample. Therefore, a highly-irregularly populated signal space will lead to a high standard deviation. We expect that the overall signal space of the self-all case will be more regularly populated than the self-self case due to the much larger number of samples from somewhat similar sets.

Fig. 7. Mean CDF +/- 1 Std of Localization Error for 3 and 25 Beacons

Figure 7 shows the mean localization error CDF $+/- 1$ standard deviation across all environments. Comparing this plot to the ranging error in Figure 6, one can see that the localization CDF is much steeper than that of the ranging algorithm. This is because the random placement of beacon motes allows for some ranging errors to effectively cancel out other ranging errors. However, this results in a much larger standard deviation than with ranging. The mean 90^{th} percentile is smaller using localization, but the 80^{th} is larger. The mean 80th percentile error was 36.97 feet, and the standard deviation was 17.95 feet for 25 beacons across all environments.

The effect of the number of beacons on accuracy is plotted in Figure 8. The open field and milking parlor error estimates

are very high, about 130% of the maximum ranging error. This is due to the poorer kNN distance approximations that can be seen in Figure 5. All of the curves seem to follow an exponential decay, showing that more beacons in general leads to better accuracy, agreeing with the results from [4].

Fig. 8. Localization Error vs. Number of Beacons

V. CONCLUSION

We conclude that received signal strength is a mediocre estimator for ranging and localization in general sparse sensor networks, with modest accuracy improvements possible if the operating environment has been pre-sampled.

VI. ACKNOWLEDGEMENTS

This research was funded in part by a Tellabs Fellowship and by the Indiana 21^{st} Century Research and Technology Fund under Award #1110030618.

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